

Data Structures – CST 201

Module ~ 5

Syllabus

■ **Sorting Techniques**

- Bubble Sort
- Selection Sort
- Insertion Sort
- Merge Sort
- Quick Sort
- Heap Sort

■ **Hashing**

- Hashing Techniques
- Collision Resolution
- Overflow handling
- Hashing functions: Mid square, Division, Folding, Digit Analysis

Bubble Sort

- In this method find the largest element in the array and this element is stored in the last position. Then find the second largest element and it is stored in the second last position, and so on

Bubble Sort

7	5	4	8	9
0	1	2	3	4

Bubble Sort

7	5	4	8	9
0	1	2	3	4

Number of Scans = $5-1 = 4$

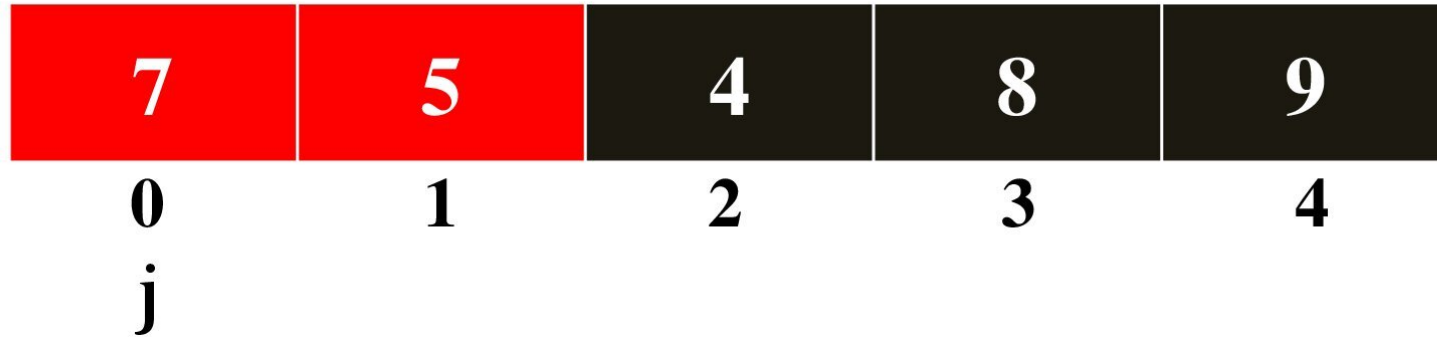
Bubble Sort



j

Scan-1

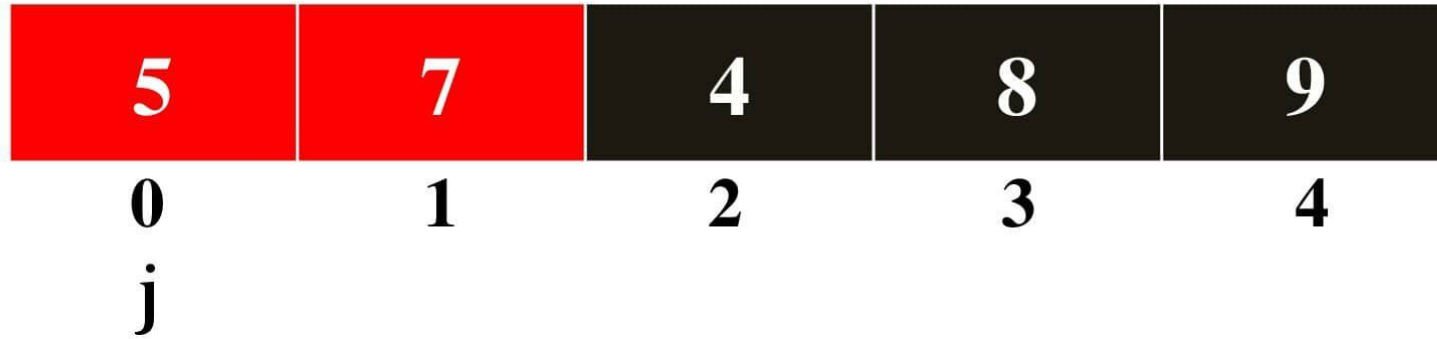
Bubble Sort



Scan-1

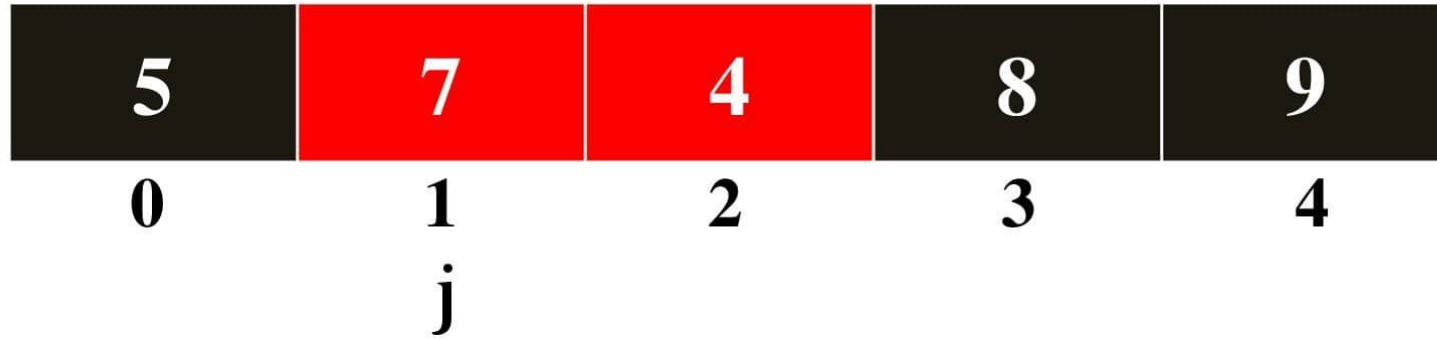
Swap the elements

Bubble Sort



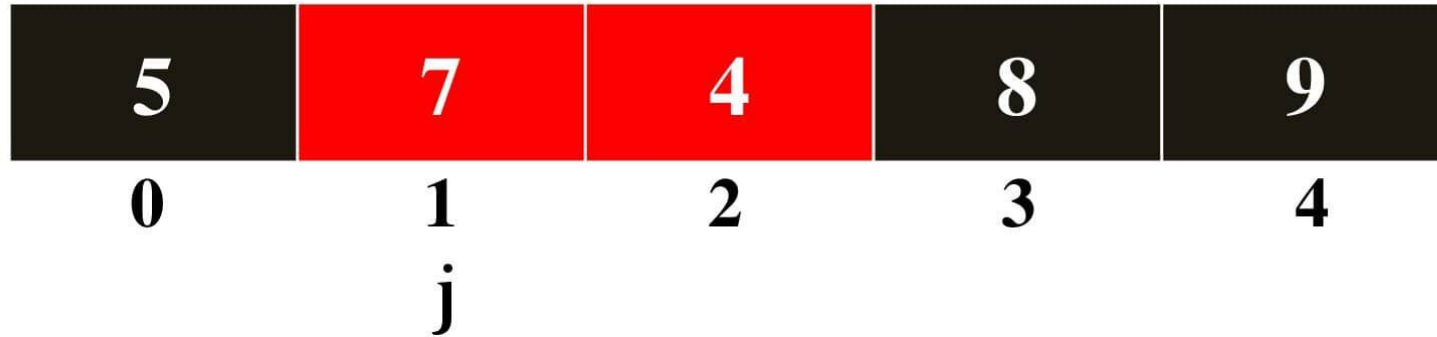
Scan-1

Bubble Sort



Scan-1

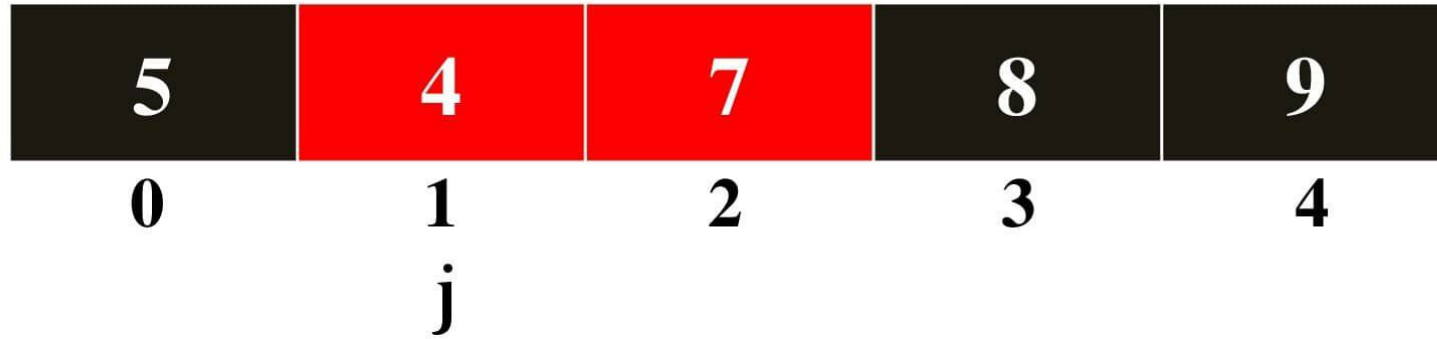
Bubble Sort



Scan-1

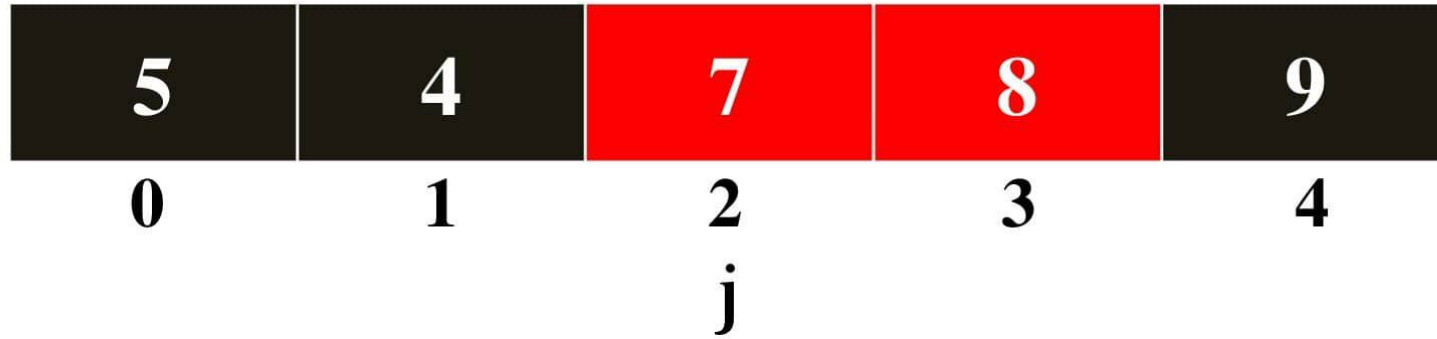
Swap the elements

Bubble Sort



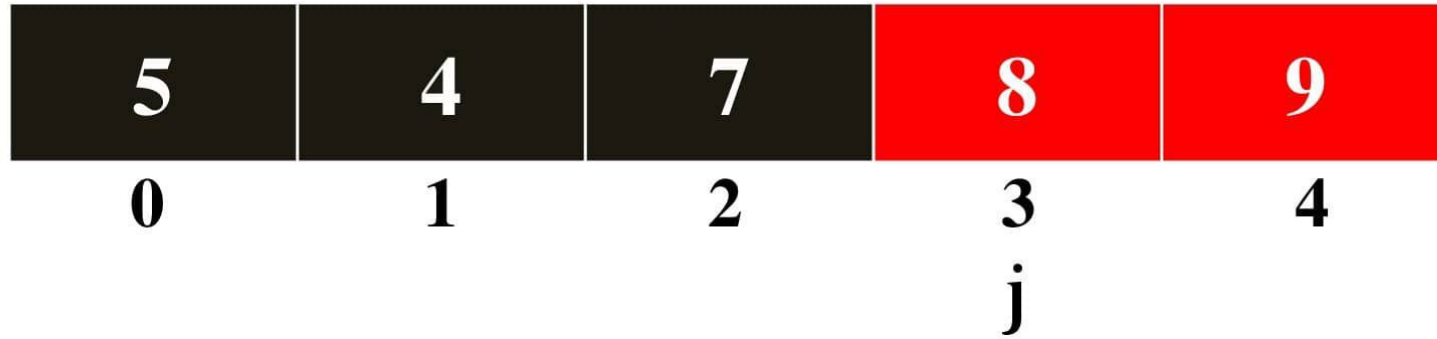
Scan-1

Bubble Sort



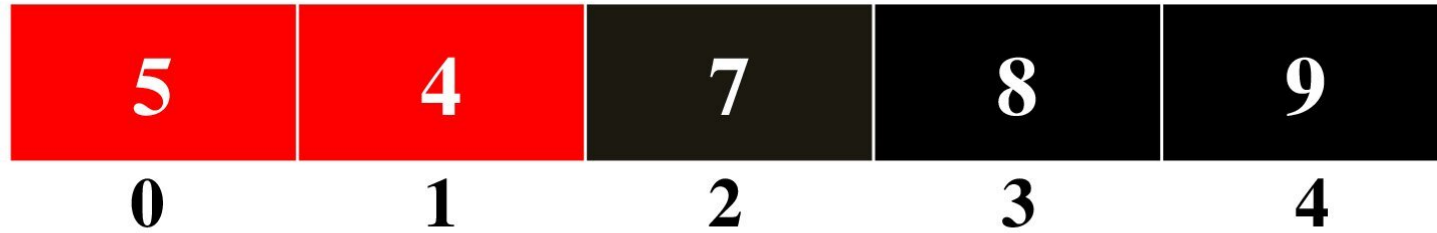
Scan-1

Bubble Sort



Scan-1

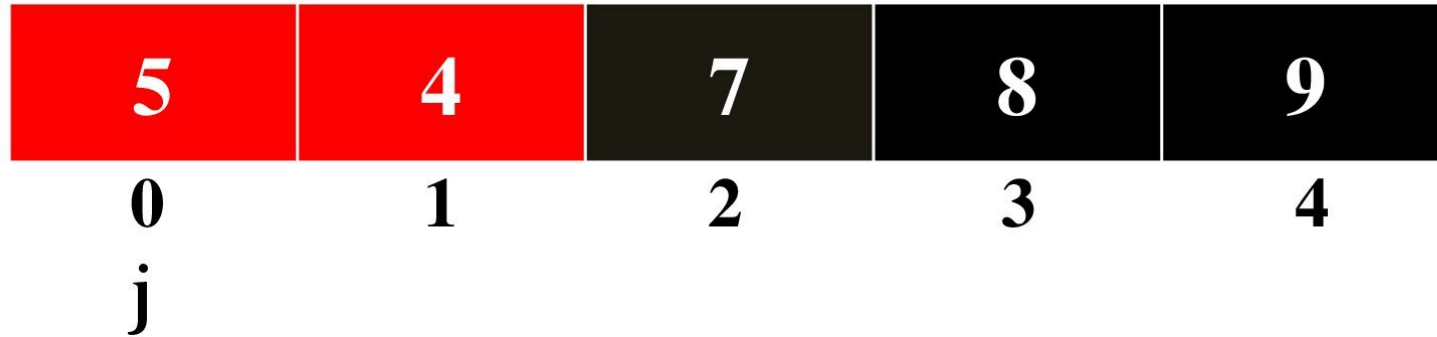
Bubble Sort



j

Scan-2

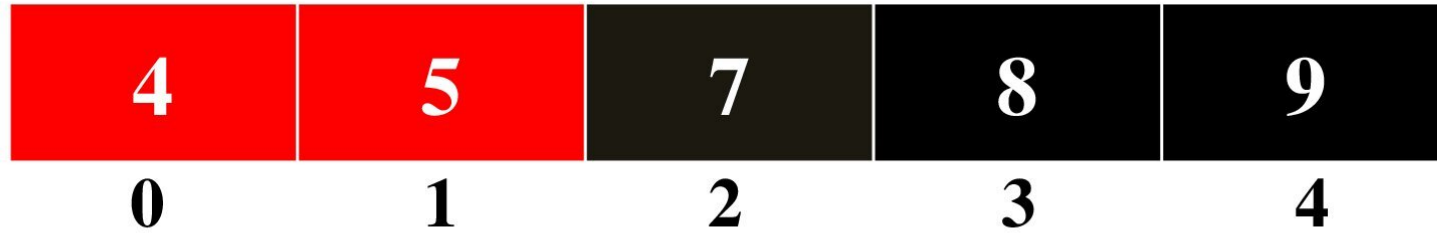
Bubble Sort



Scan-2

Swap the elements

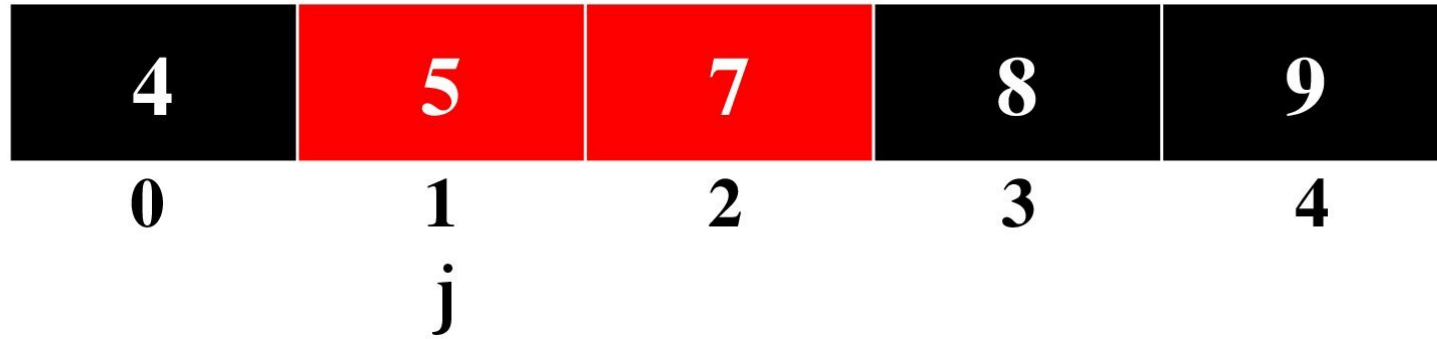
Bubble Sort



j

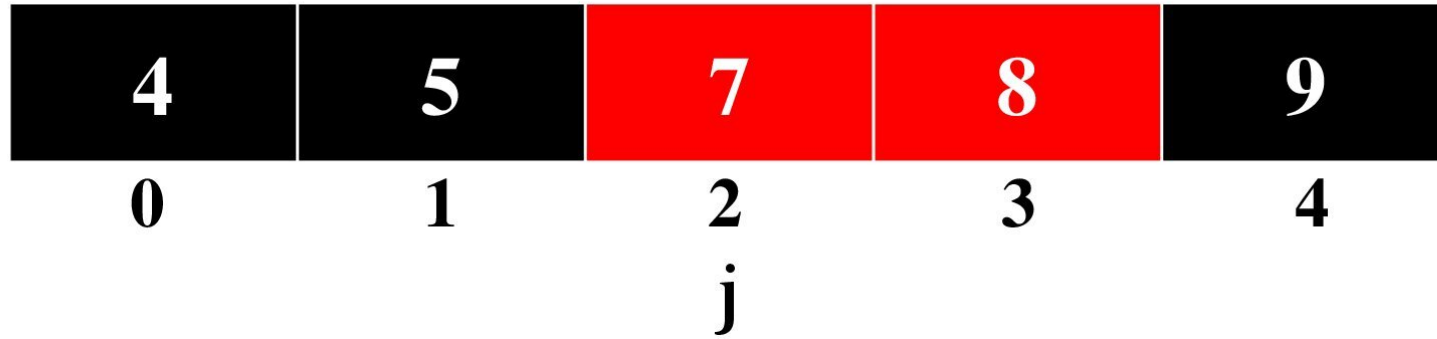
Scan-2

Bubble Sort



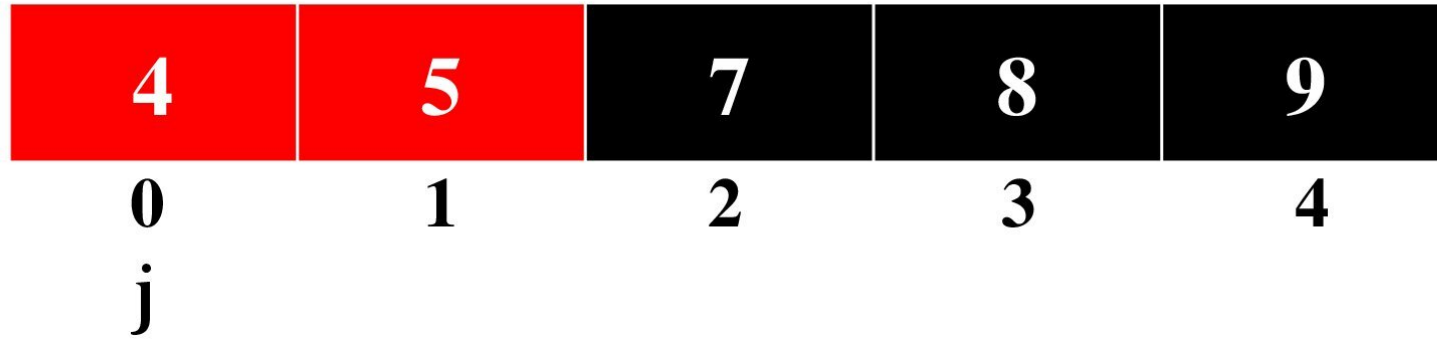
Scan-2

Bubble Sort



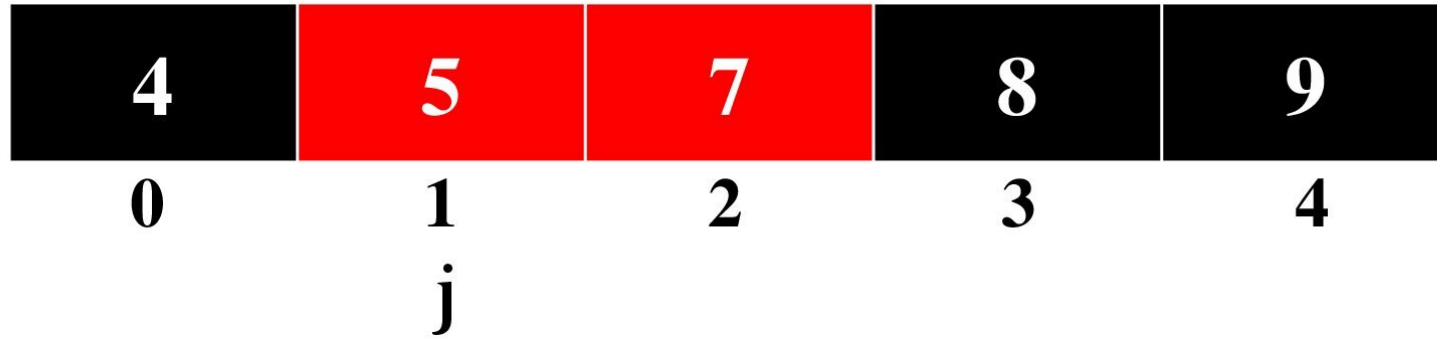
Scan-2

Bubble Sort



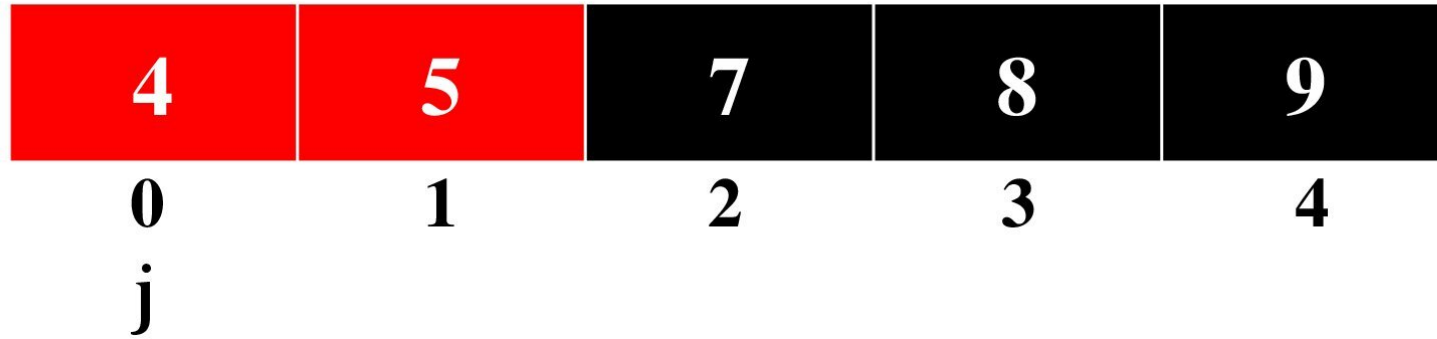
Scan-3

Bubble Sort



Scan-3

Bubble Sort



Scan-4

Bubble Sort

4	5	7	8	9
0	1	2	3	4

Final Sorted Array

Bubble Sort

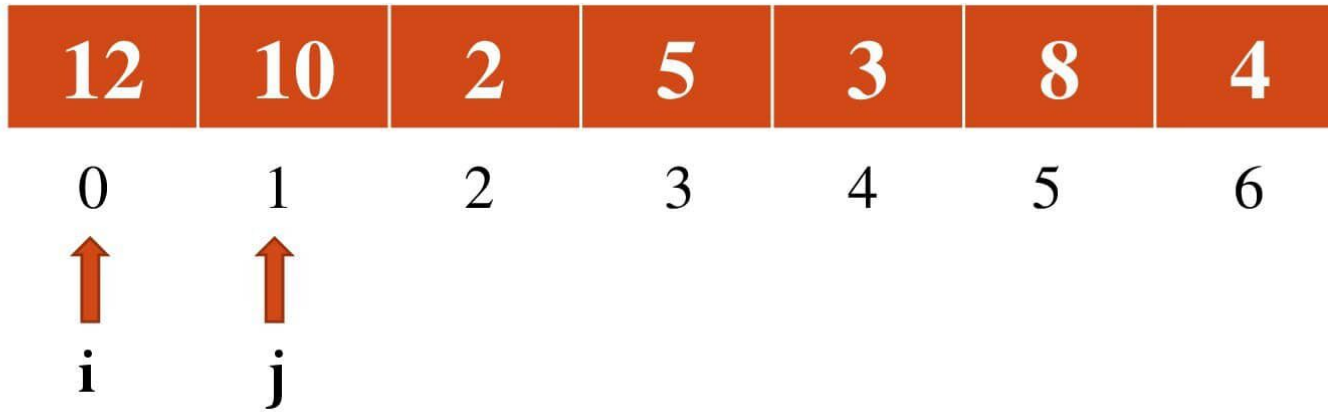
Algorithm BubbleSort(A, n)

```
{
  for i=1 to n-1 do
    {
      for j=0 to n-i-1 do
        {
          If  $A[j] > A[j+1]$  then
            Swap  $A[j]$  and  $A[j+1]$ 
        }
      }
    }
}
```

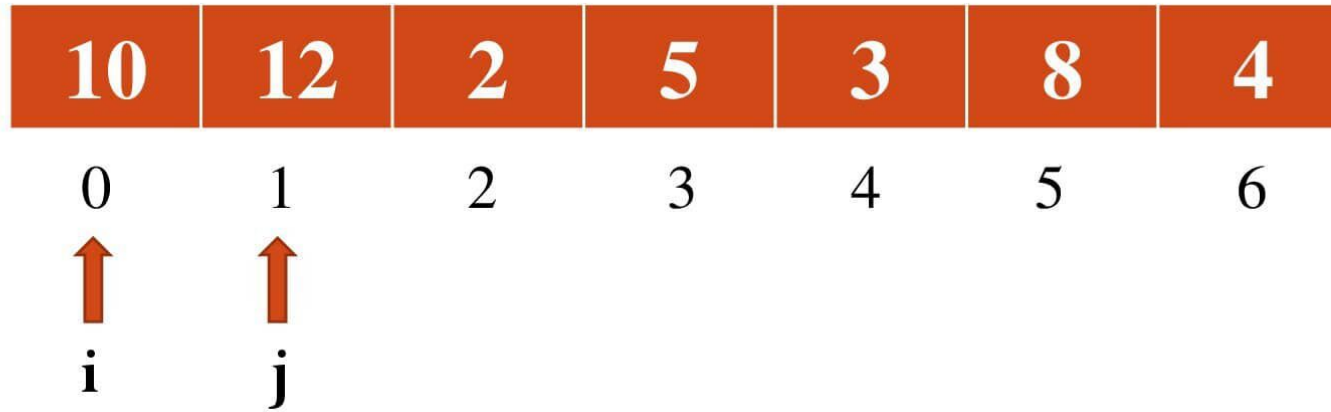
Selection Sort

- In this method find the smallest element in the array and this element is stored in the first position. Then find the second smallest element and it is stored in the second position, and so on

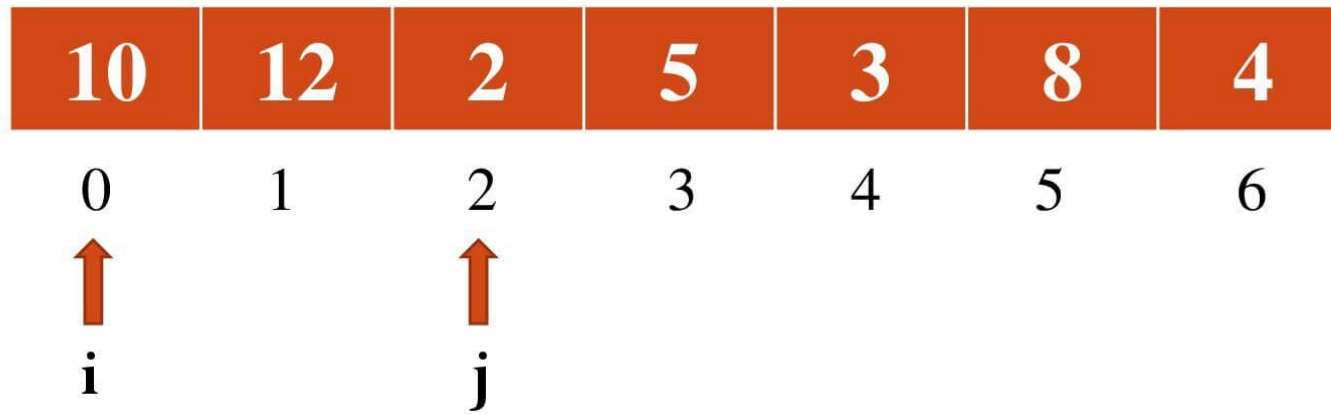
Selection Sort



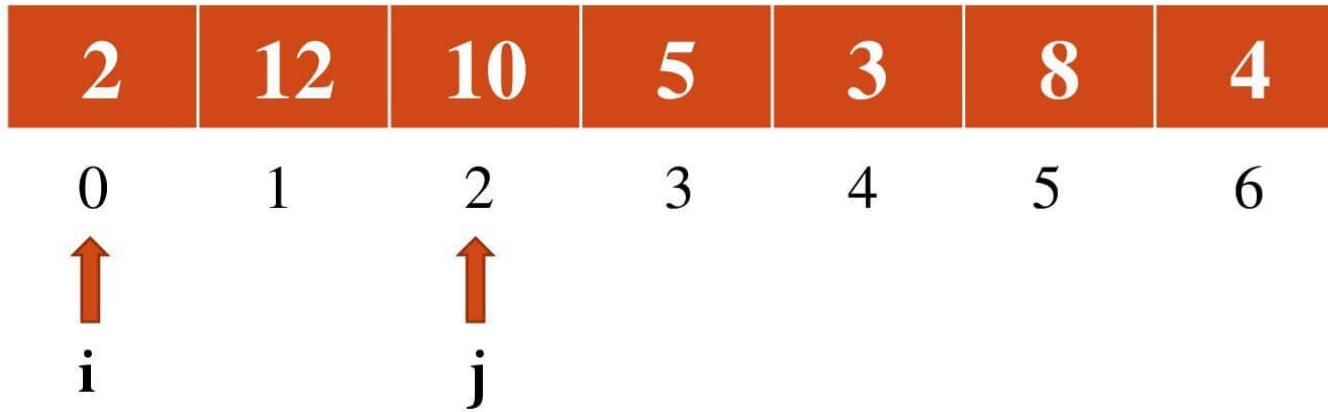
Selection Sort



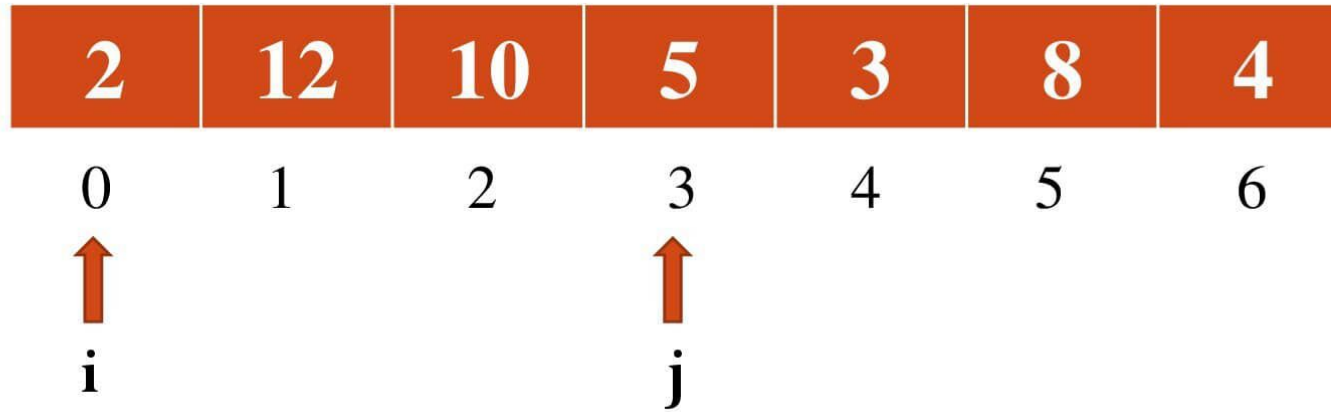
Selection Sort



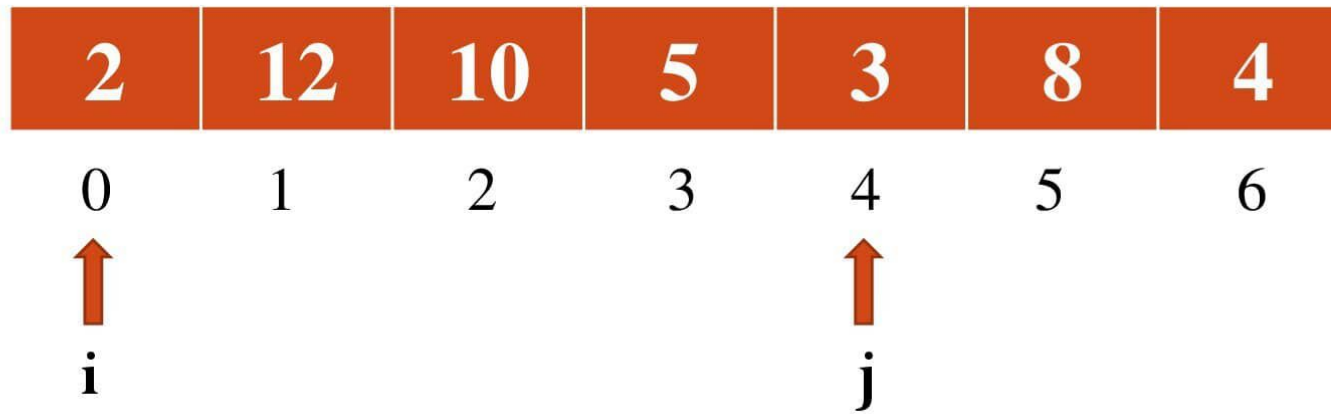
Selection Sort



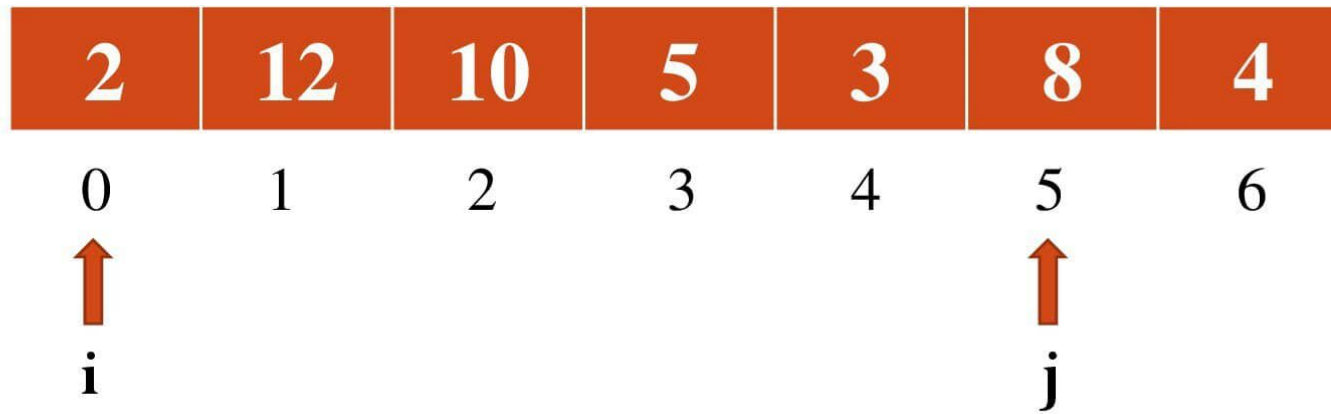
Selection Sort



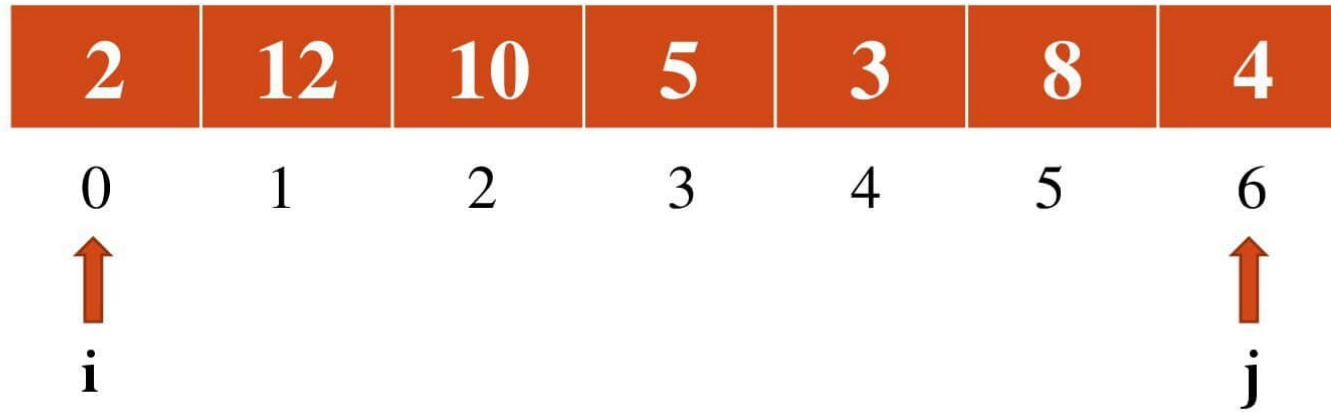
Selection Sort



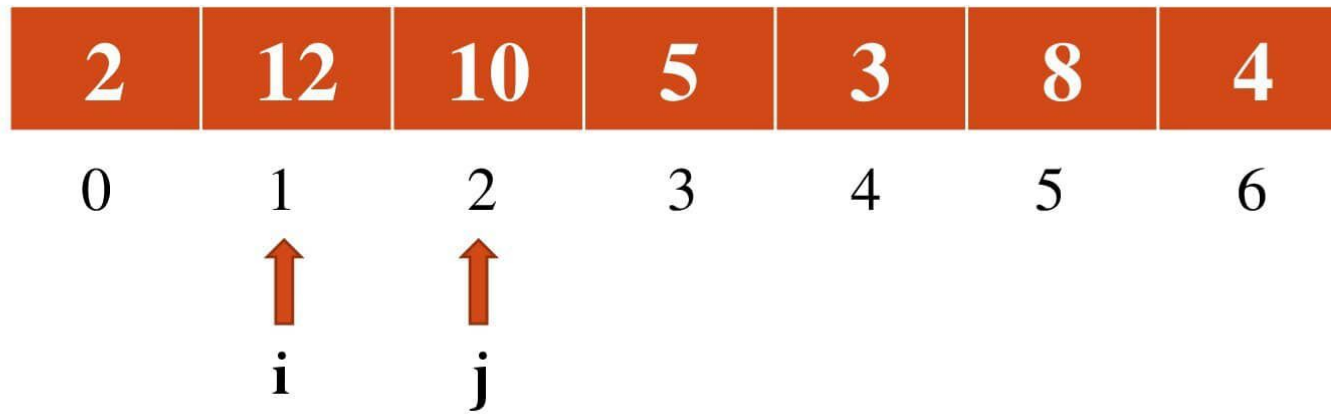
Selection Sort



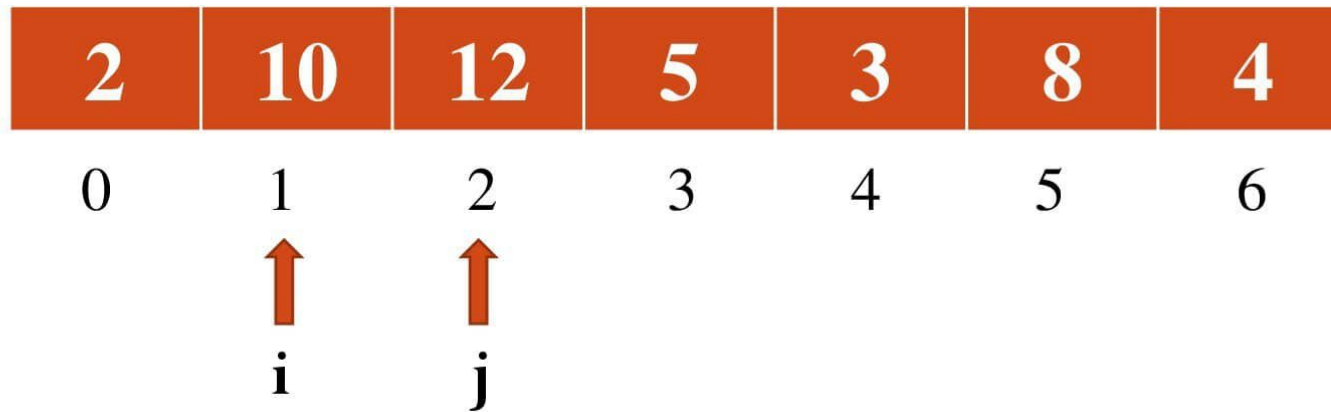
Selection Sort



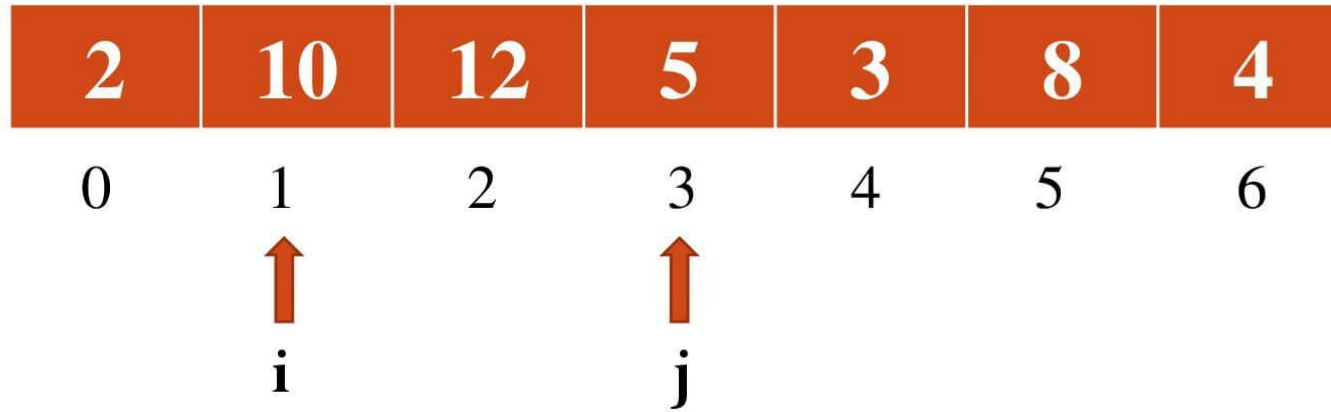
Selection Sort



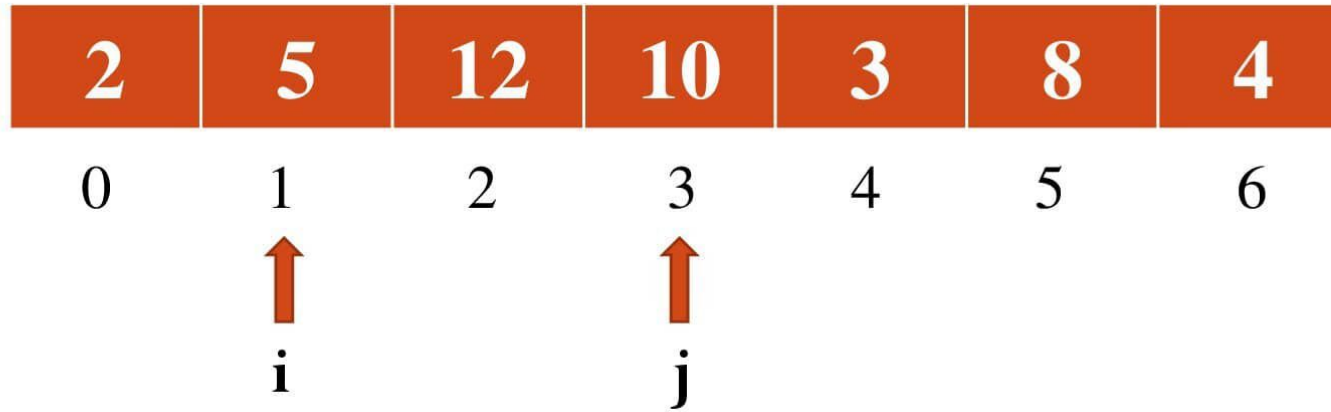
Selection Sort



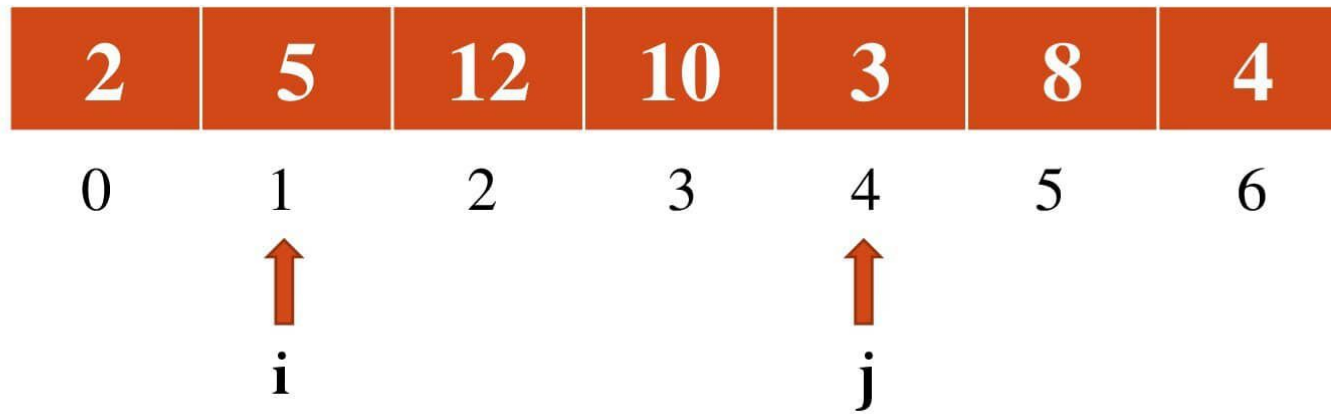
Selection Sort



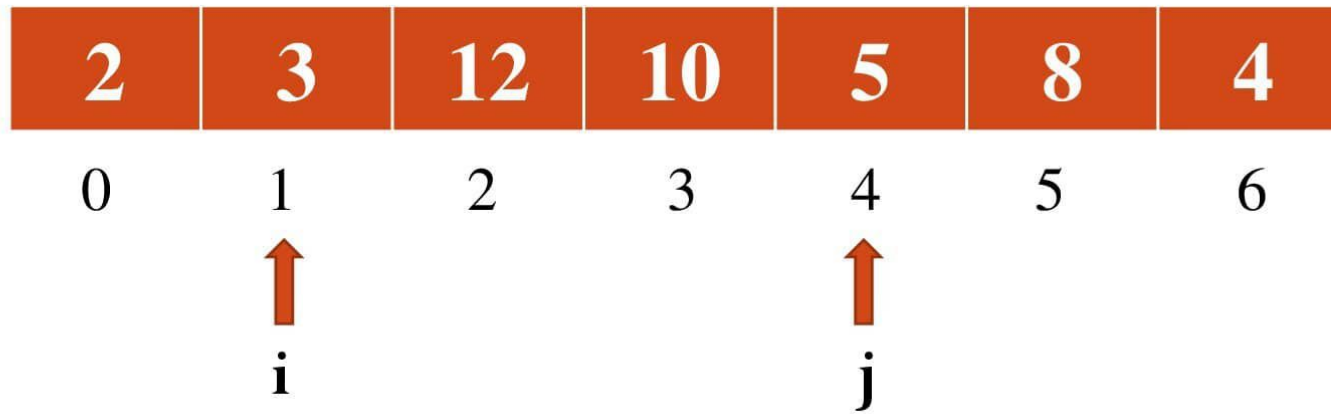
Selection Sort



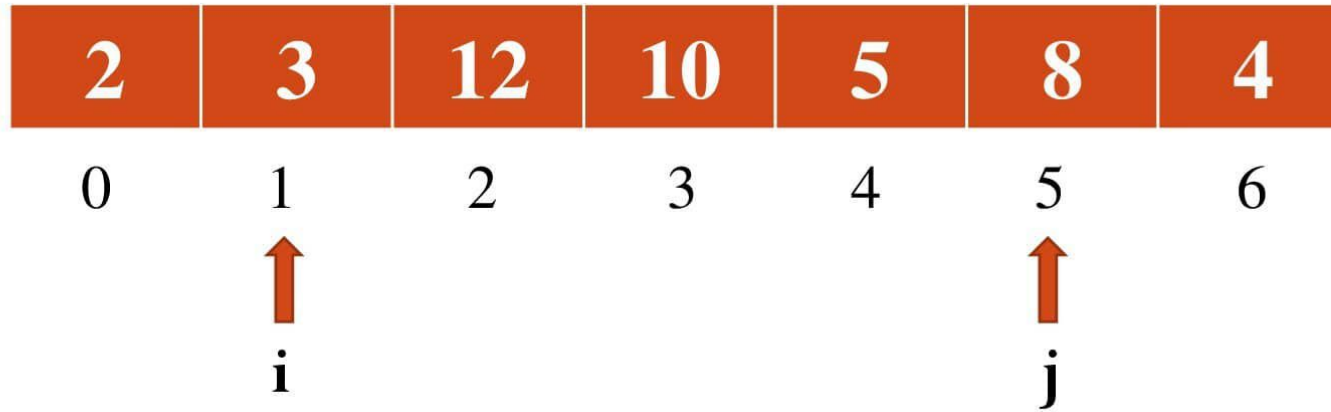
Selection Sort



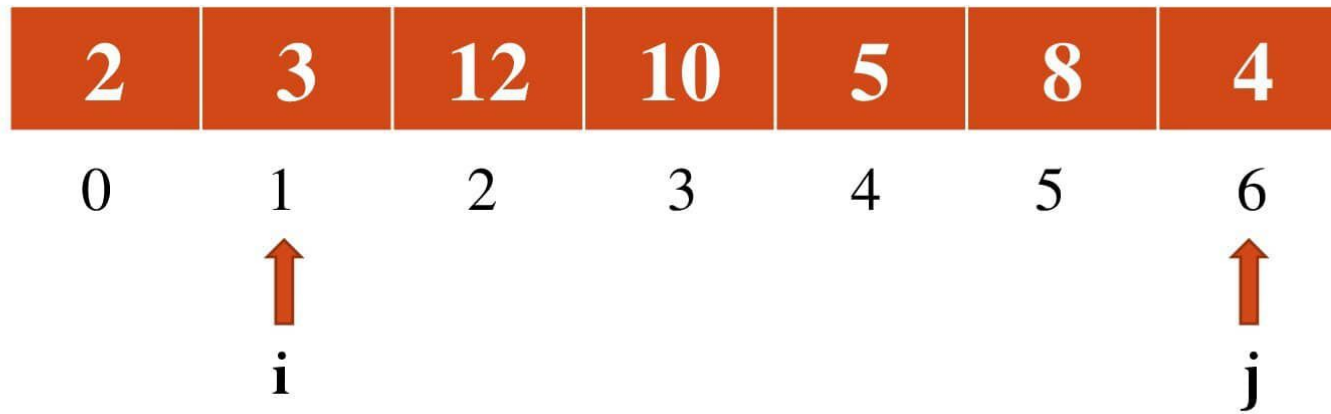
Selection Sort



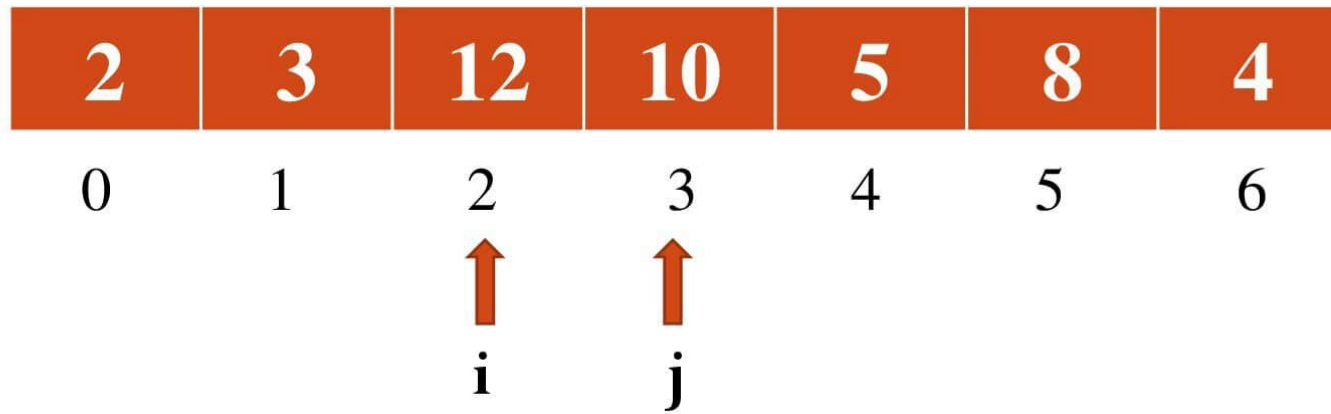
Selection Sort



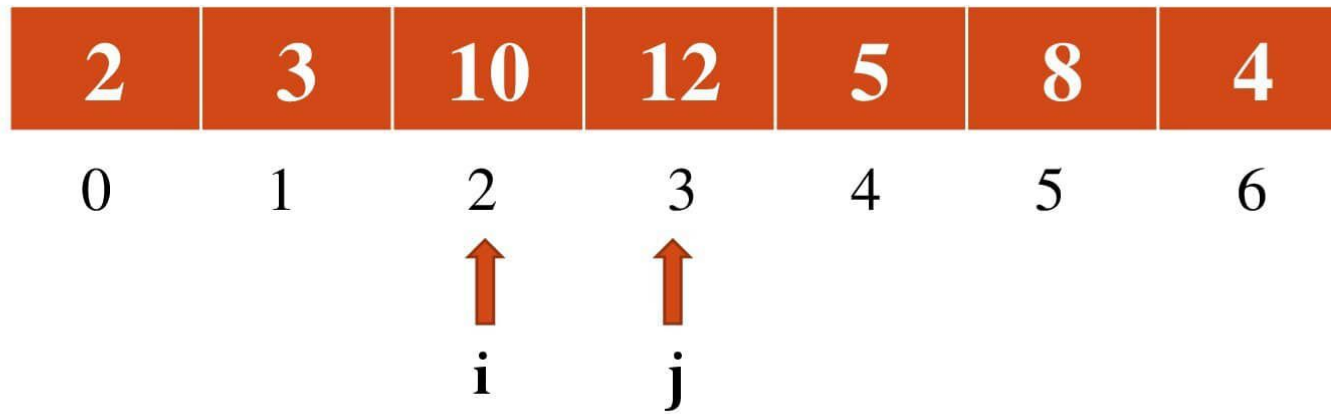
Selection Sort



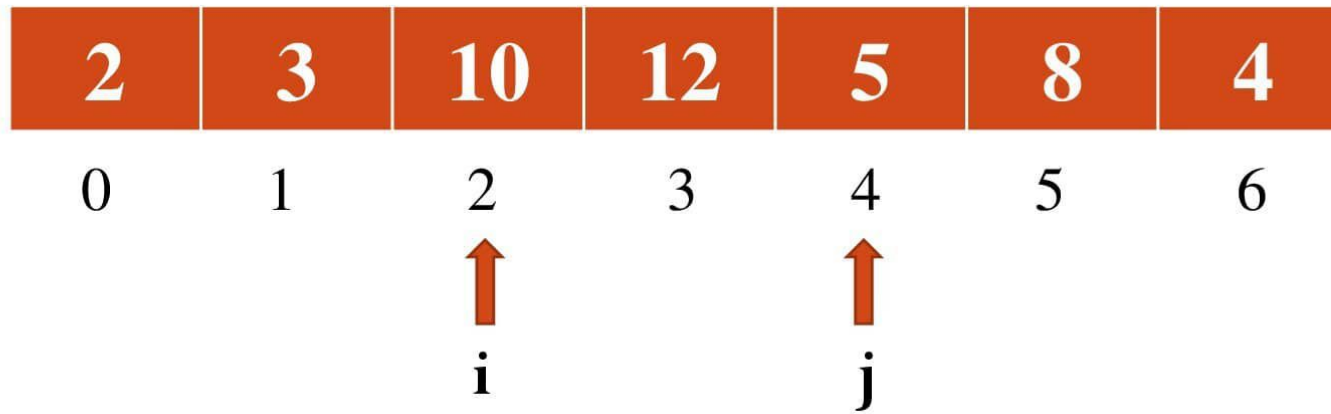
Selection Sort



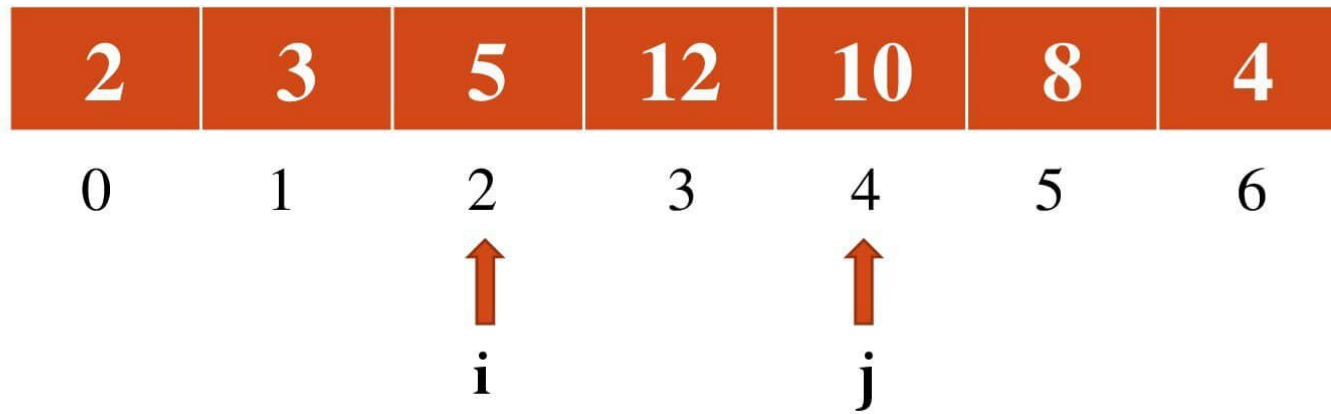
Selection Sort



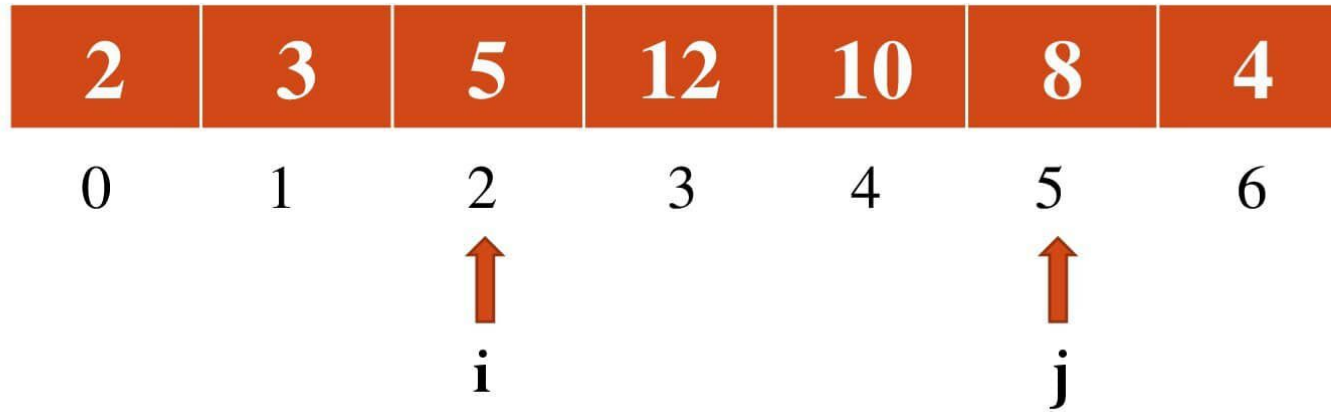
Selection Sort



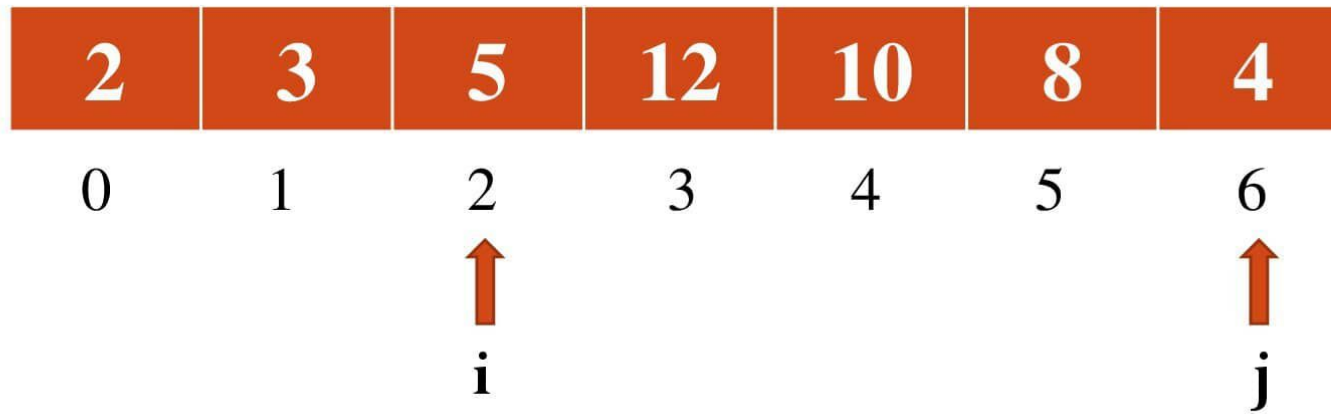
Selection Sort



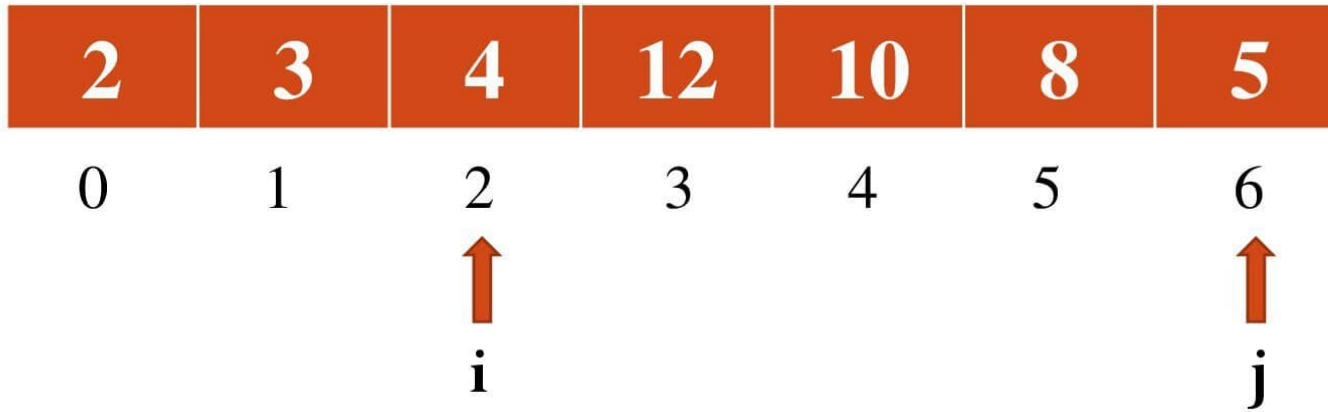
Selection Sort



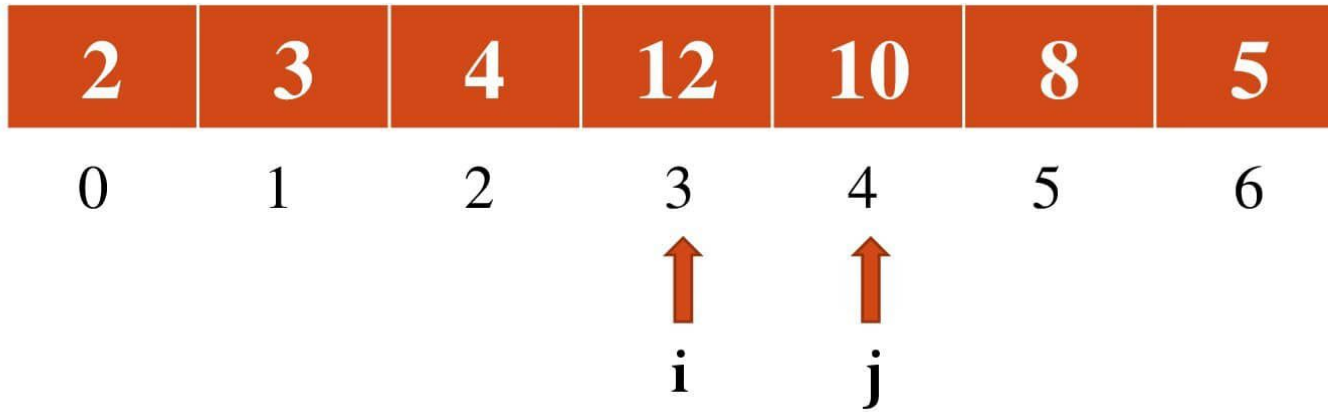
Selection Sort



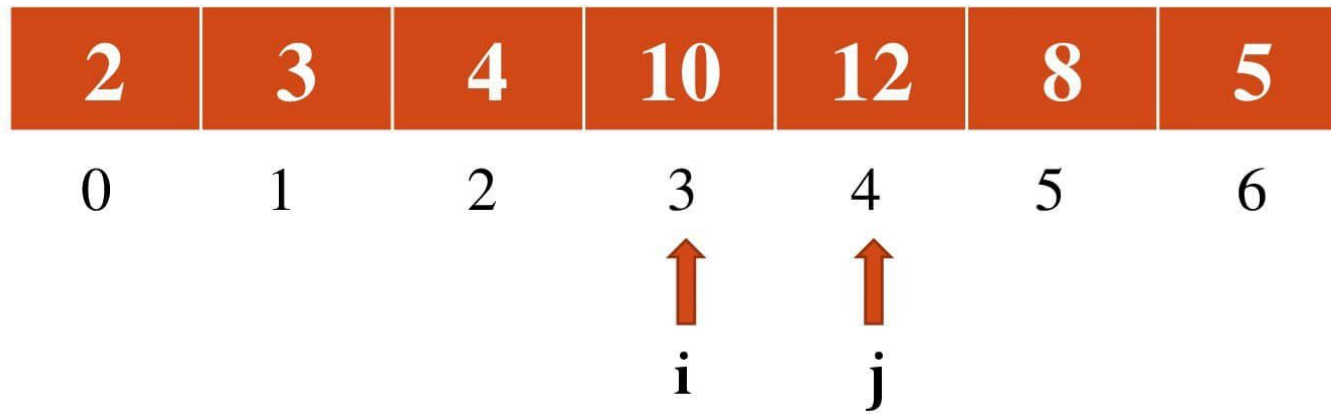
Selection Sort



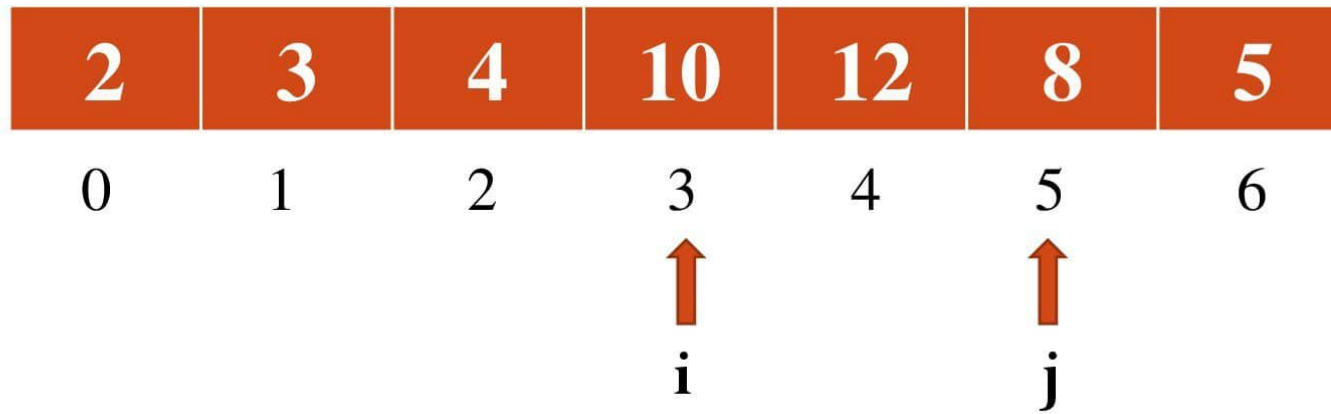
Selection Sort



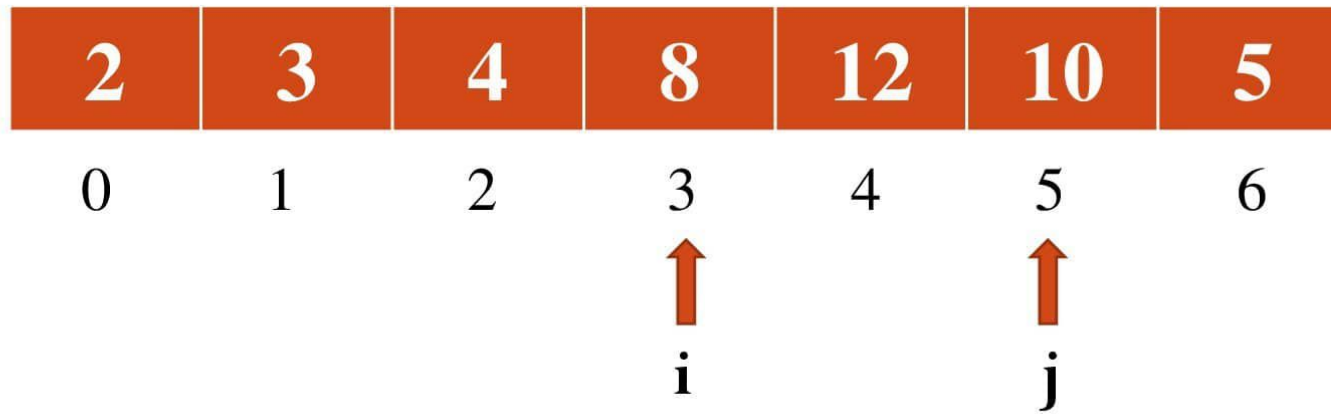
Selection Sort



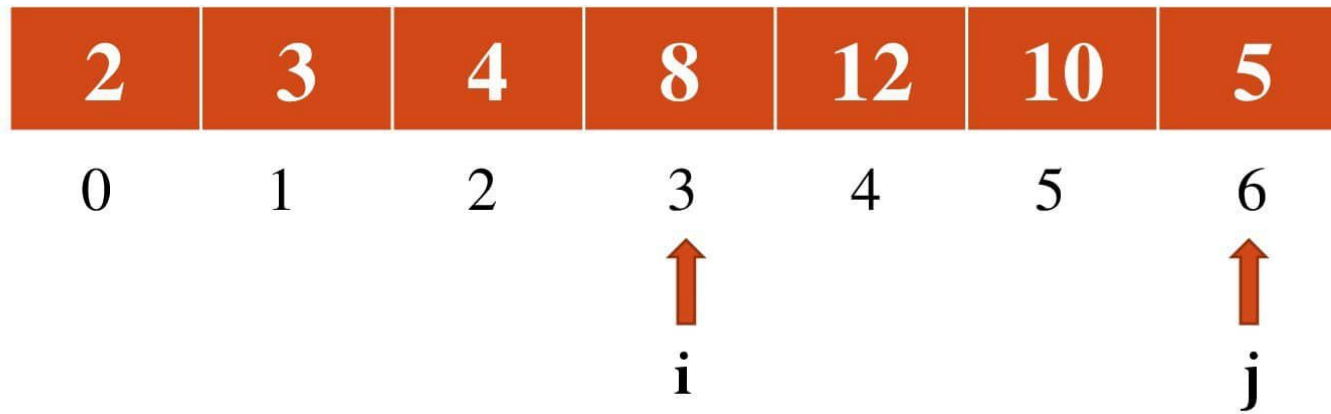
Selection Sort



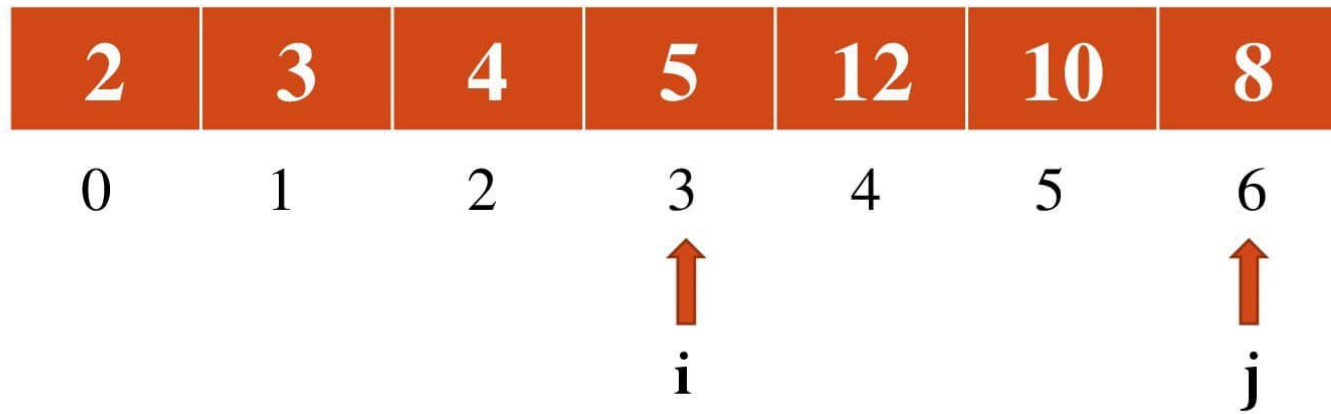
Selection Sort



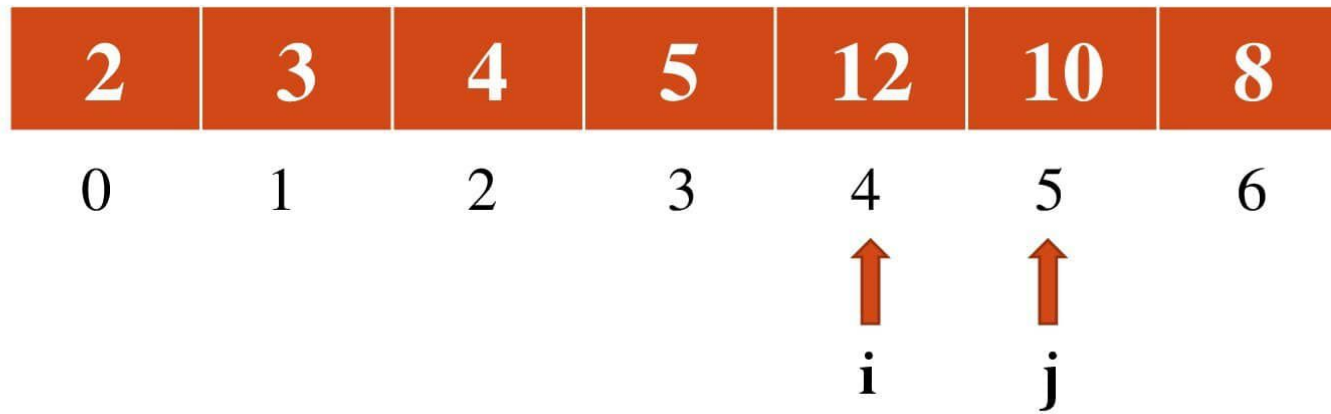
Selection Sort



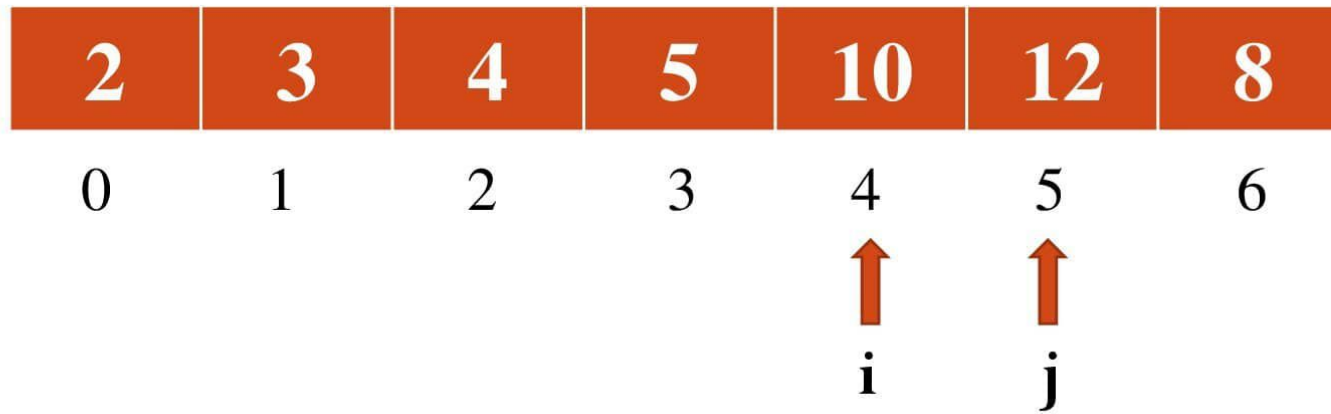
Selection Sort



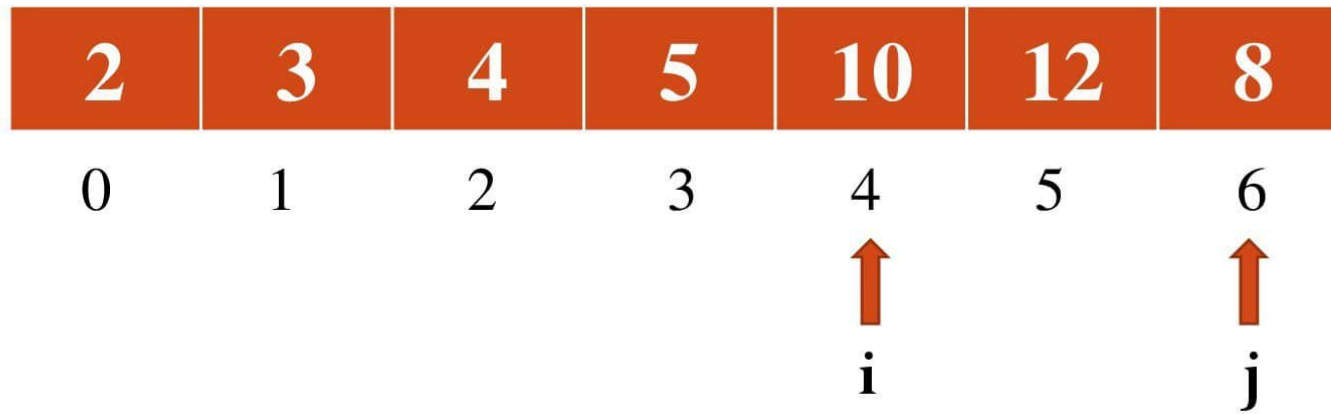
Selection Sort



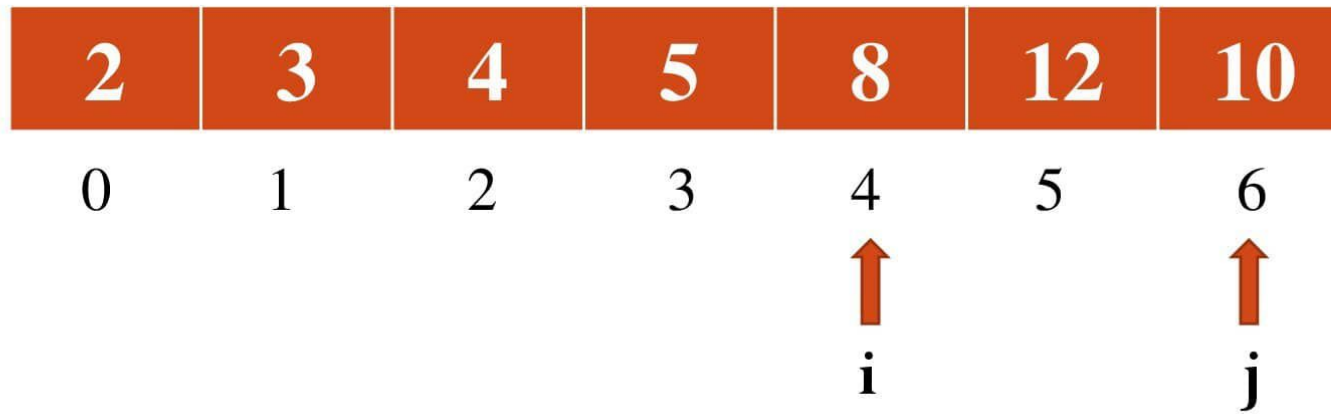
Selection Sort



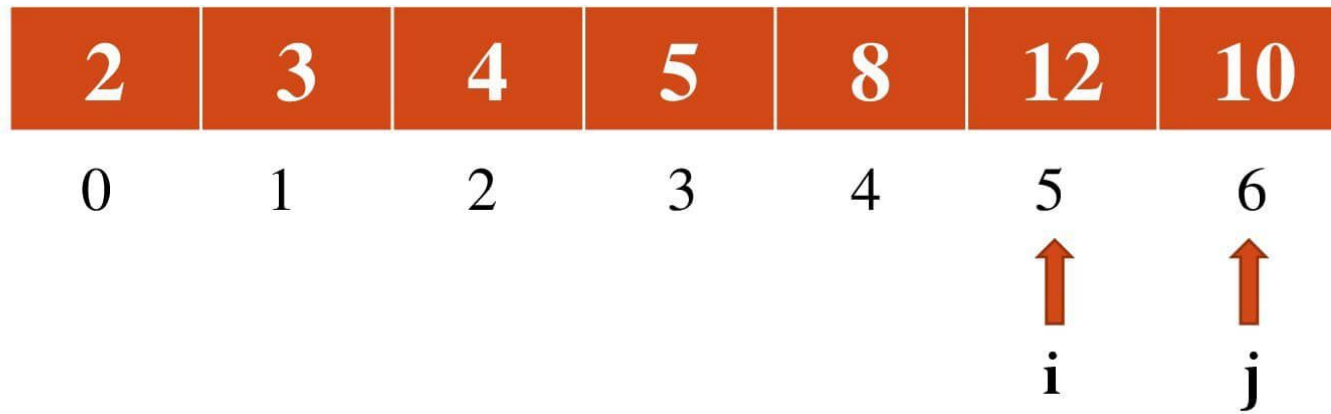
Selection Sort



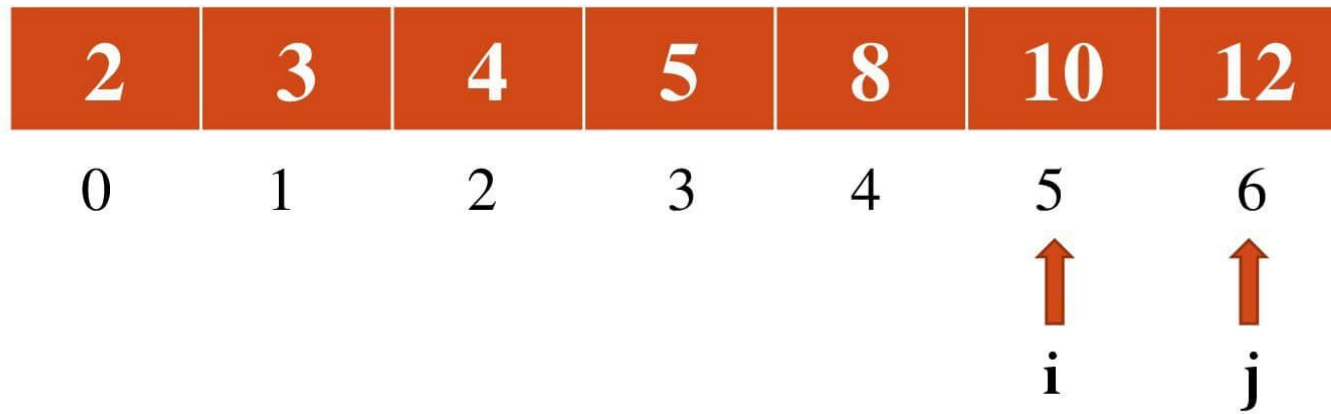
Selection Sort



Selection Sort



Selection Sort



Selection Sort

2	3	4	5	8	10	12
0	1	2	3	4	5	6

Selection Sort

Algorithm SelectionSort(A, n)

```
{
  for i=0 to n-2 do
    {
      for j=i+1 to n-1 do
        {
          If  $A[i] > A[j]$  then
            Swap  $A[i]$  and  $A[j]$ 
        }
      }
    }
}
```

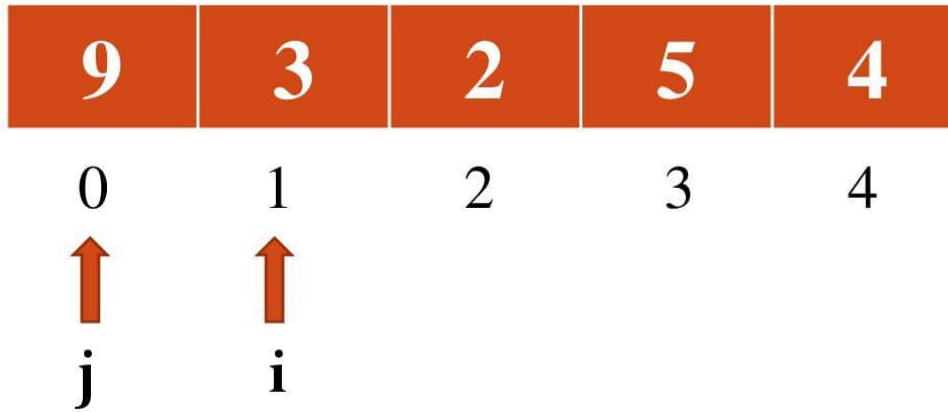
Insertion Sort

- During first iteration element in pos 1 is compared with element in pos 0 and insert it in the proper position
- During second iteration element in pos 2 is compared with elements in pos 1 and 0 and insert it in the proper position
- During third iteration element in pos 3 is compared with elements in pos 2, 1 and 0. Then insert it in the proper position
- In general in every iteration an element is compared with all elements before it and insert it in the proper position

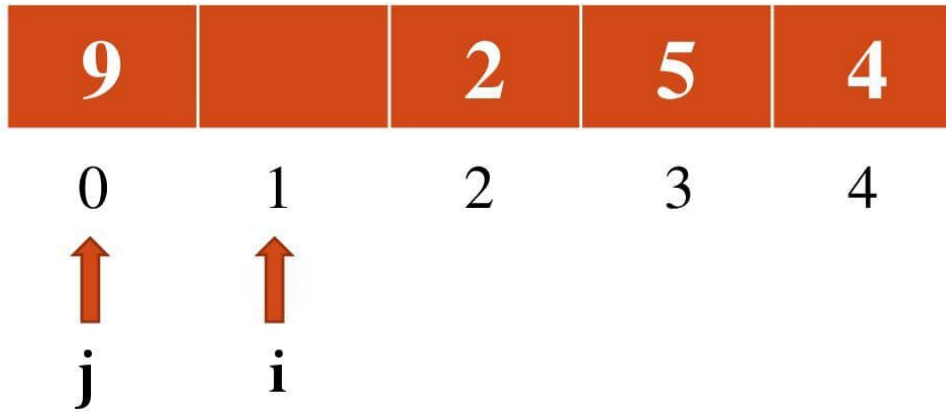
Insertion Sort

9	3	2	5	4
0	1	2	3	4

Insertion Sort

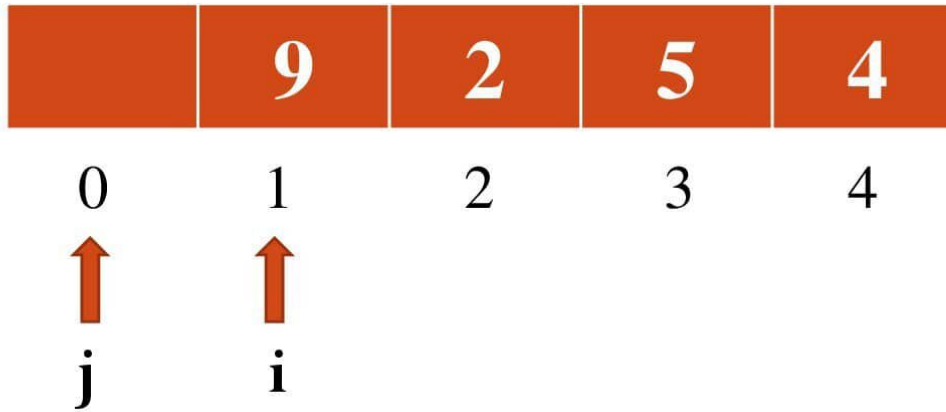


Insertion Sort



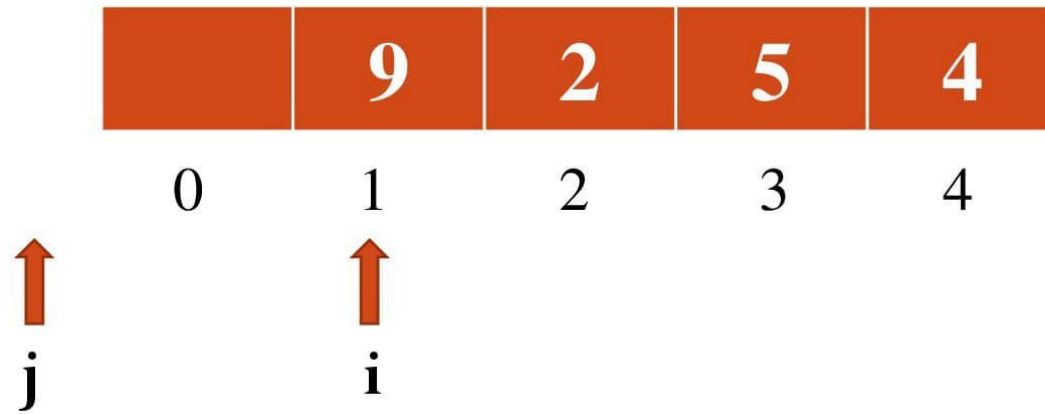
temp=3

Insertion Sort



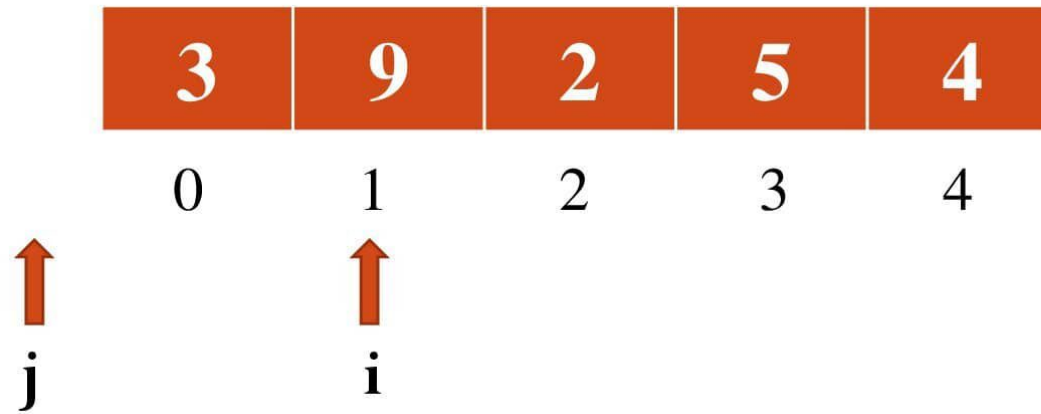
temp=3

Insertion Sort



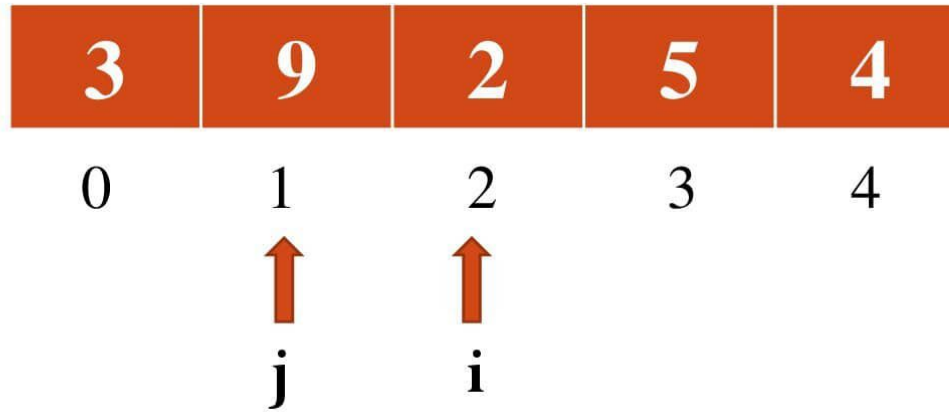
temp=3

Insertion Sort

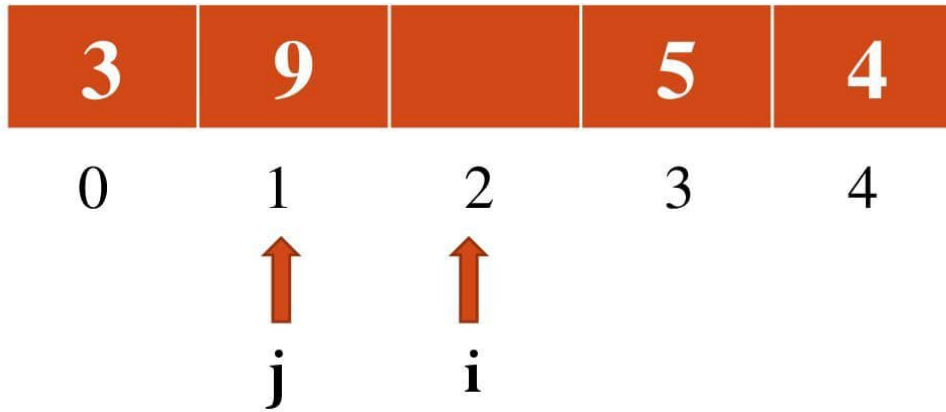


temp=3

Insertion Sort

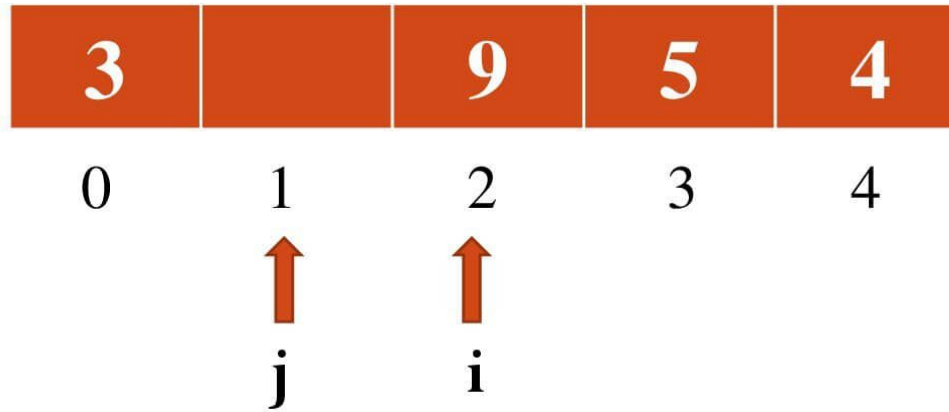


Insertion Sort



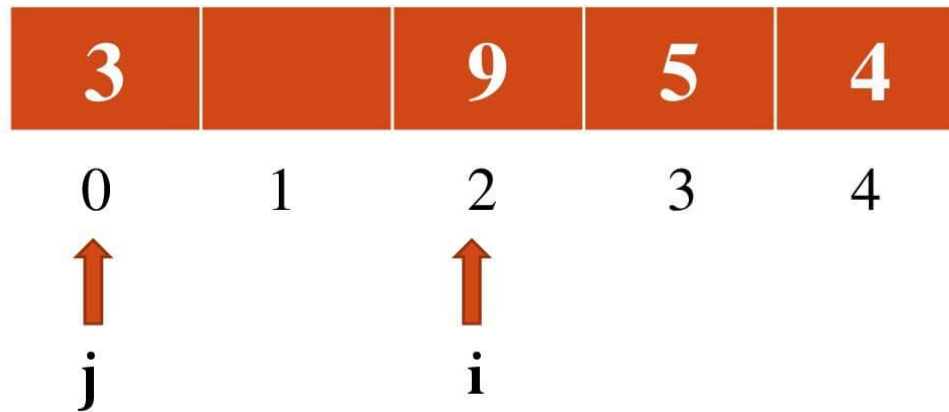
temp=2

Insertion Sort

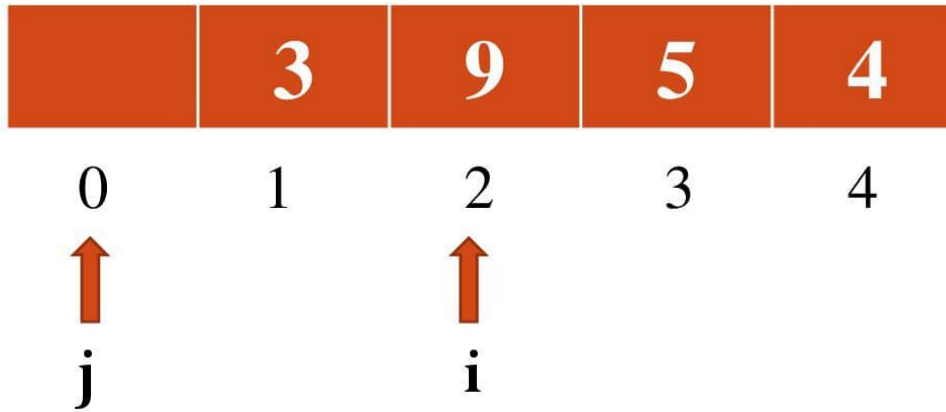


temp=2

Insertion Sort

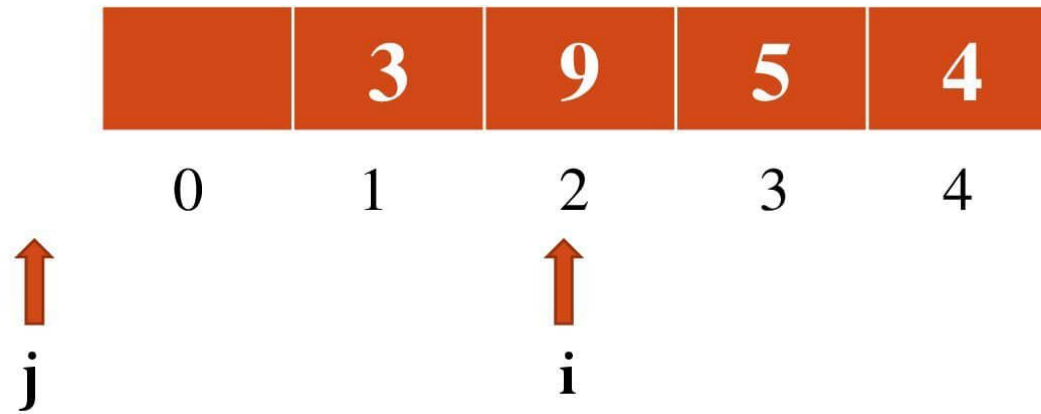


Insertion Sort



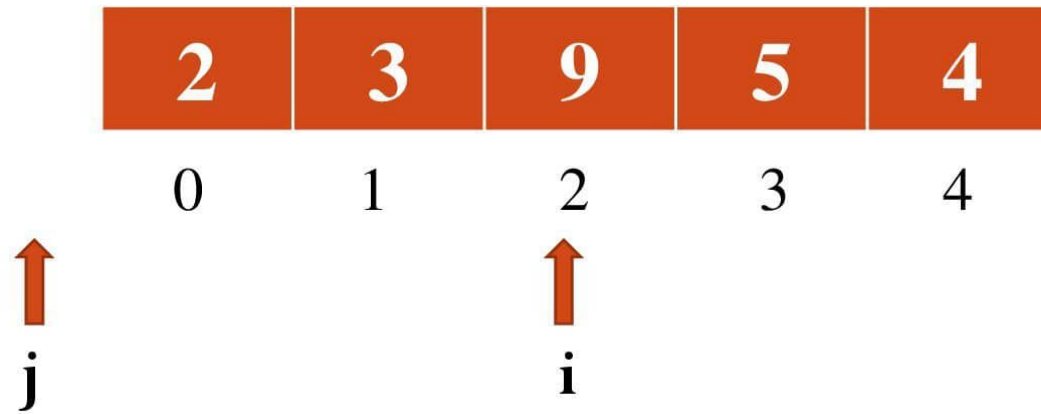
temp=2

Insertion Sort



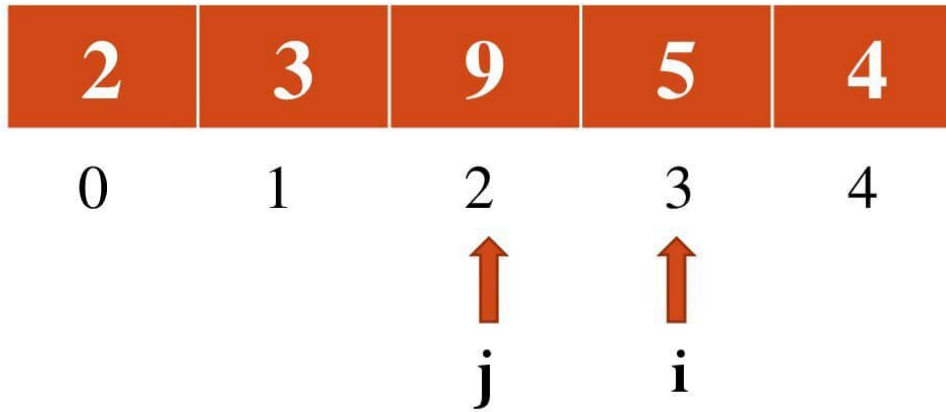
temp=2

Insertion Sort

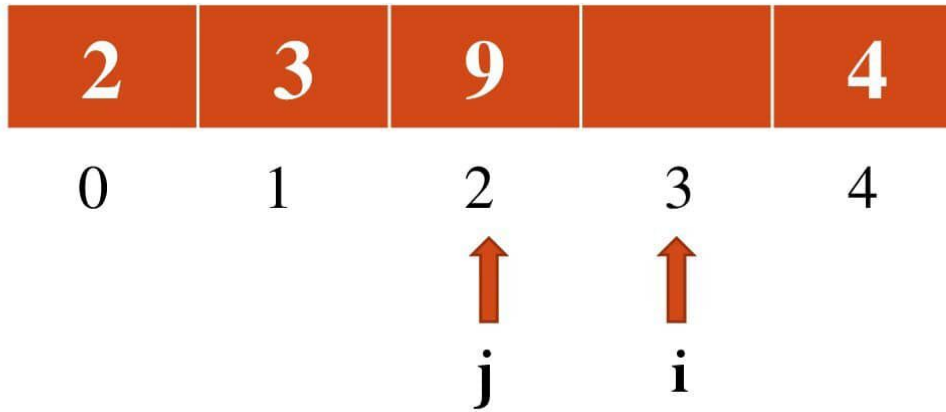


temp=2

Insertion Sort

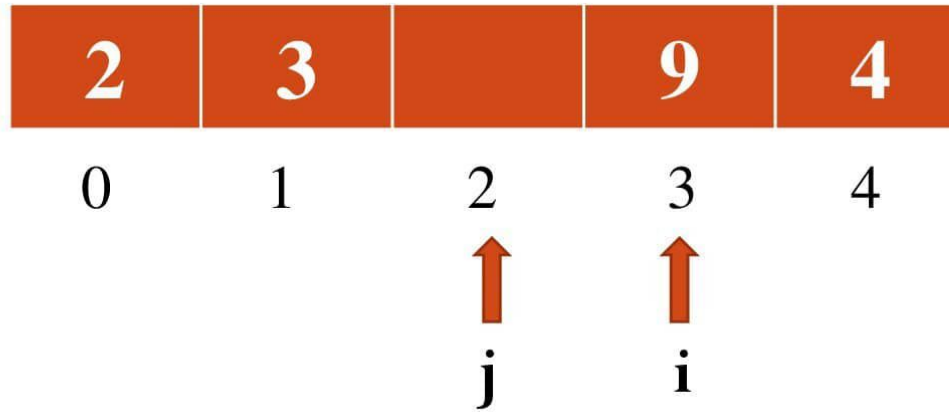


Insertion Sort



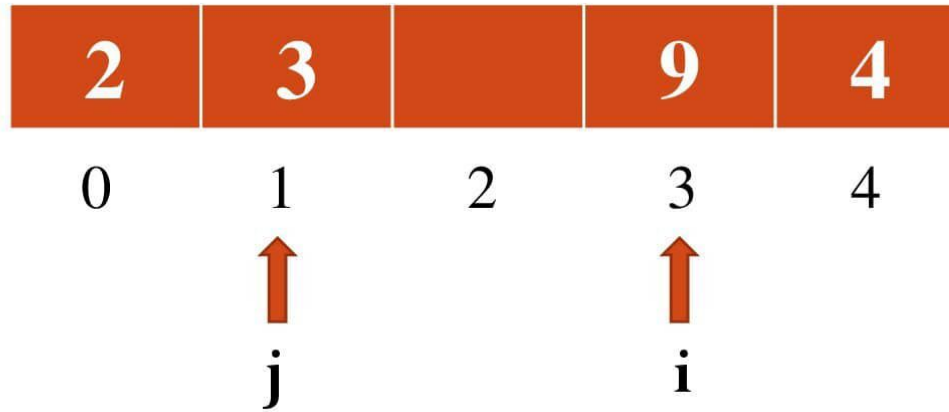
temp= 5

Insertion Sort



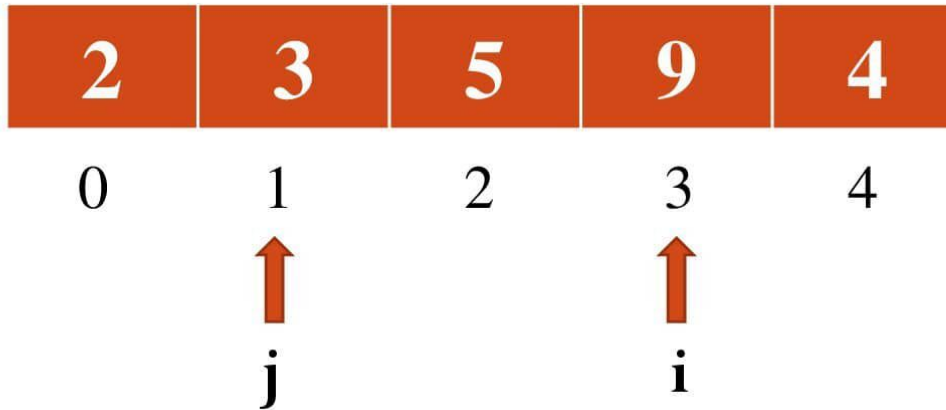
temp= 5

Insertion Sort



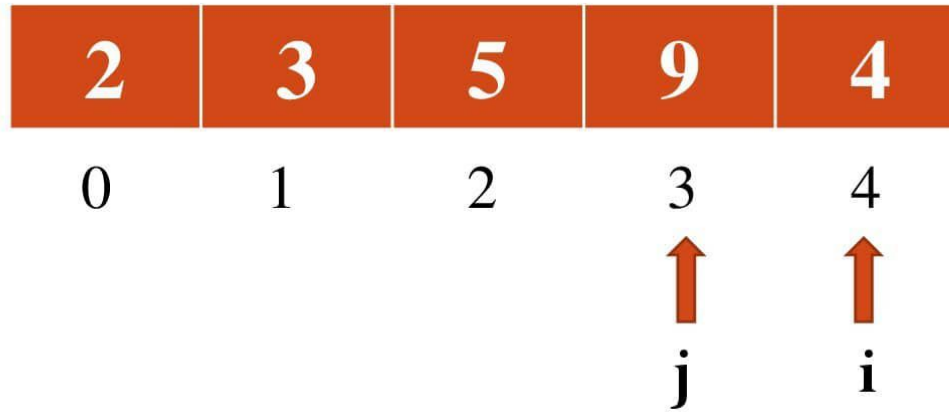
temp= 5

Insertion Sort

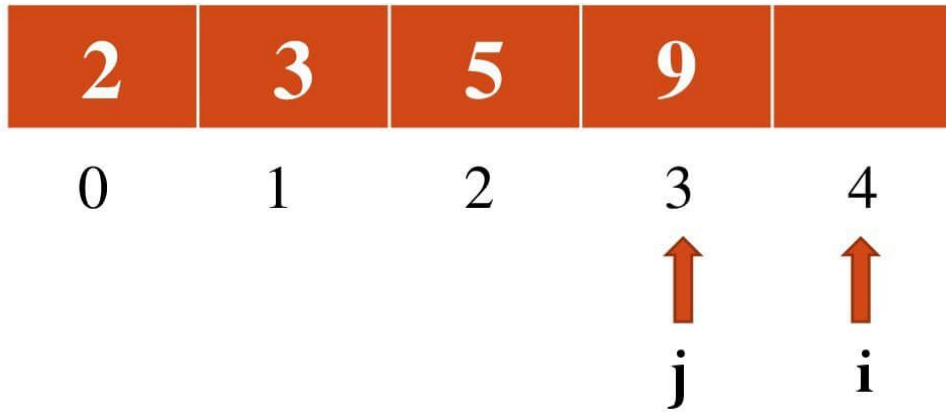


temp= 5

Insertion Sort

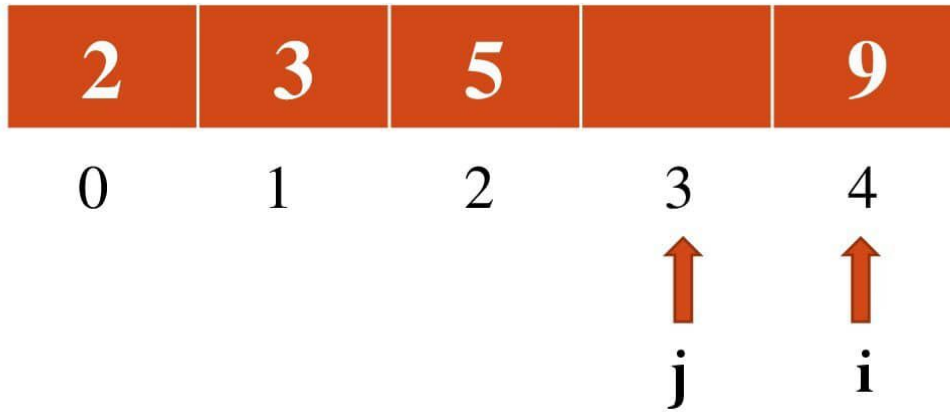


Insertion Sort



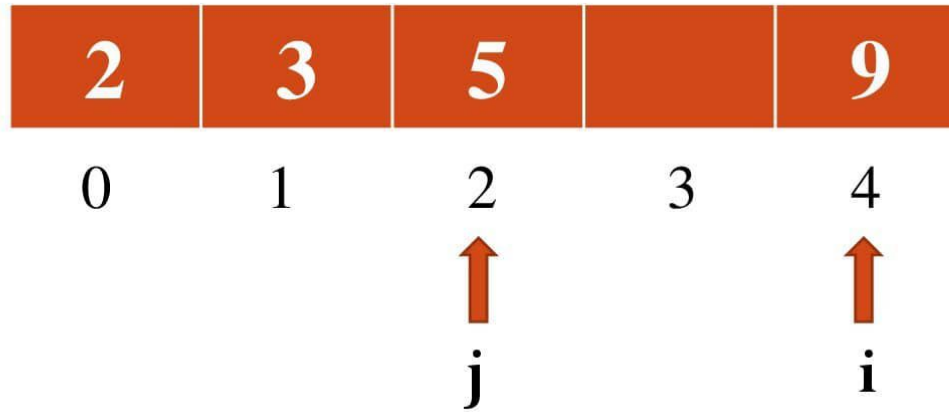
temp= 4

Insertion Sort



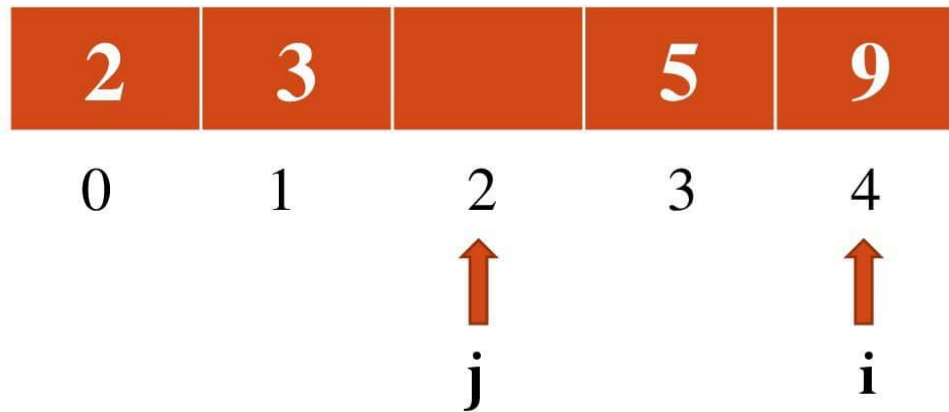
temp= 4

Insertion Sort



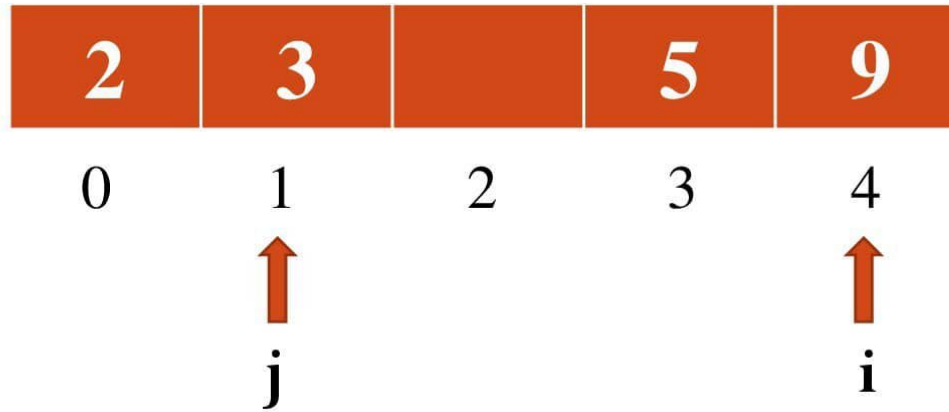
temp= 4

Insertion Sort



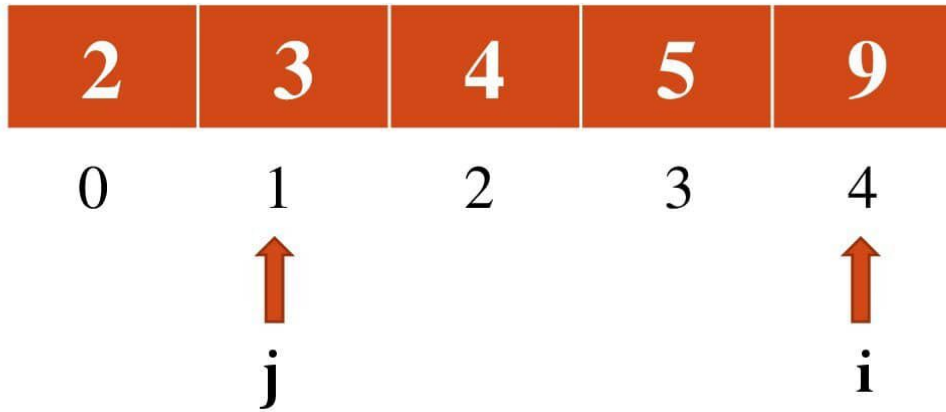
temp= 4

Insertion Sort



temp= 4

Insertion Sort



temp= 4

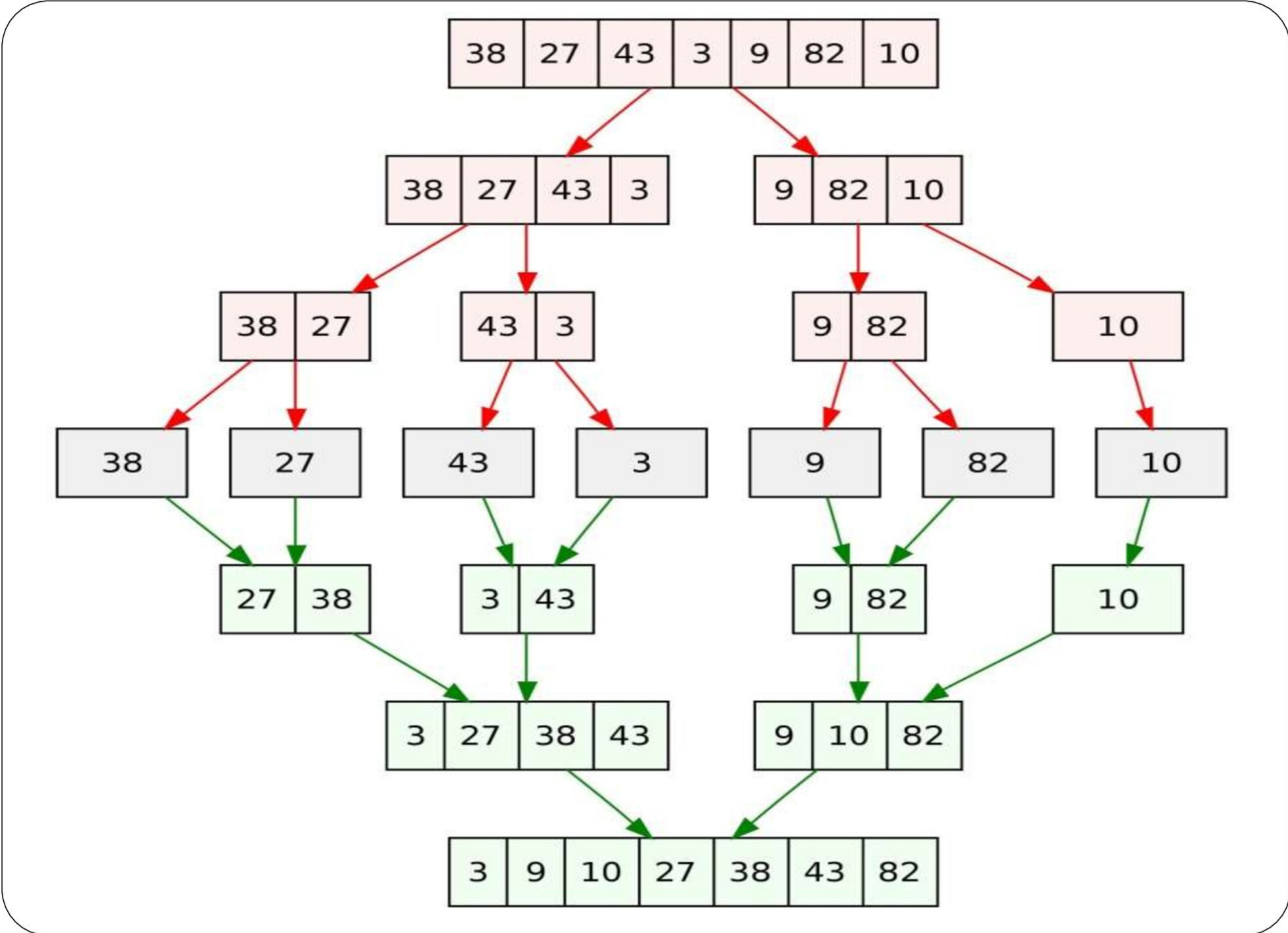
Insertion Sort

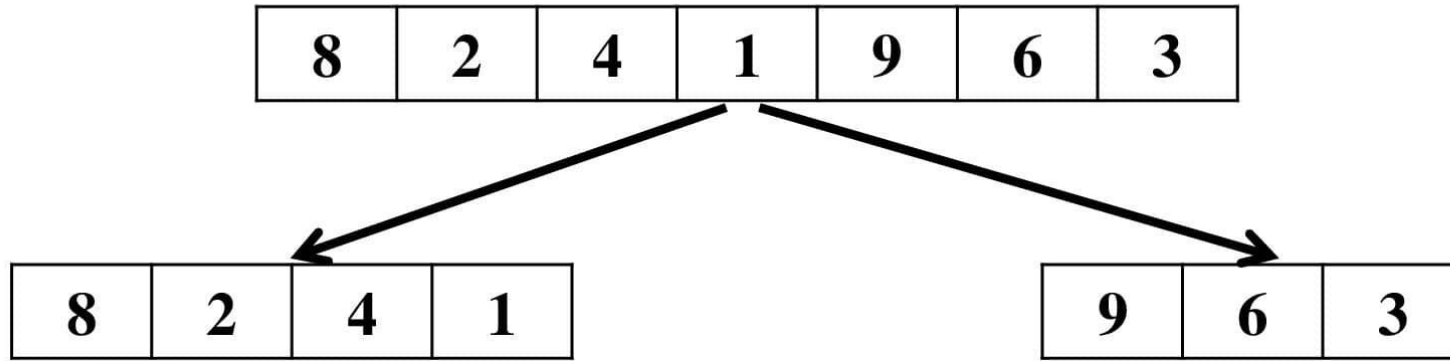
Algorithm InsertionSort(A, n)

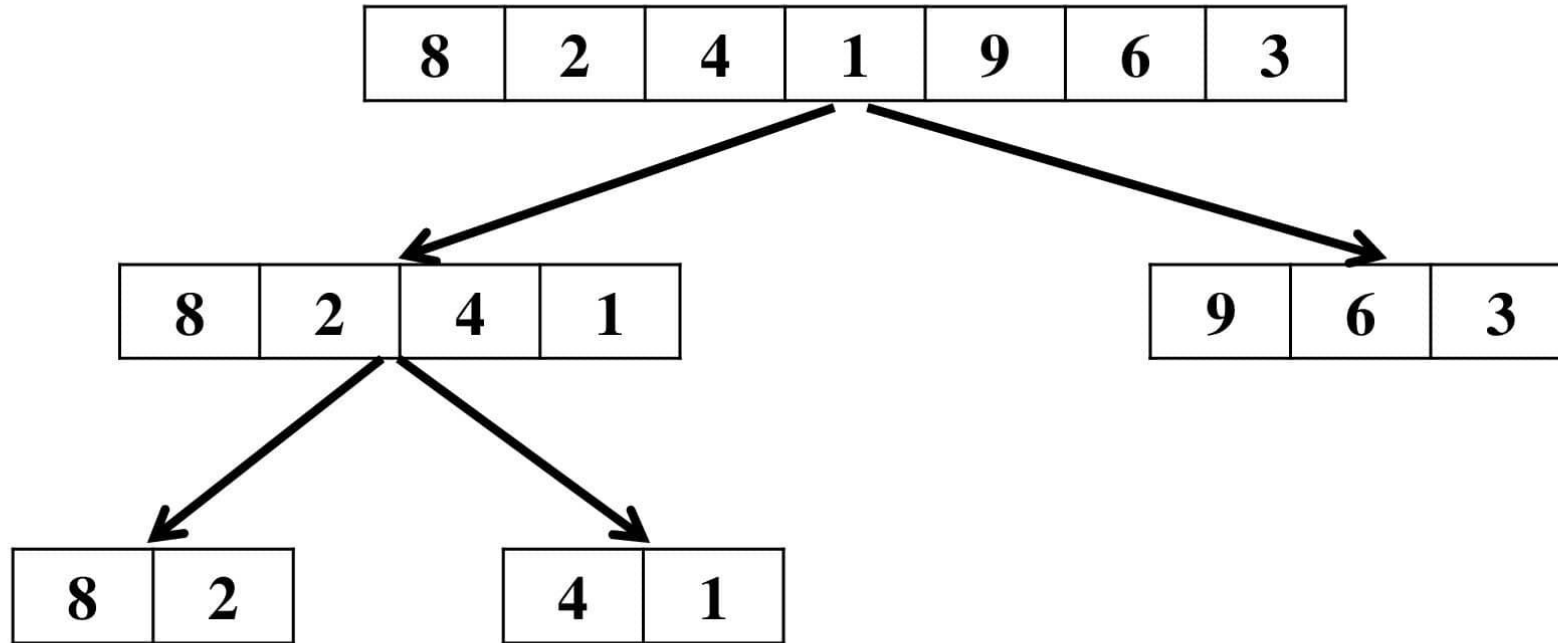
1. for $i=1$ to $n-1$ do
 1. $temp=A[i]$
 2. $j=i-1$
 3. While $j \geq 0$ and $A[j] > temp$ do
 1. $A[j+1] = A[j]$
 2. $j=j-1$
 4. $A[j+1] = temp$

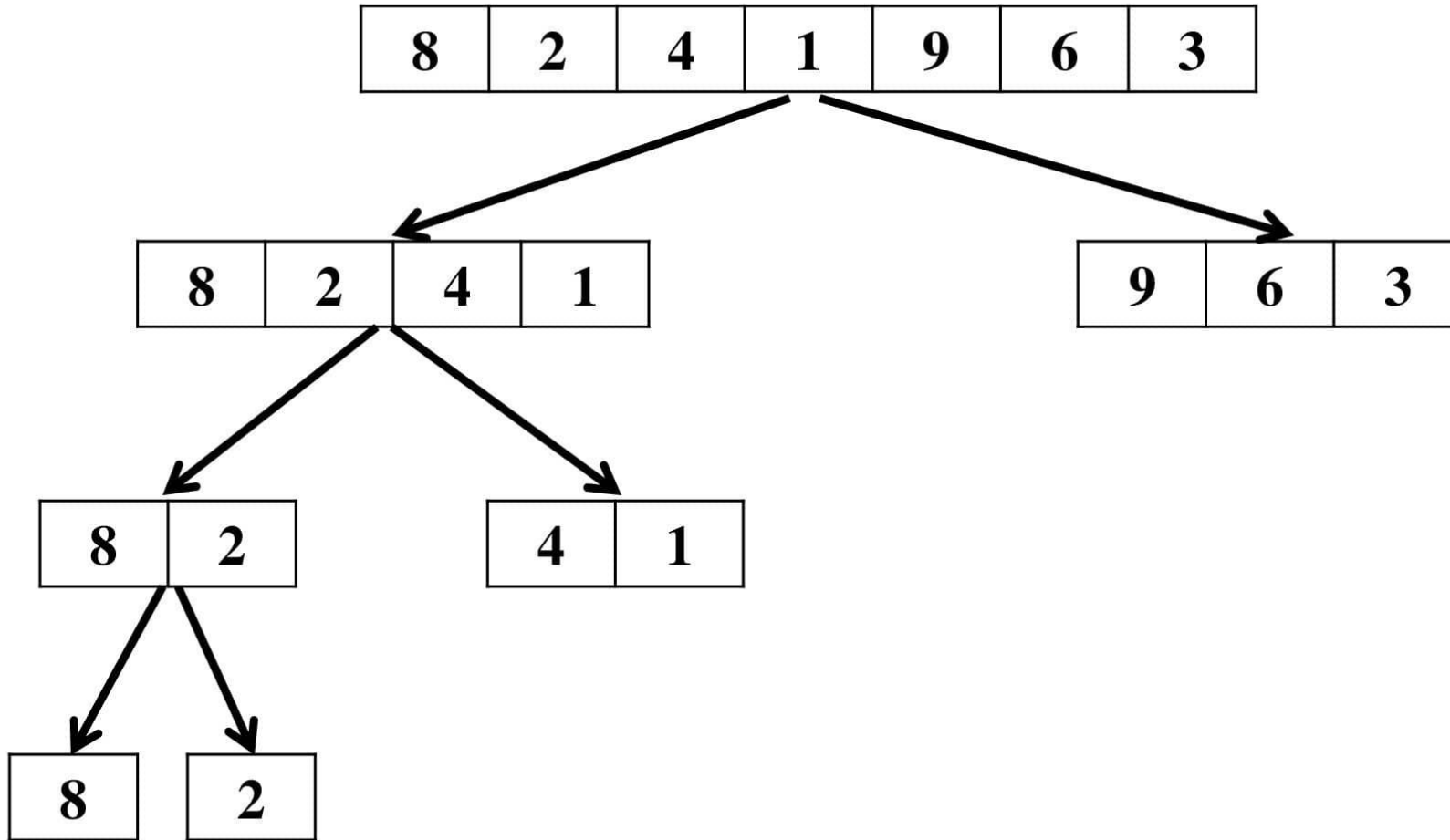
Merge Sort

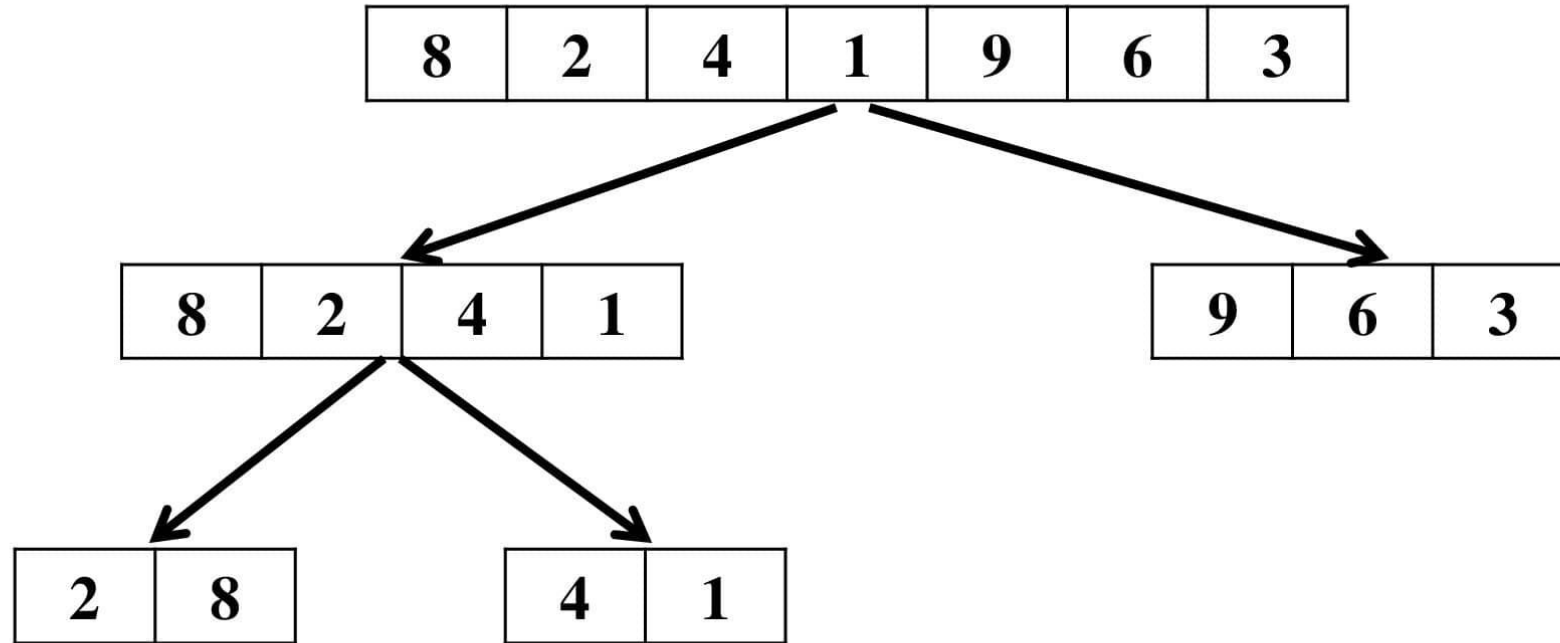
- Given a sequence of n elements $a[1], \dots, a[n]$. Split this array into two sets $a[1], \dots, a[n/2]$ and $a[(n/2)+1], \dots, a[n]$. Each set is individually sorted, and the resulting sorted sequences are merged to produce a single sorted sequence of n elements.

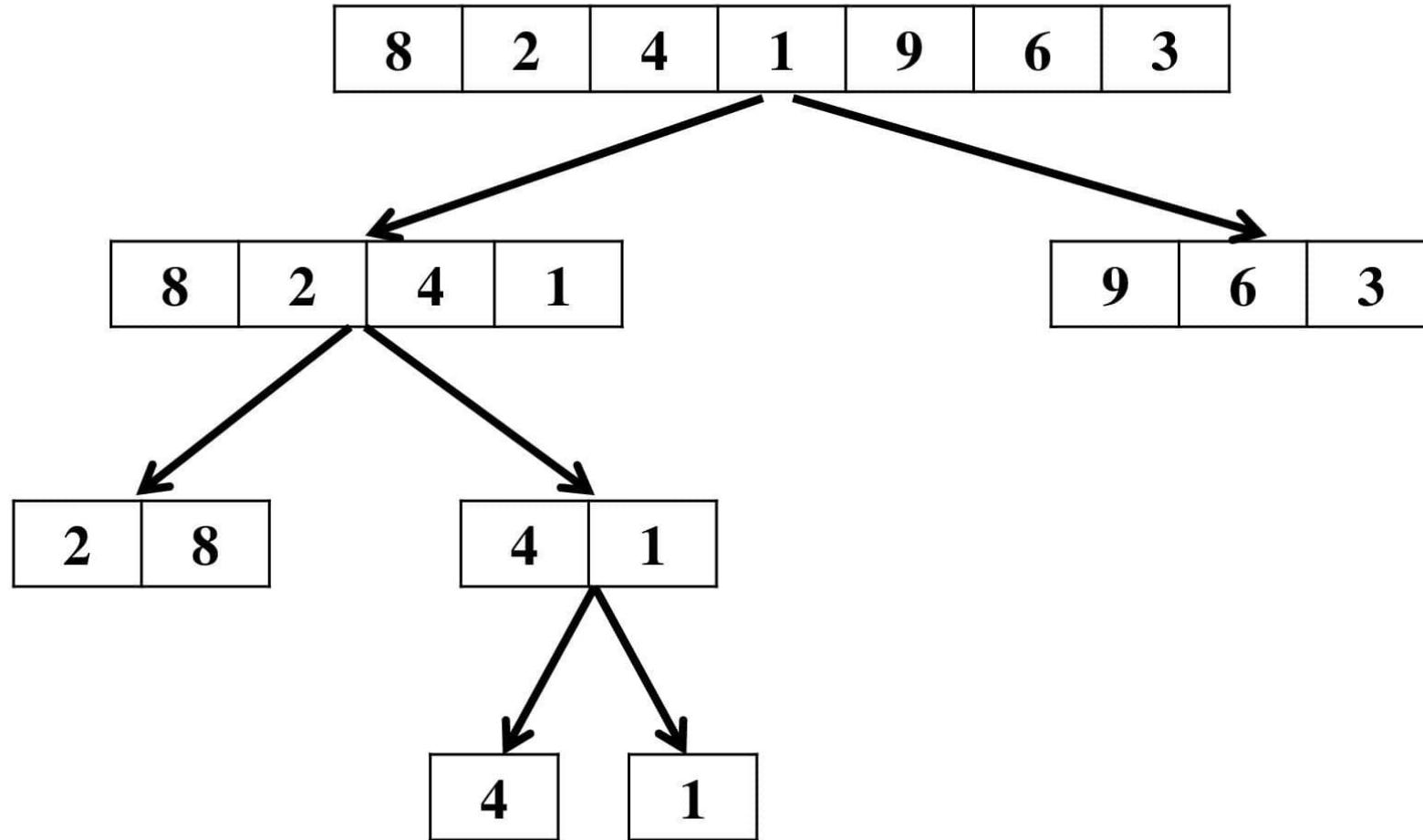


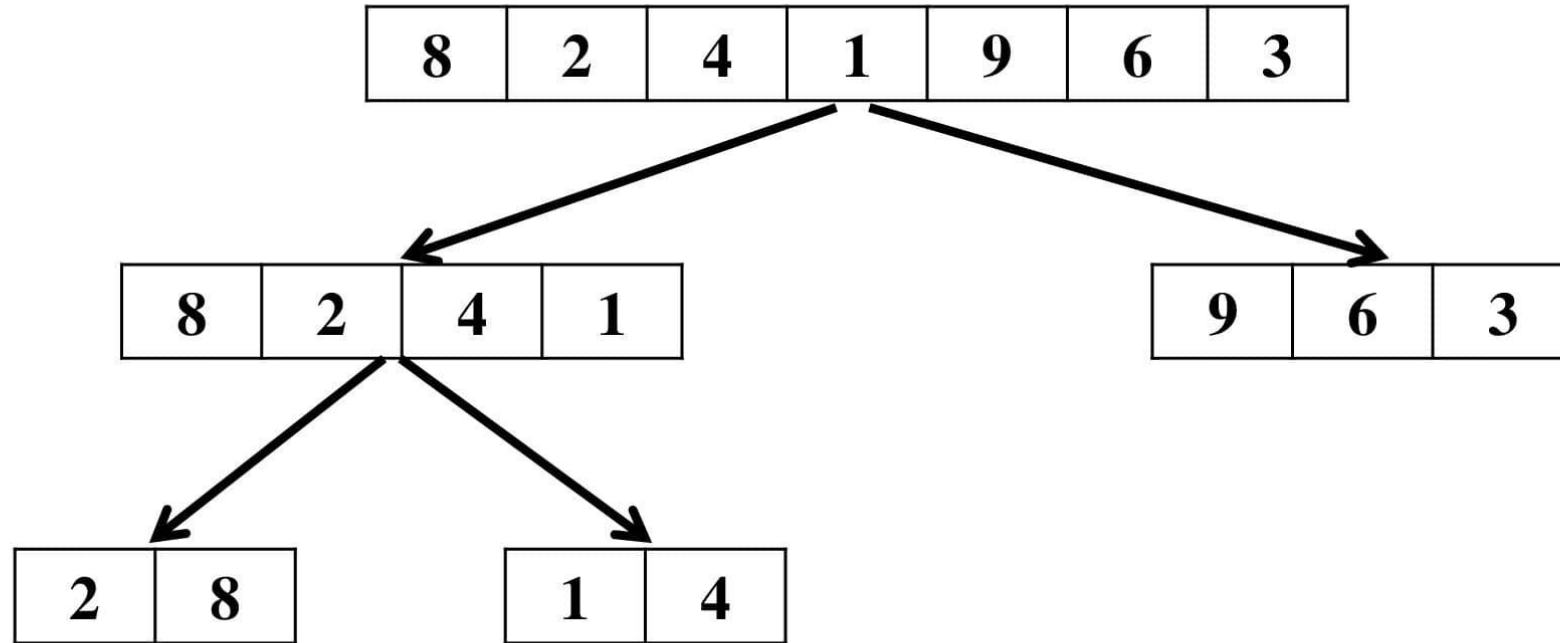


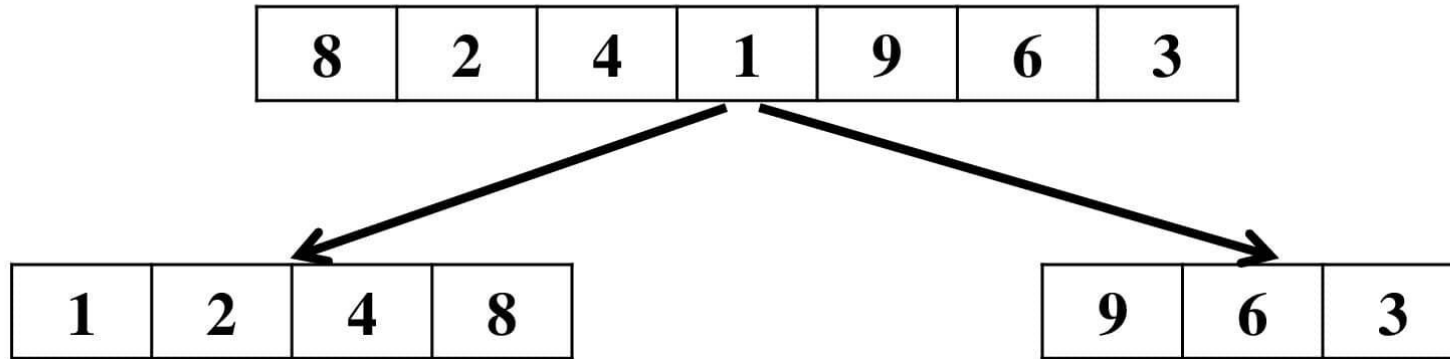


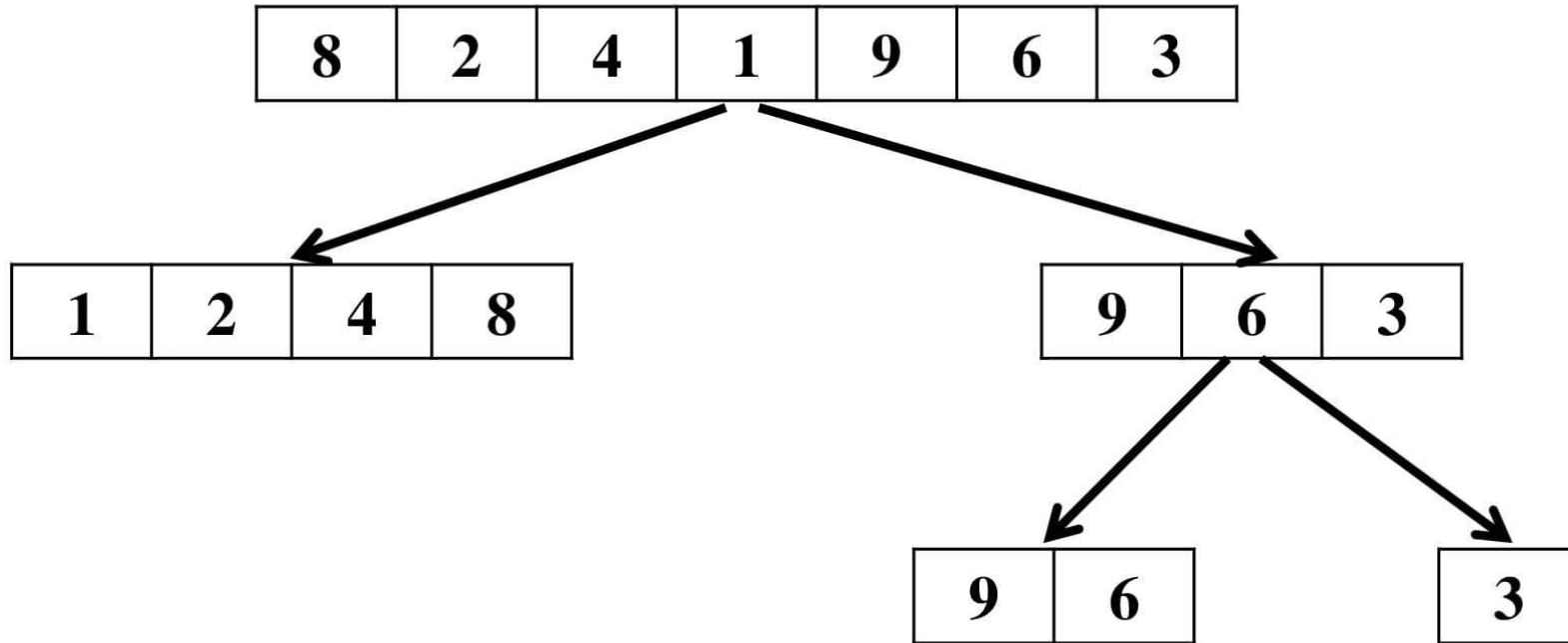


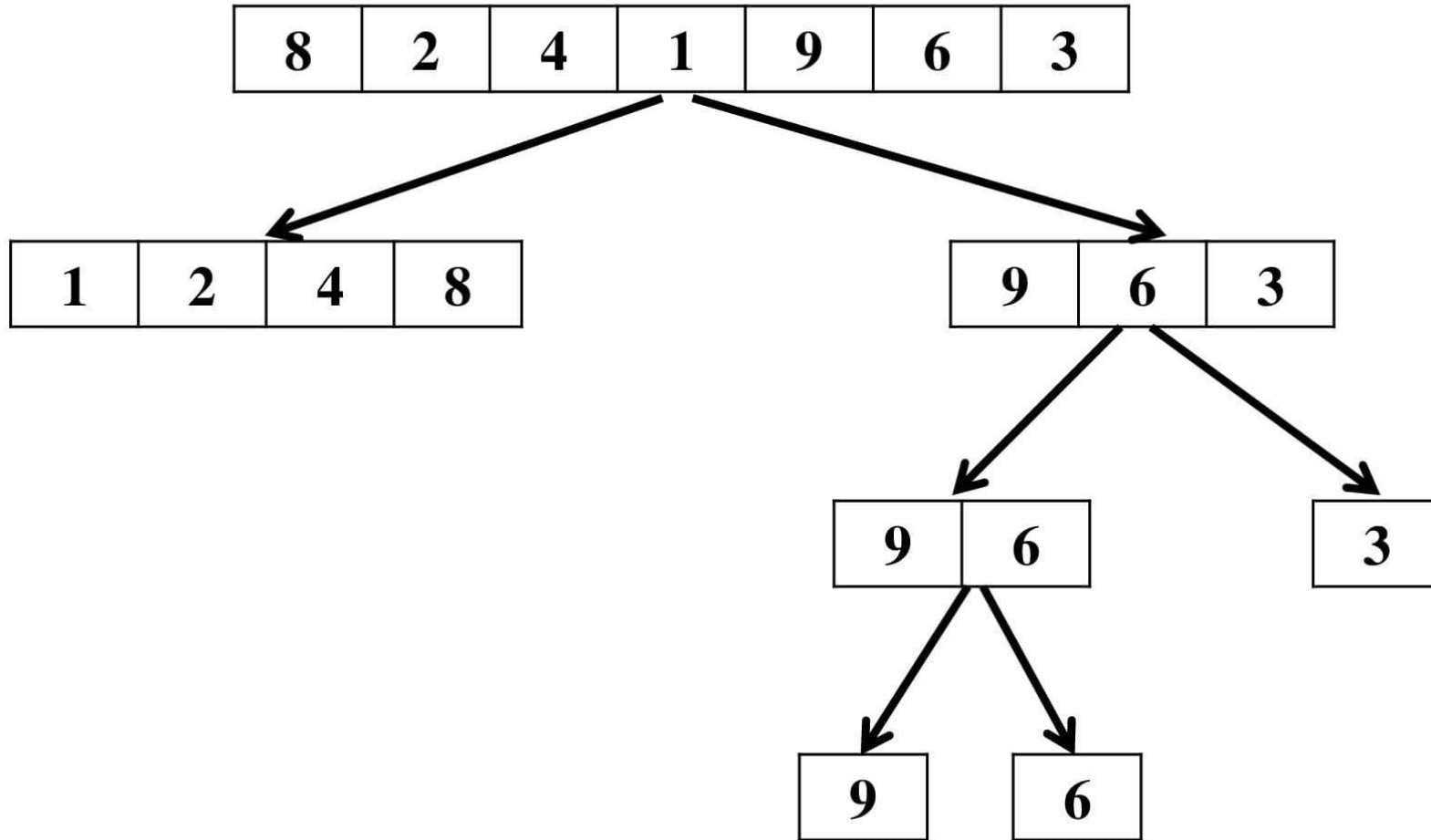


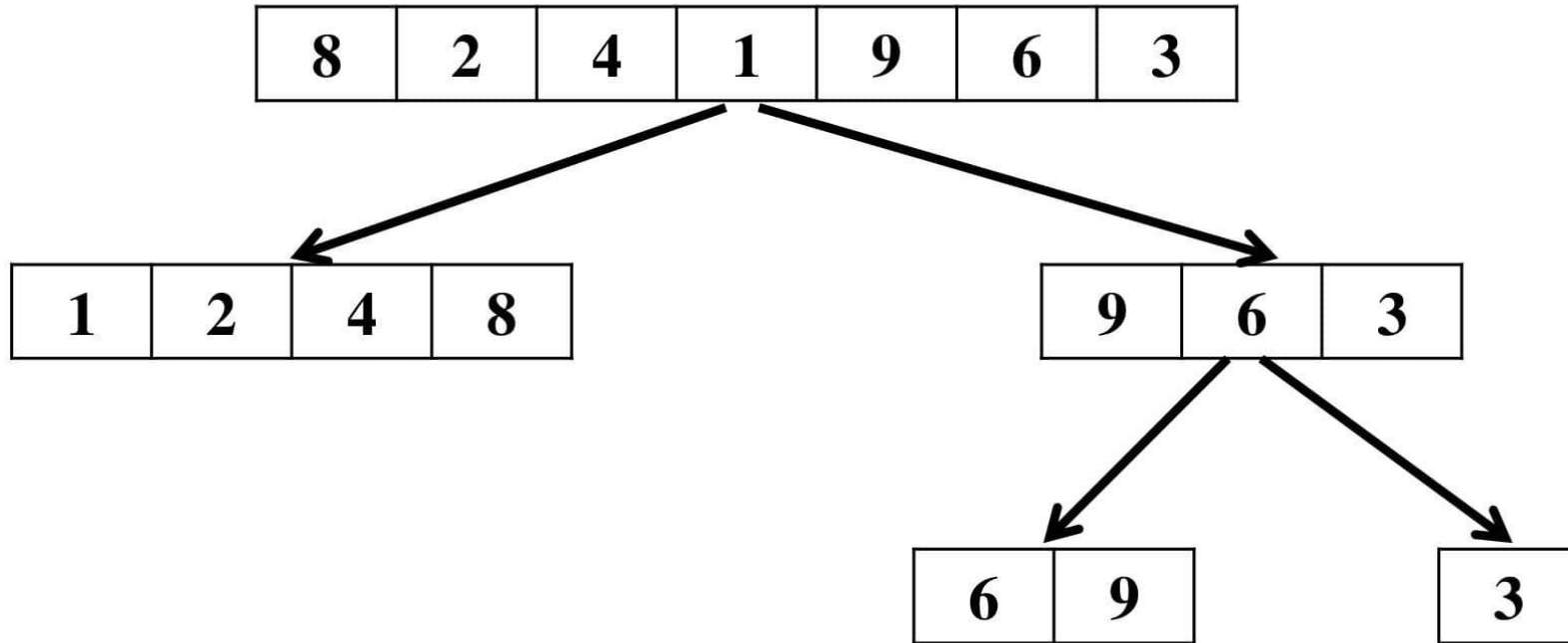


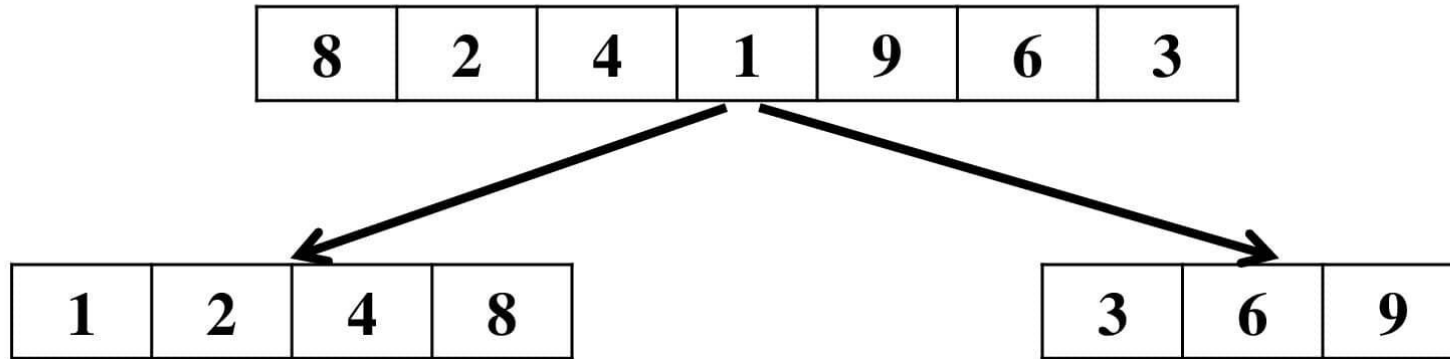












1	2	3	4	6	8	9
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Merge Sort

Algorithm MergeSort(low, high)

{

 mid = (low + high) / 2;

 MergeSort(low, mid);

 MergeSort(mid+1, high);

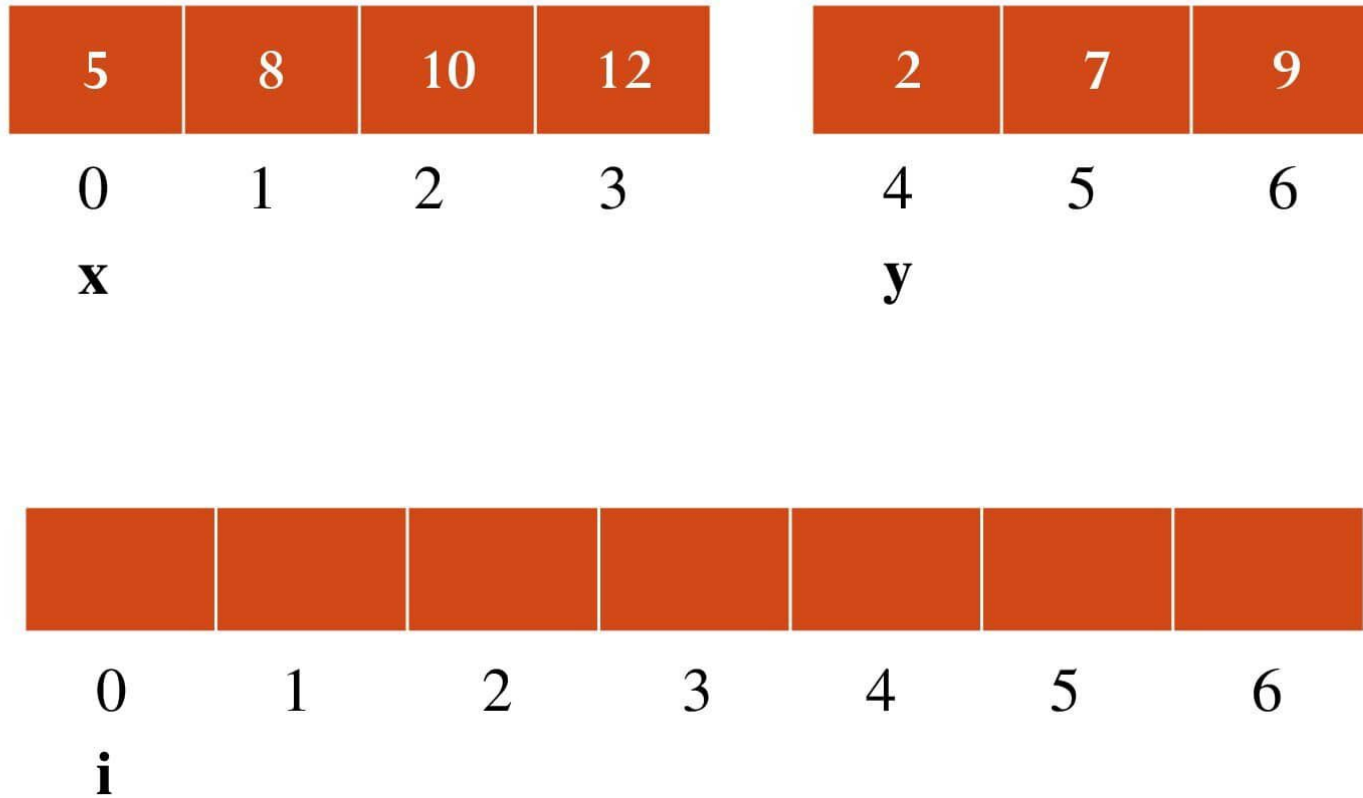
 Merge(low, mid, high);

}

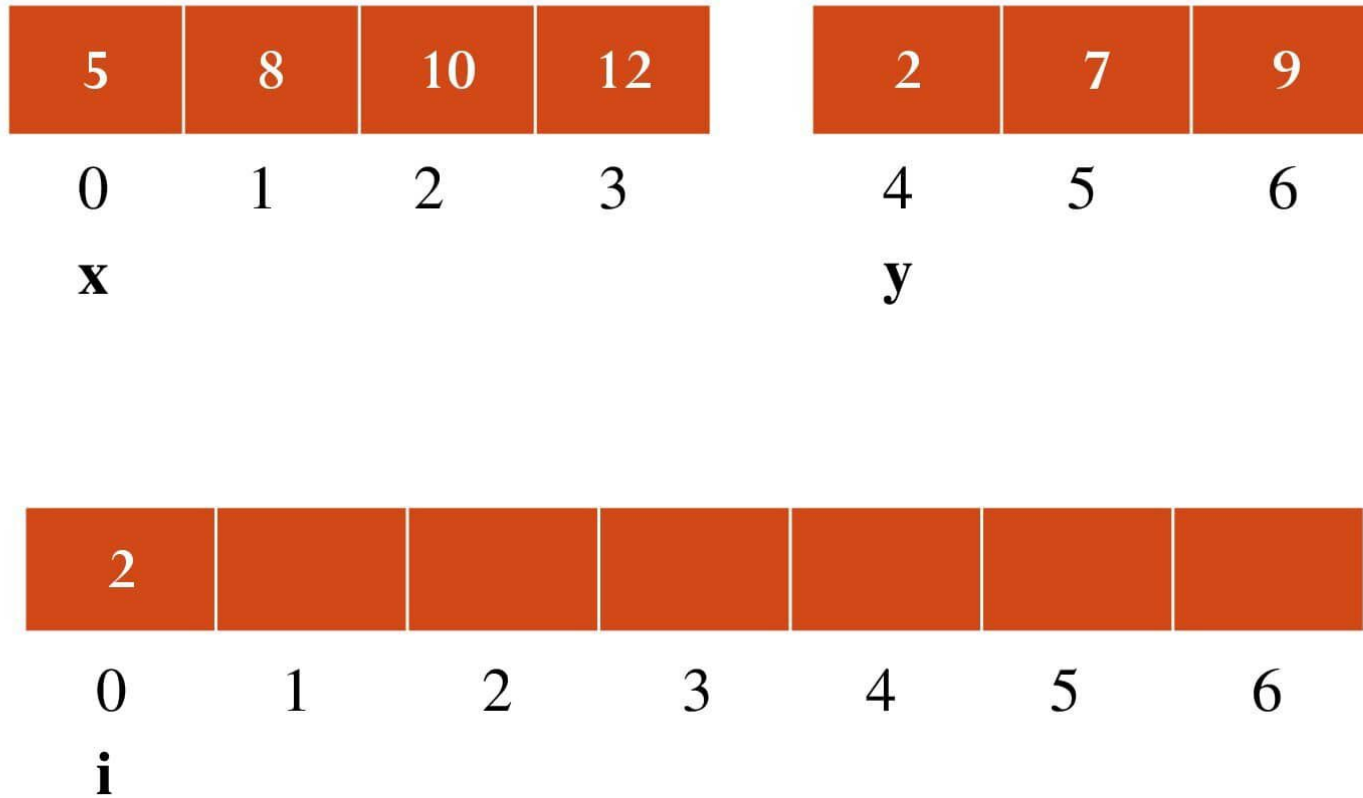
Merge



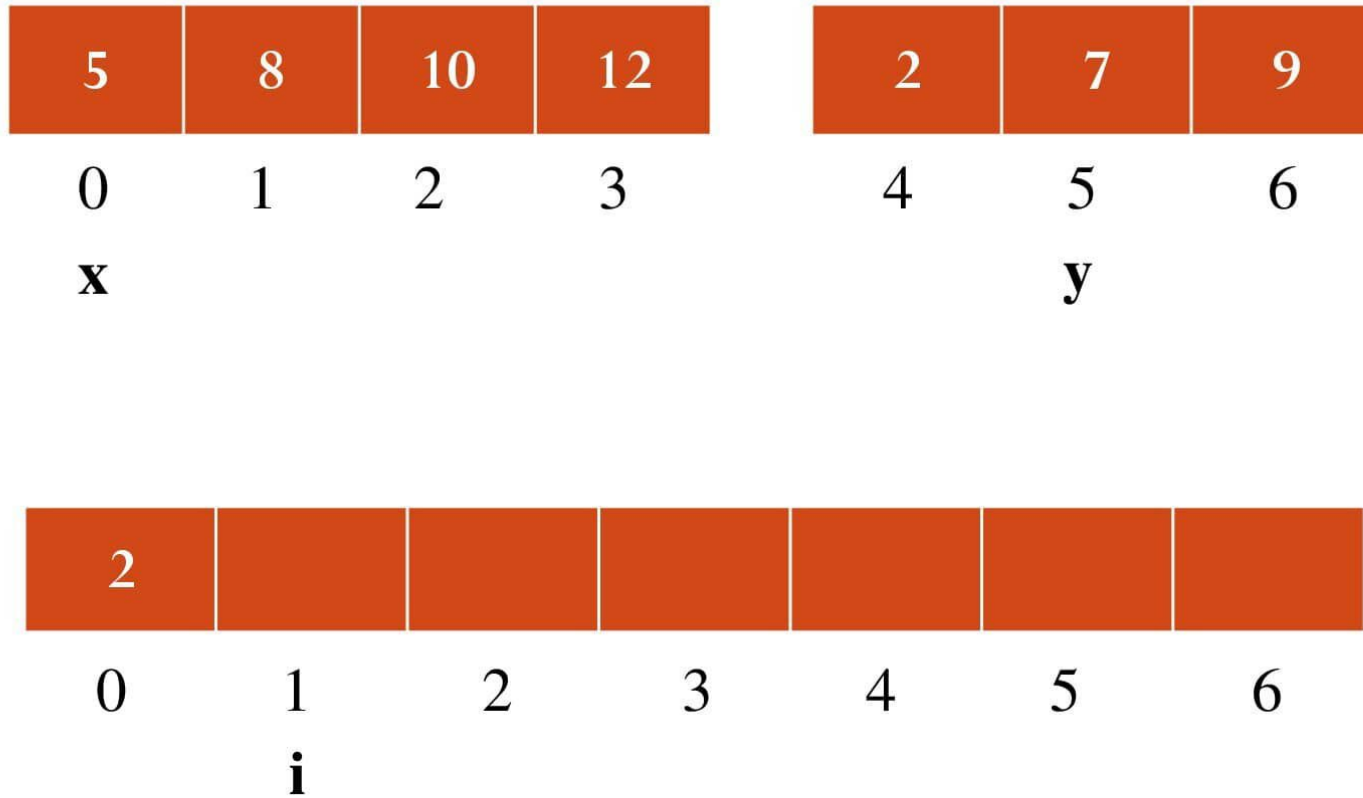
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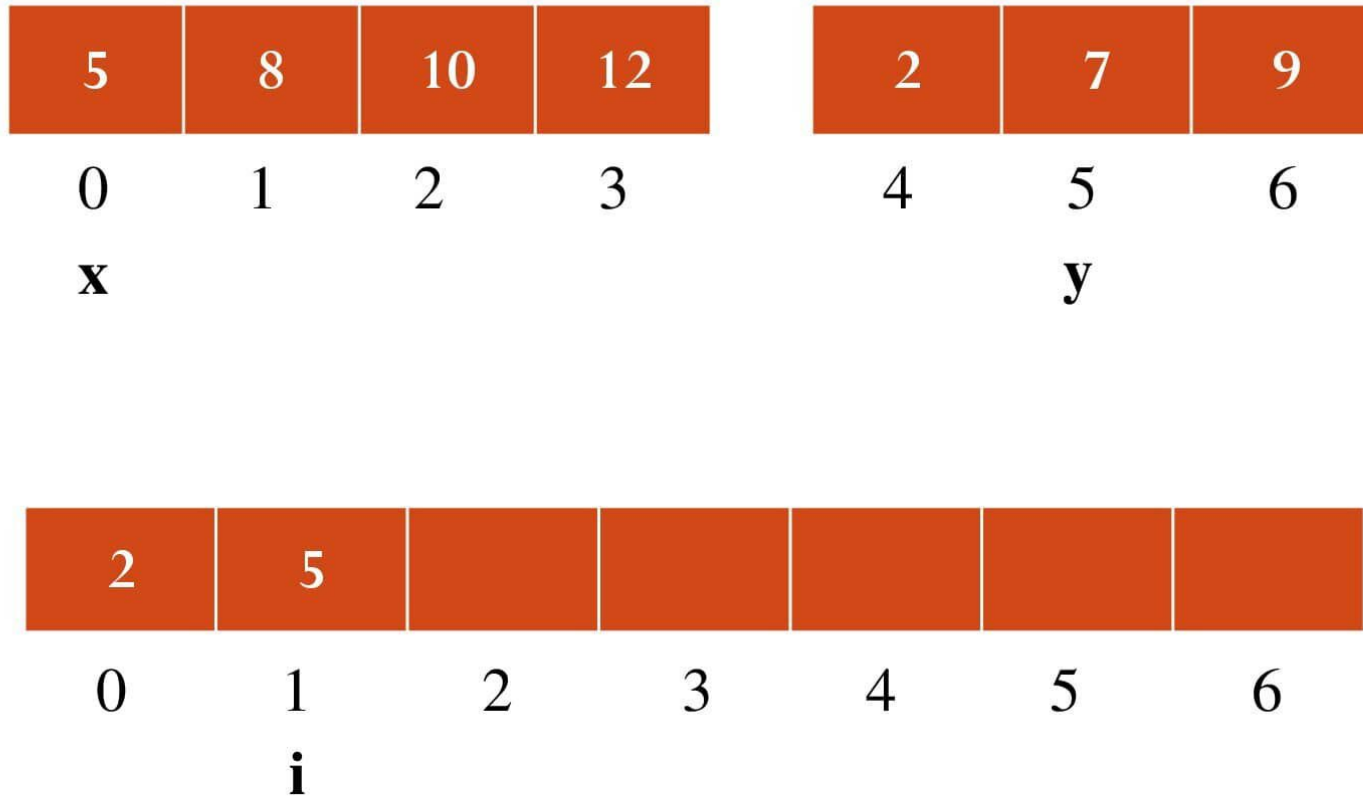
Merge



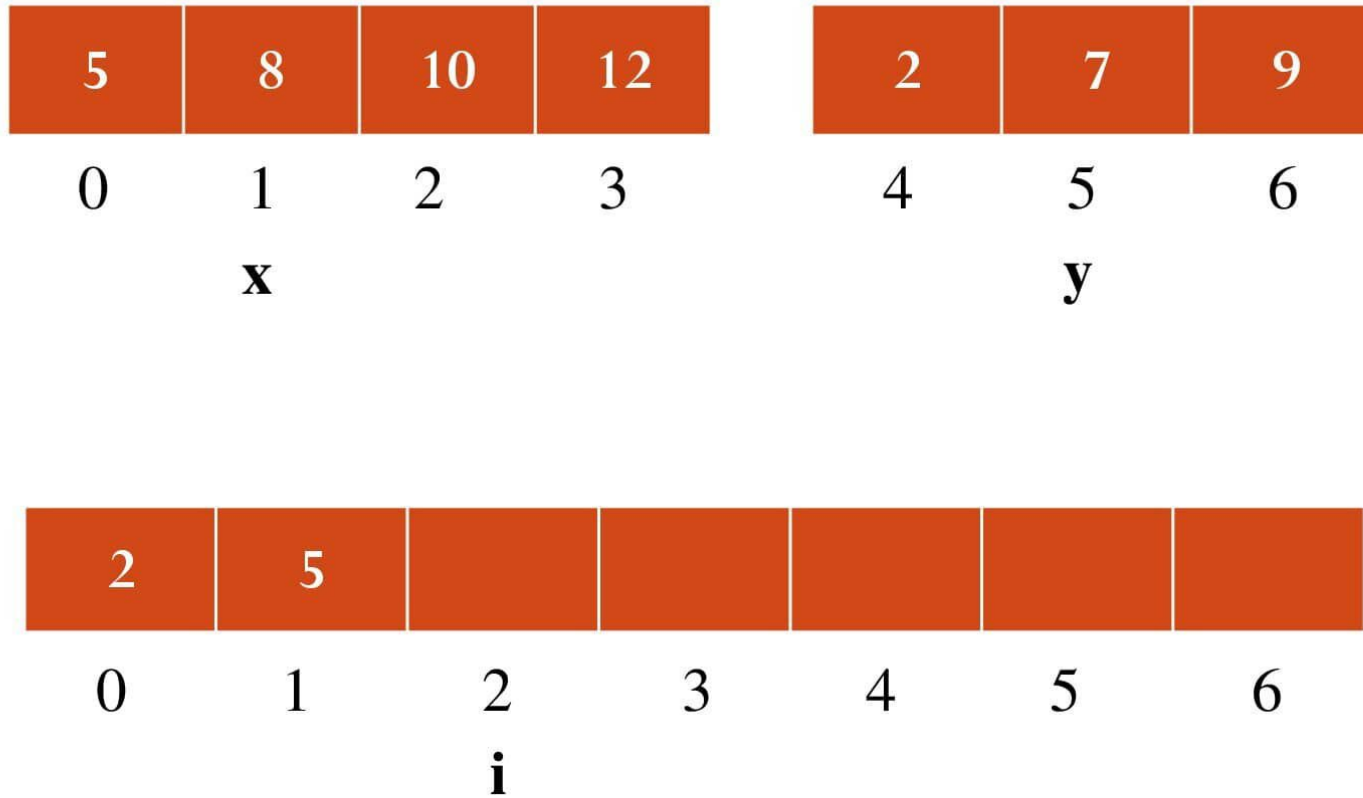
Merge



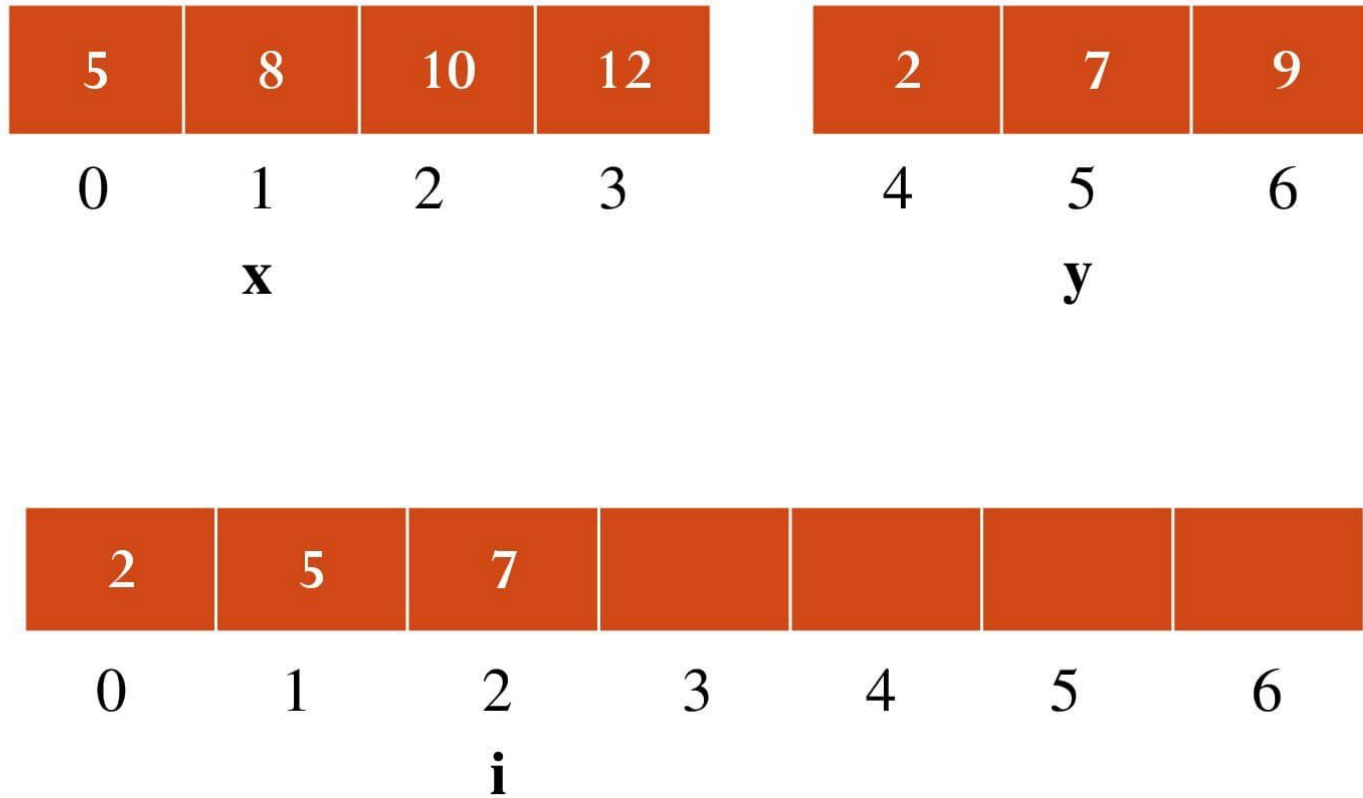
Merge



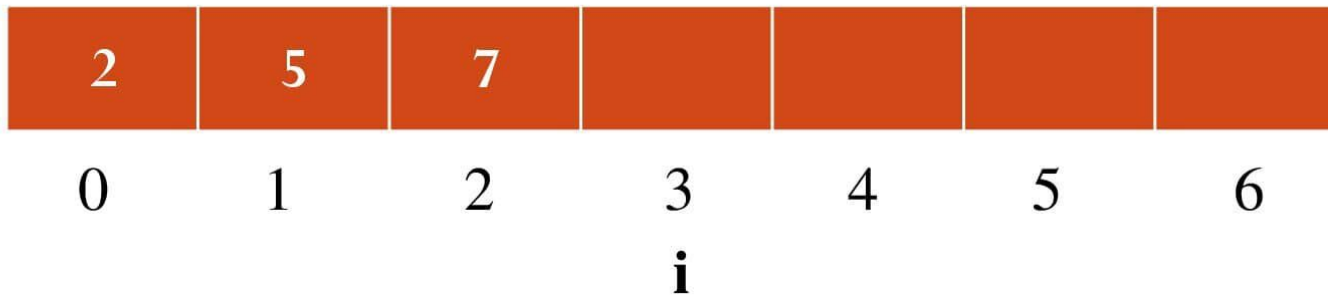
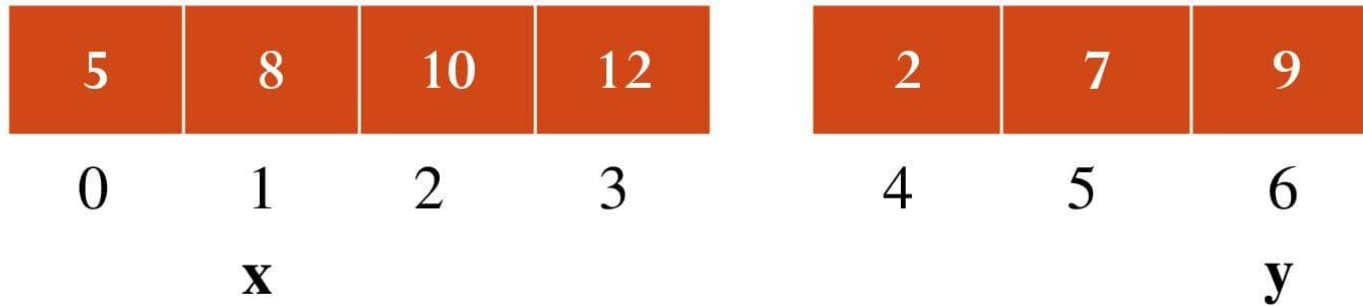
Merge



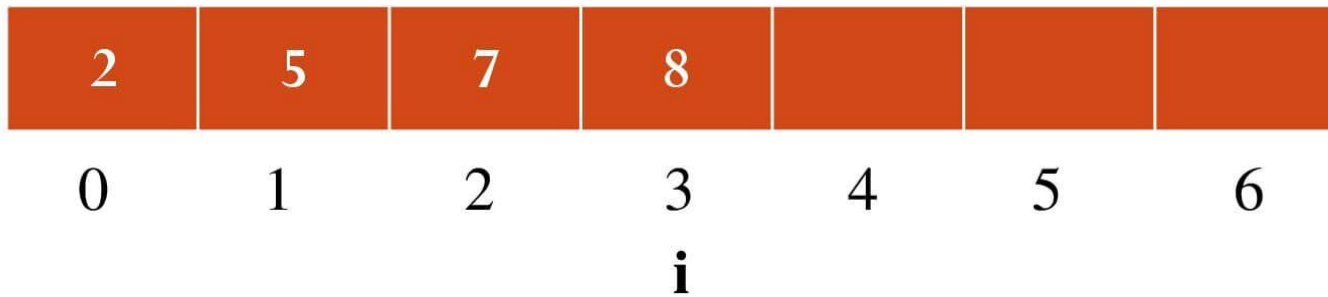
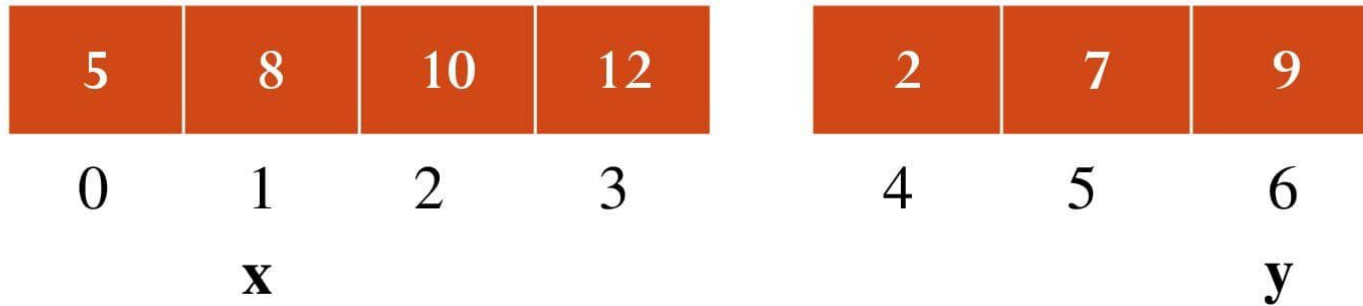
Merge



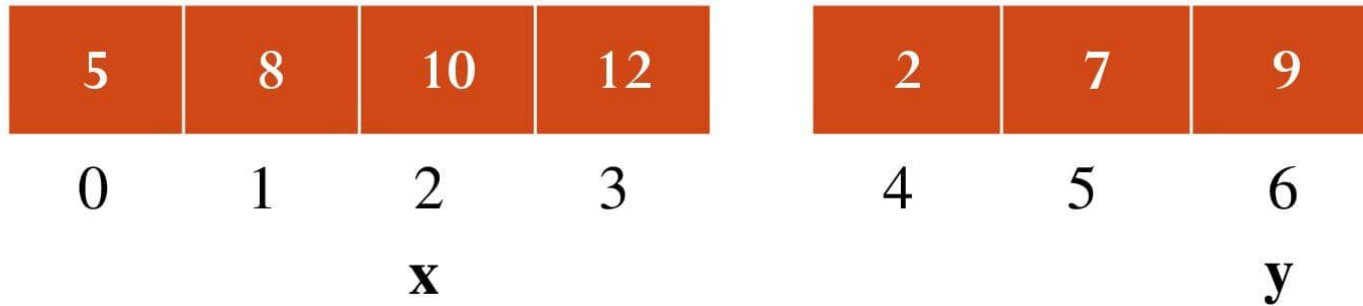
Merge



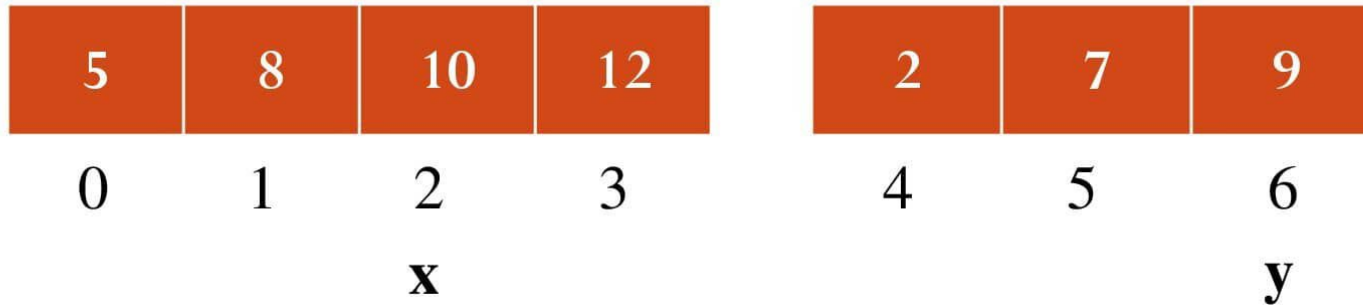
Merge



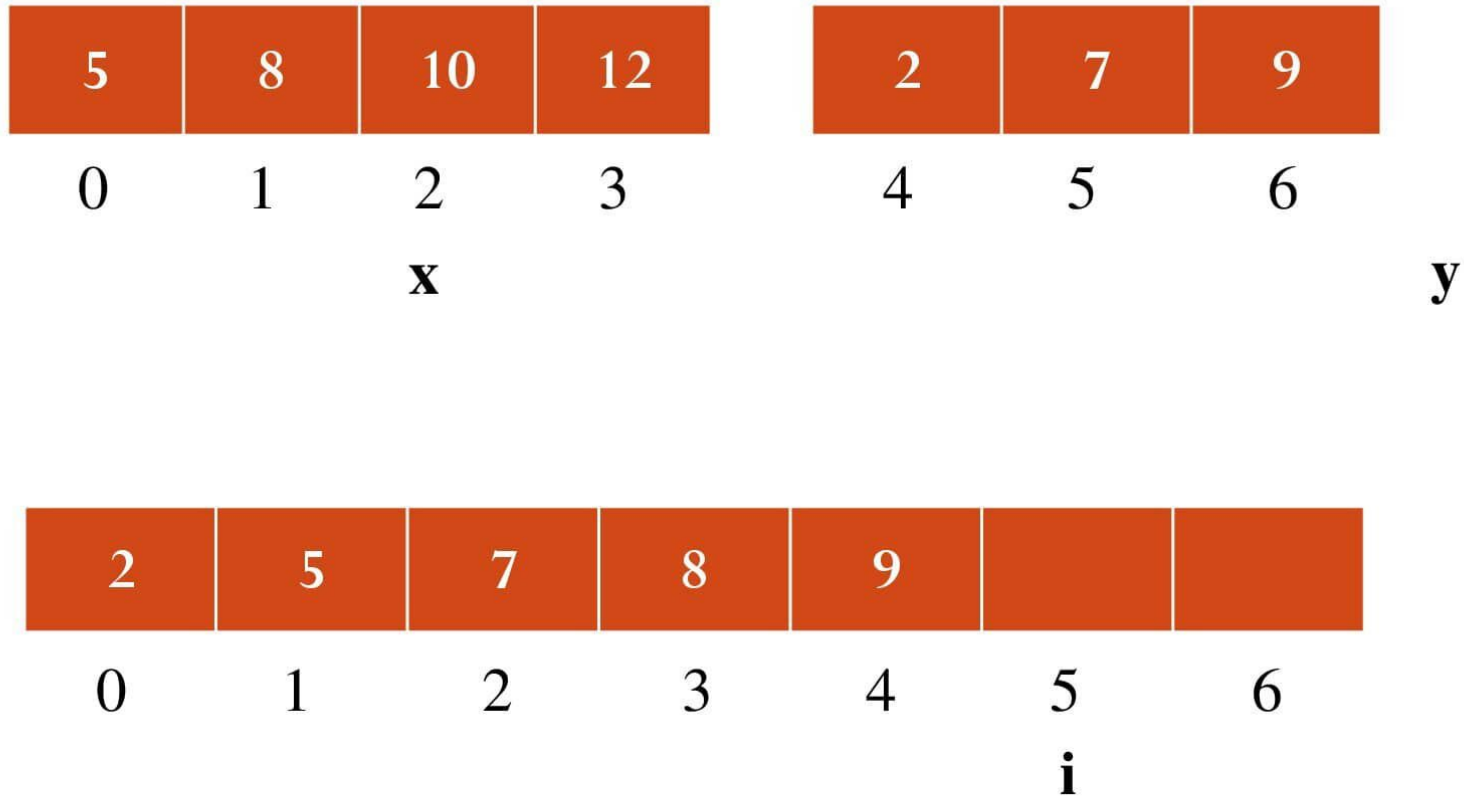
Merge



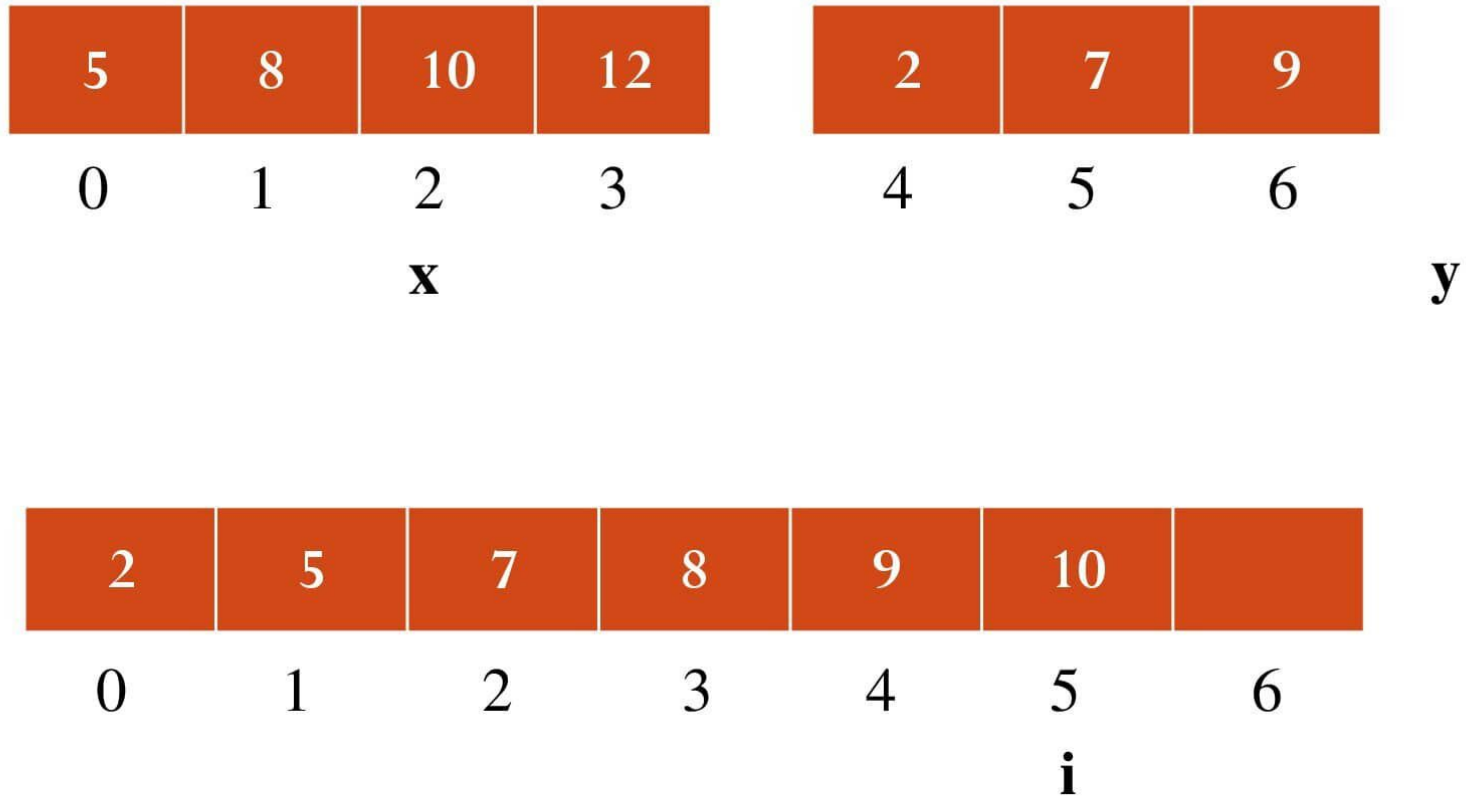
Merge



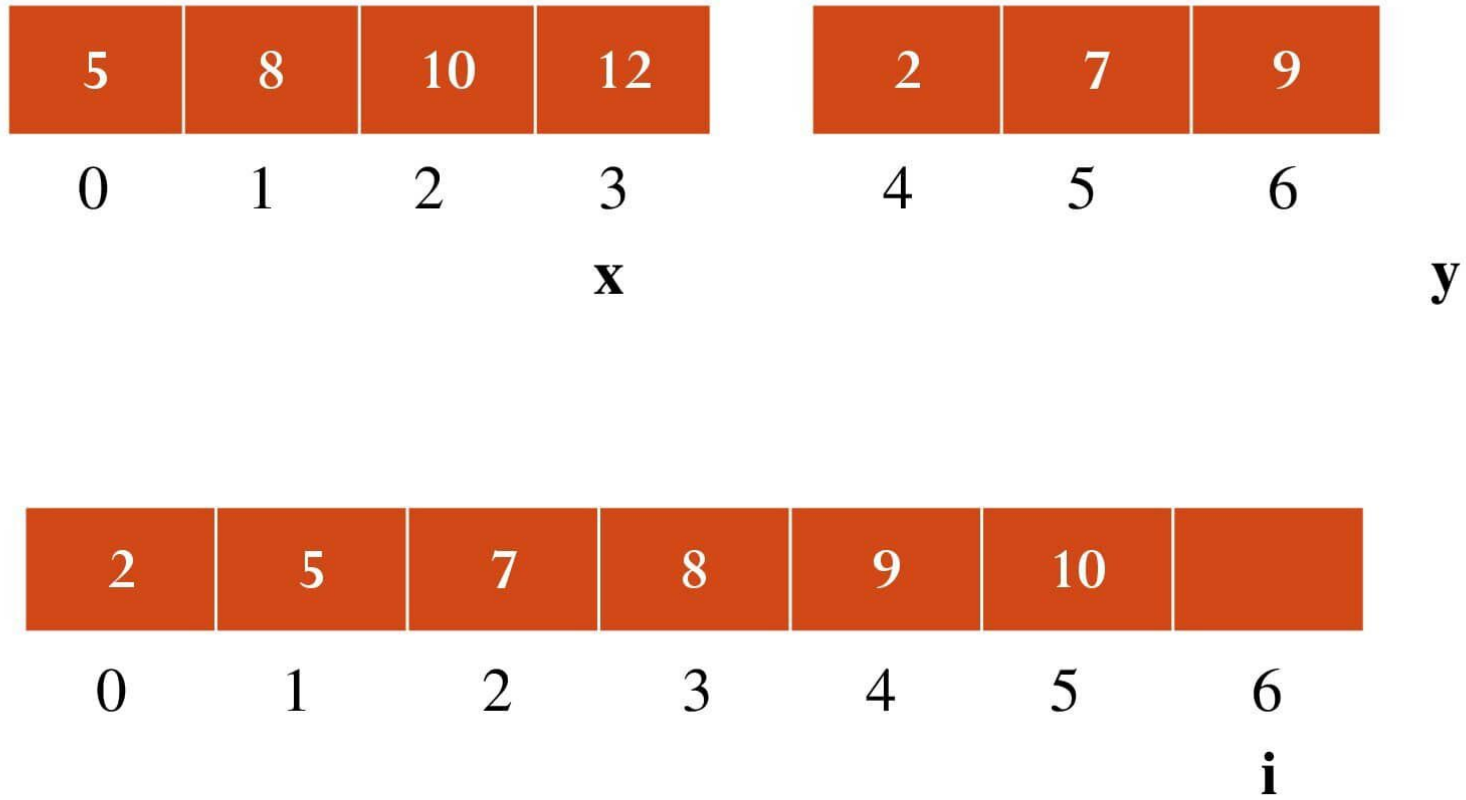
Merge



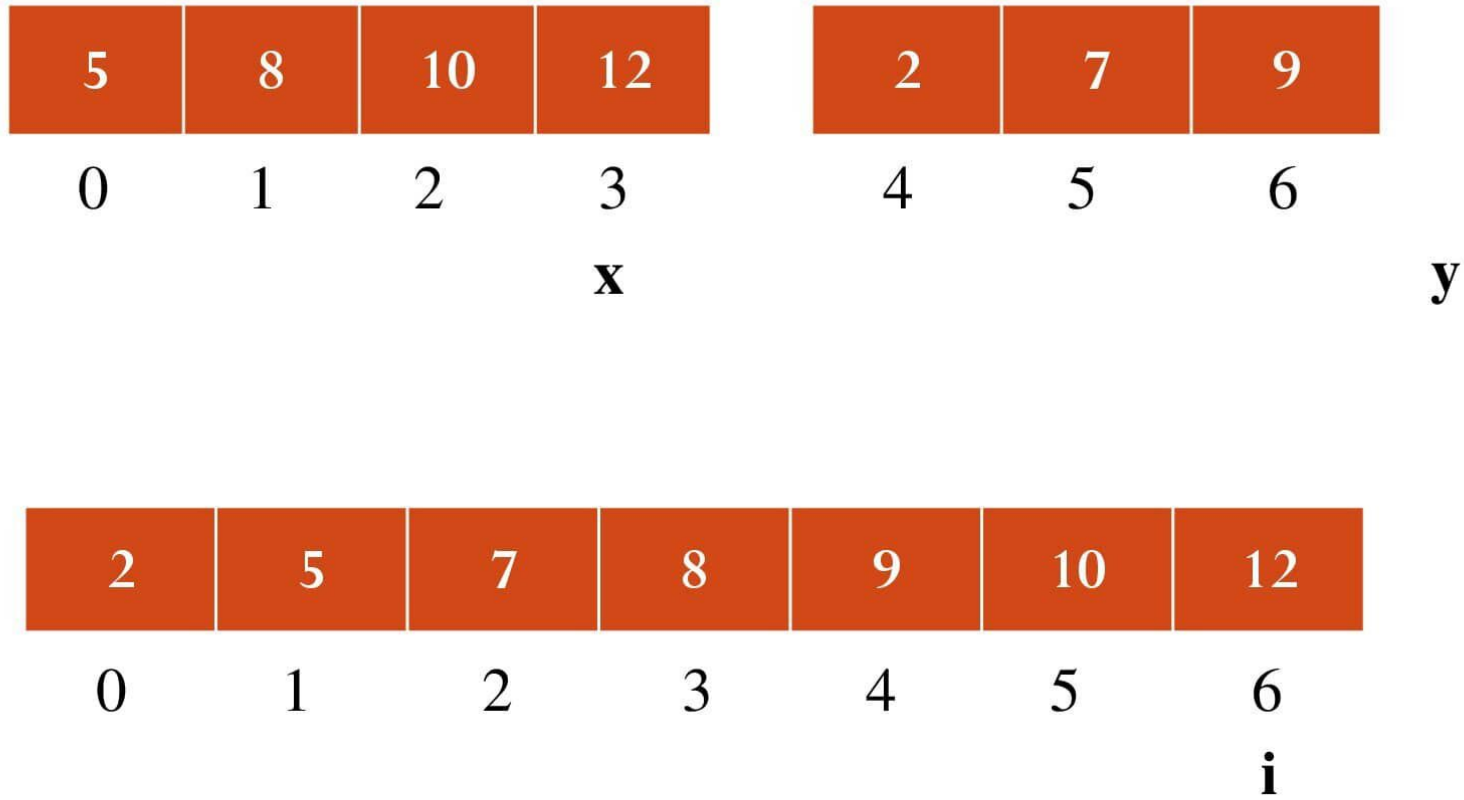
Merge



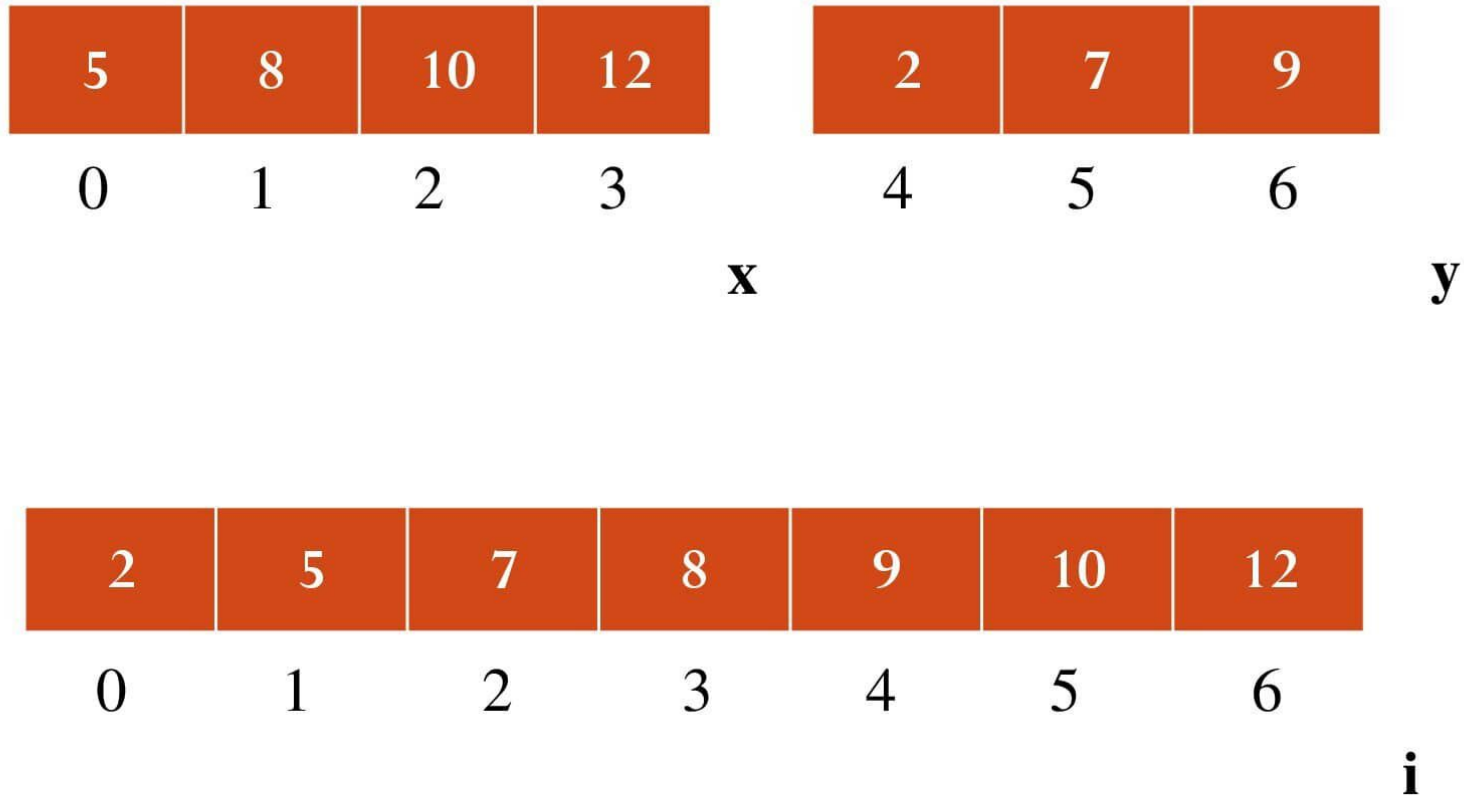
Merge



Merge



Merge



Algorithm Merge(low, mid, high)

```
{   i= low; x= low; y= mid + 1;
  while(( $x \leq mid$ ) and ( $y \leq high$ )) do
  {   if (  $a[x] \leq a[y]$  ) then
      {       b[i] = a[x];
              x = x+1;
            }
      else
      {       b[i] = a[y];
              y = y+1;
            }
        i=i+1;
  }
}
```

```
  if(  $x \leq mid$ ) then
  {   for  $k=x$  to  $mid$  do
      {       b[i] = a[k];
              i =i+1;
            }
  }
  else
  {   for  $k=y$  to  $high$  do
      {       b[i] = a[k];
              i =i+1;
            }
  }
  for  $k= low$  to  $high$  do
    a[k] = b[k];
}
```

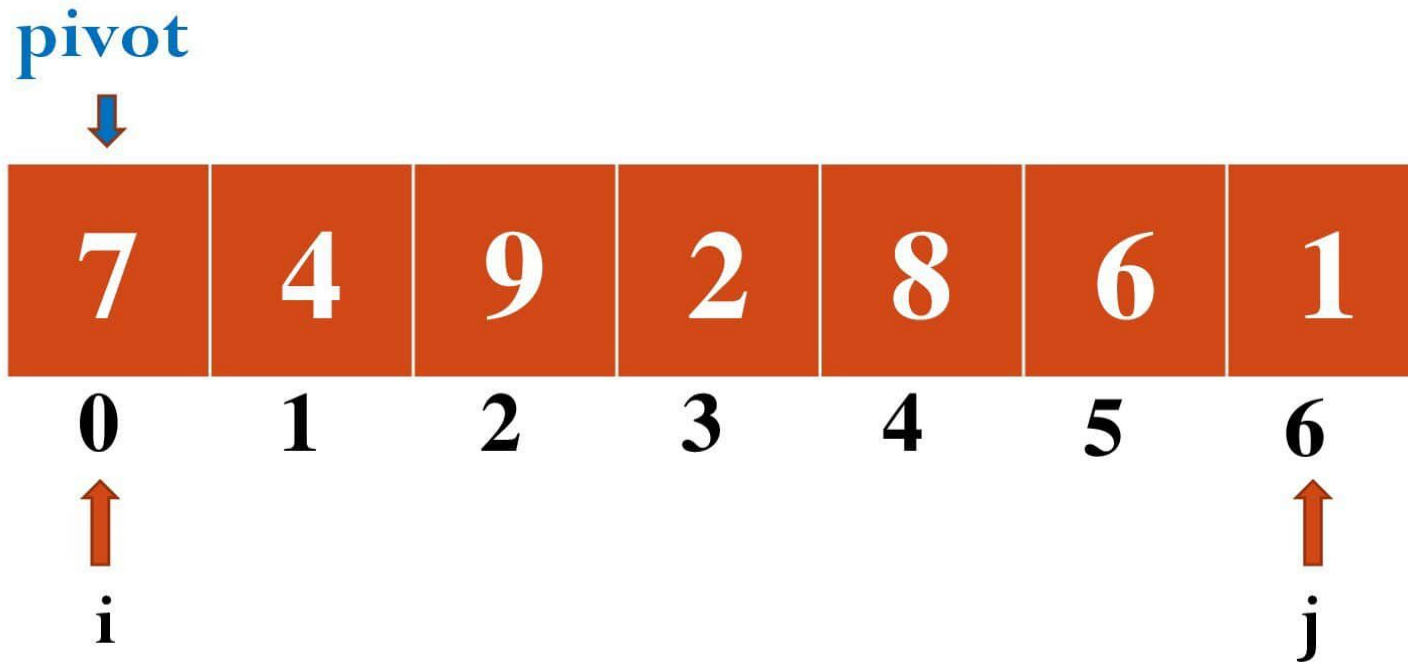

Quick Sort

1. Let **A** be an array with n elements.
2. Select an element 'x' from any position, called a **pivot**, from the list.
3. Partition **A** into two, so that x is placed in position j such that
 1. Elements from position 0 to $j-1$ is less than x
 2. Elements from position $j+1$ to n is greater than xThis is called a partition operation
After this partitioning, the pivot is in its final position
4. Recursively sort the sub-list of lesser elements and the sub-list of greater elements

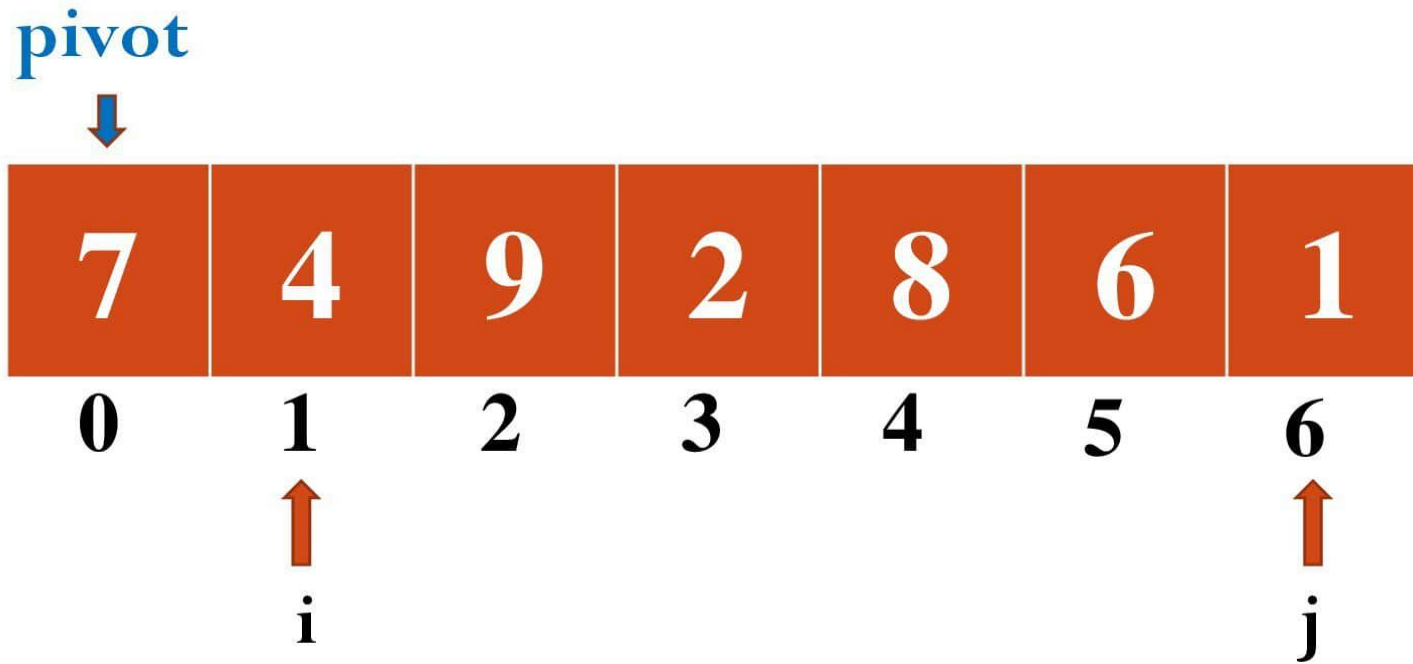
Quick Sort

7	4	9	2	8	6	1
0	1	2	3	4	5	6

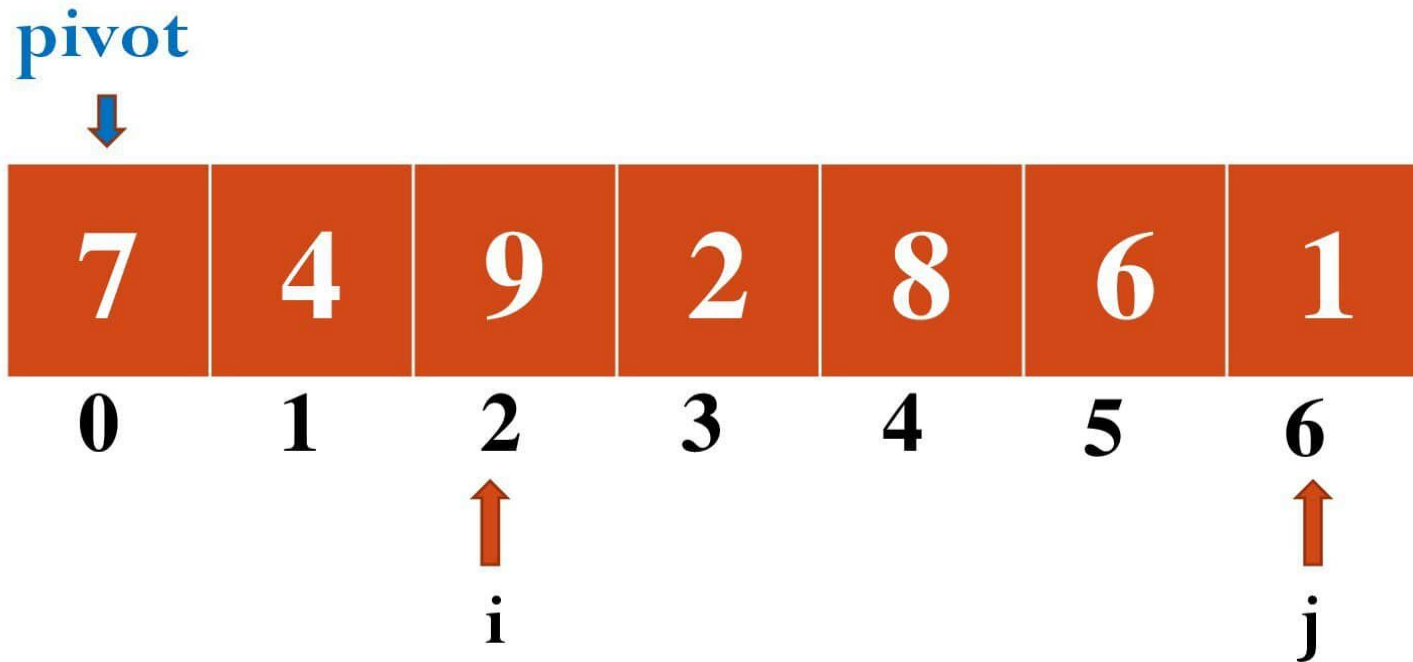
Quick Sort



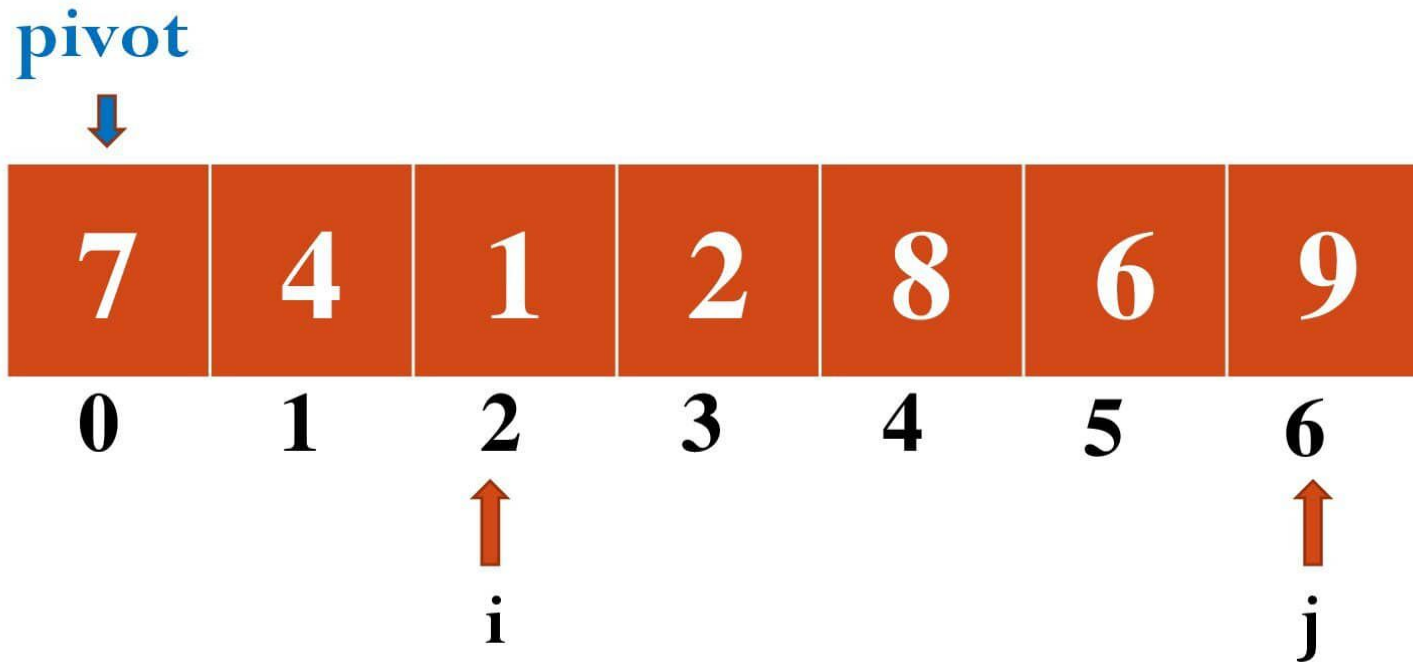
Quick Sort



Quick Sort

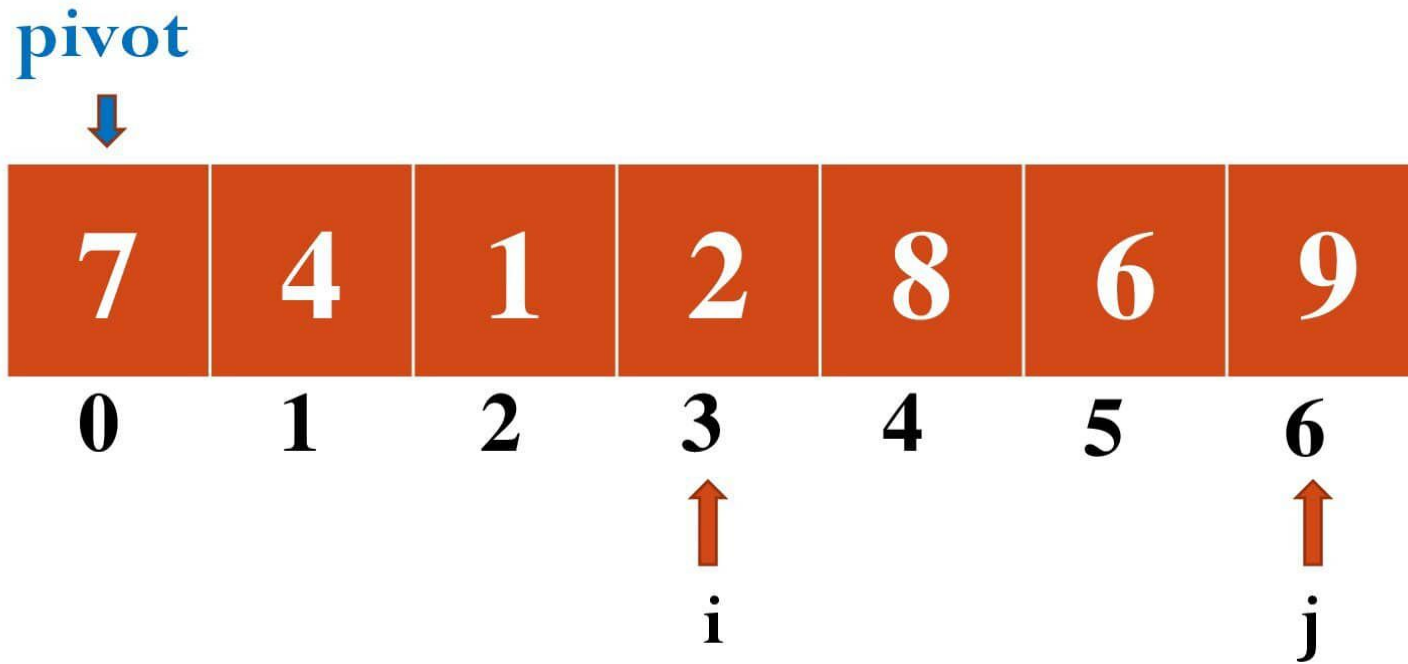


Quick Sort

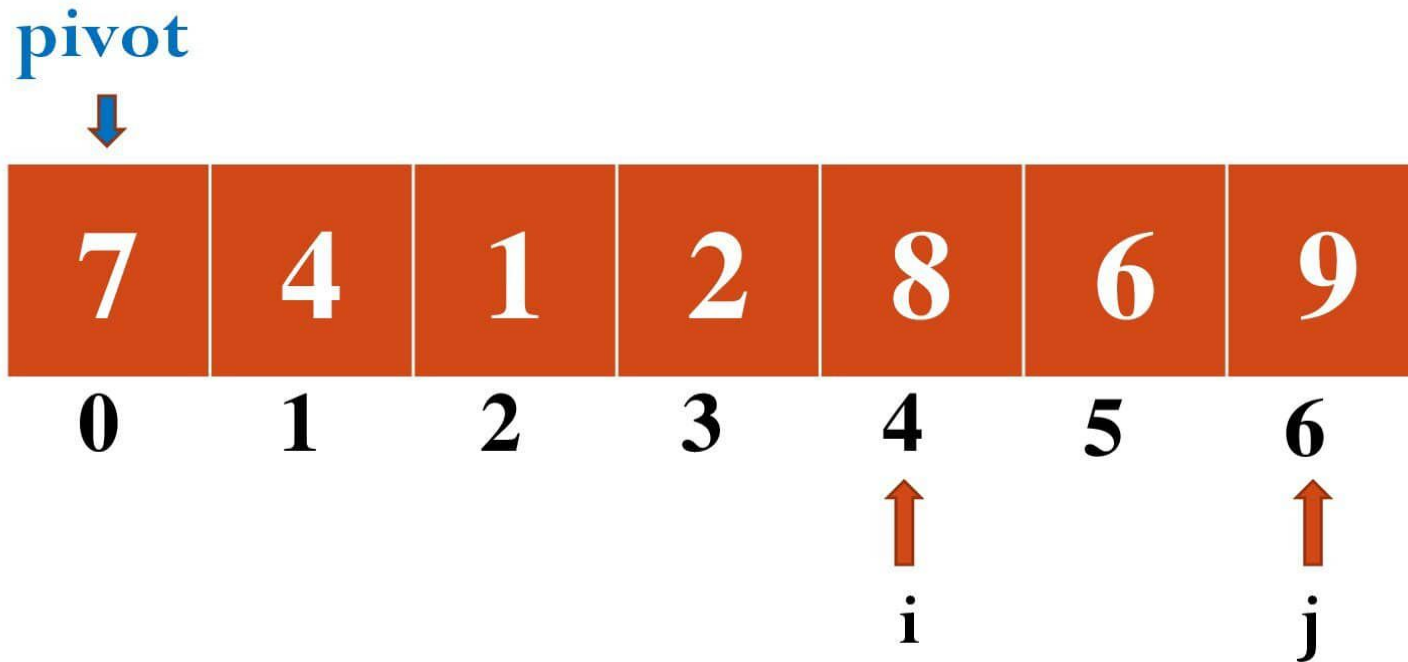


Swap $A[i]$ and $A[j]$

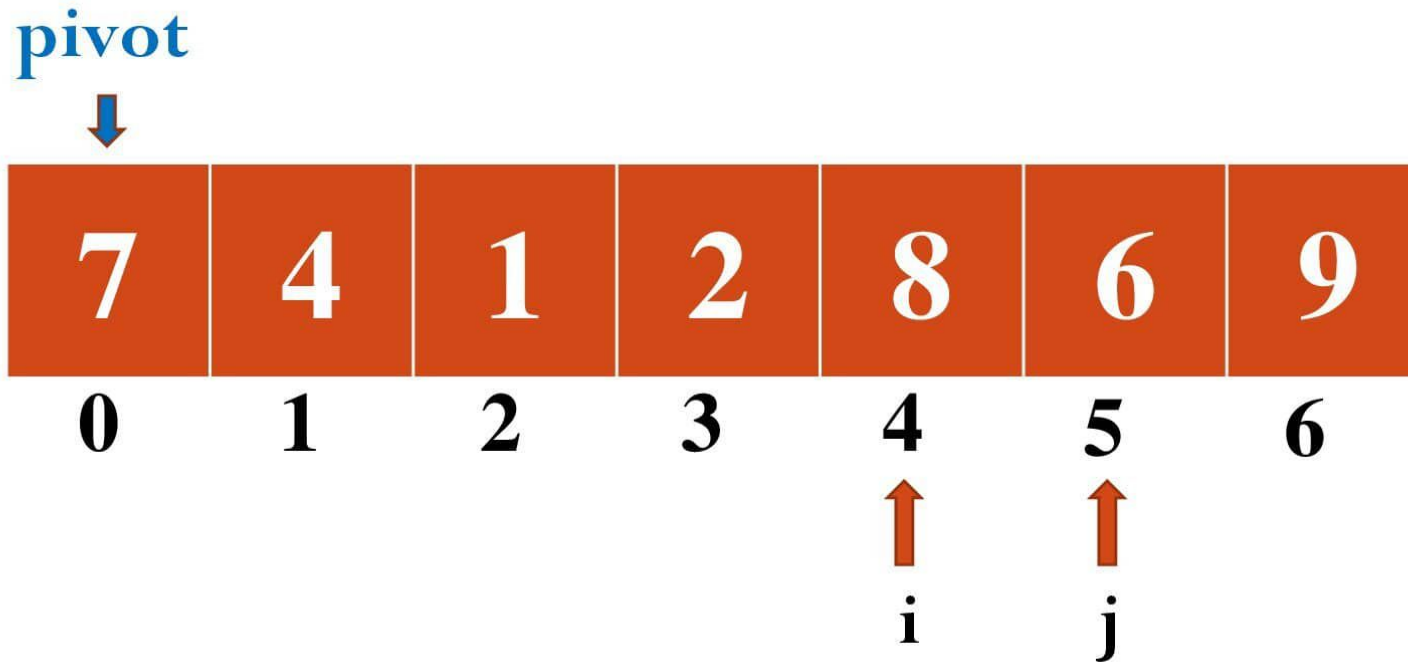
Quick Sort



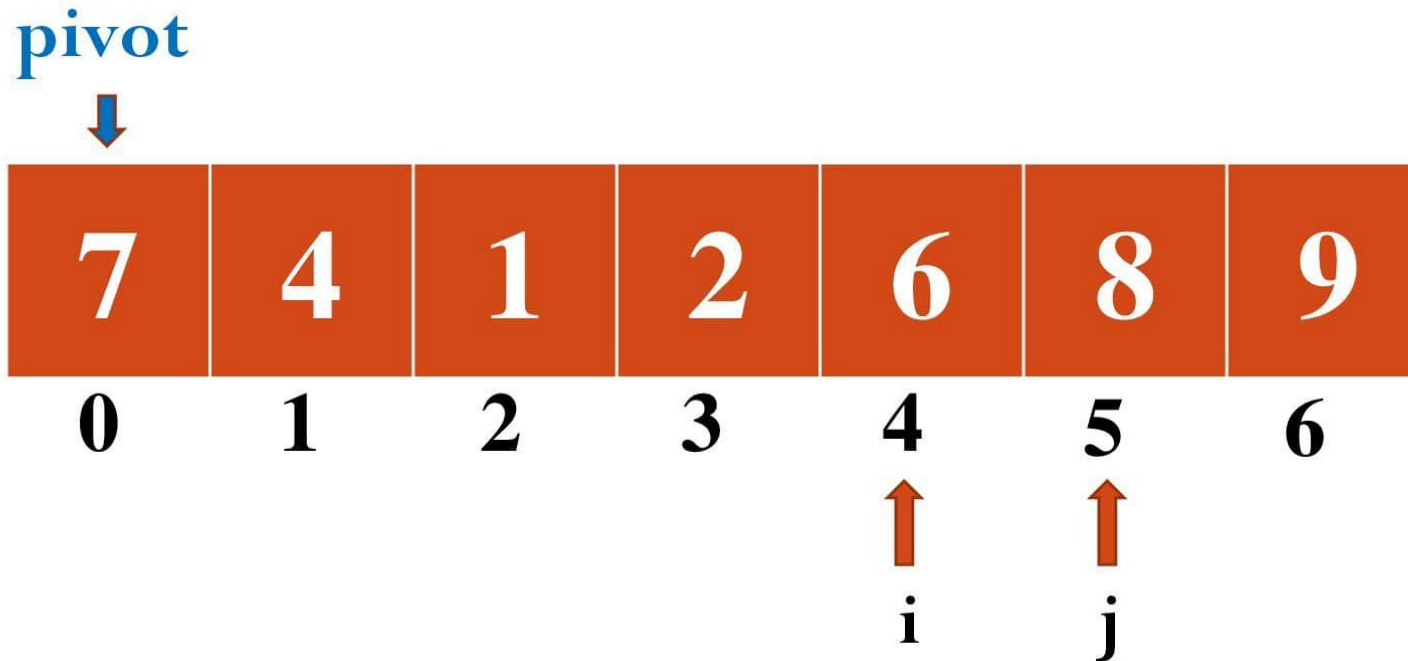
Quick Sort



Quick Sort

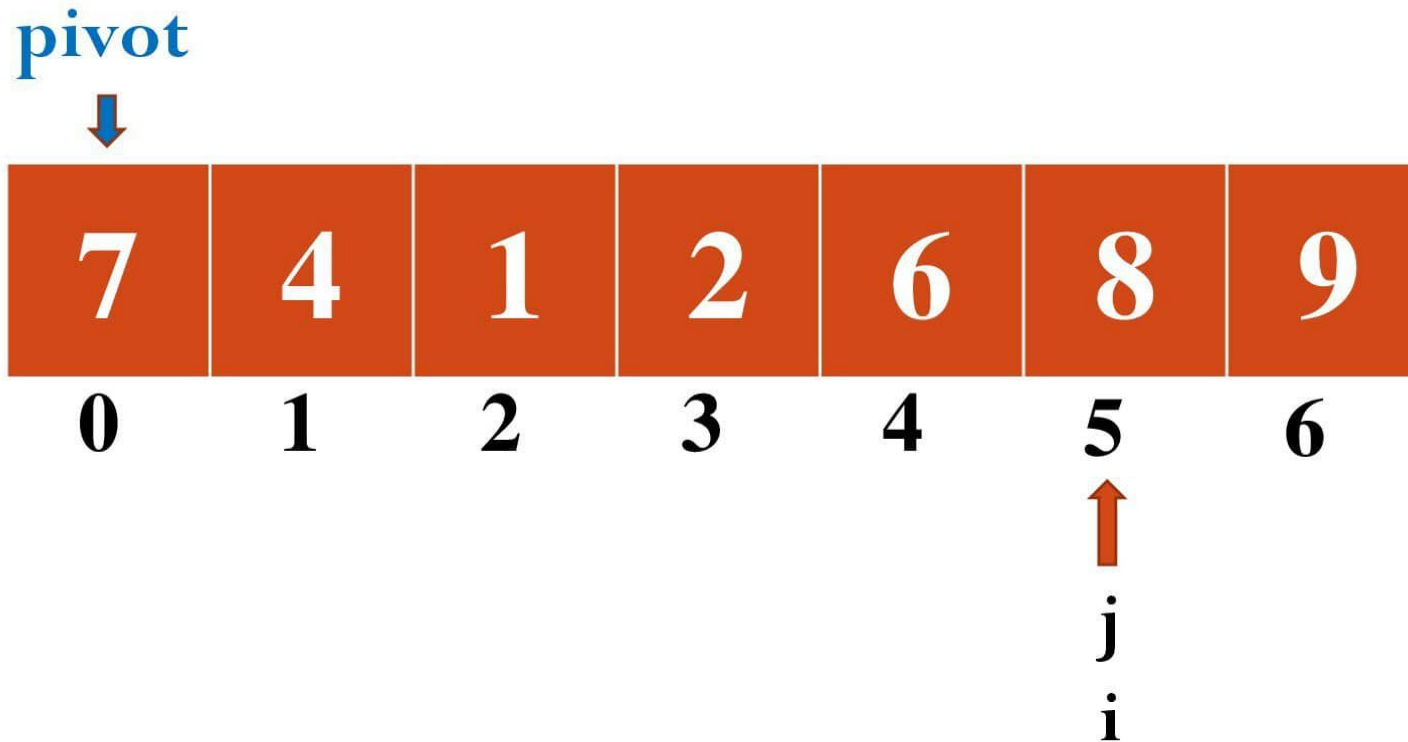


Quick Sort

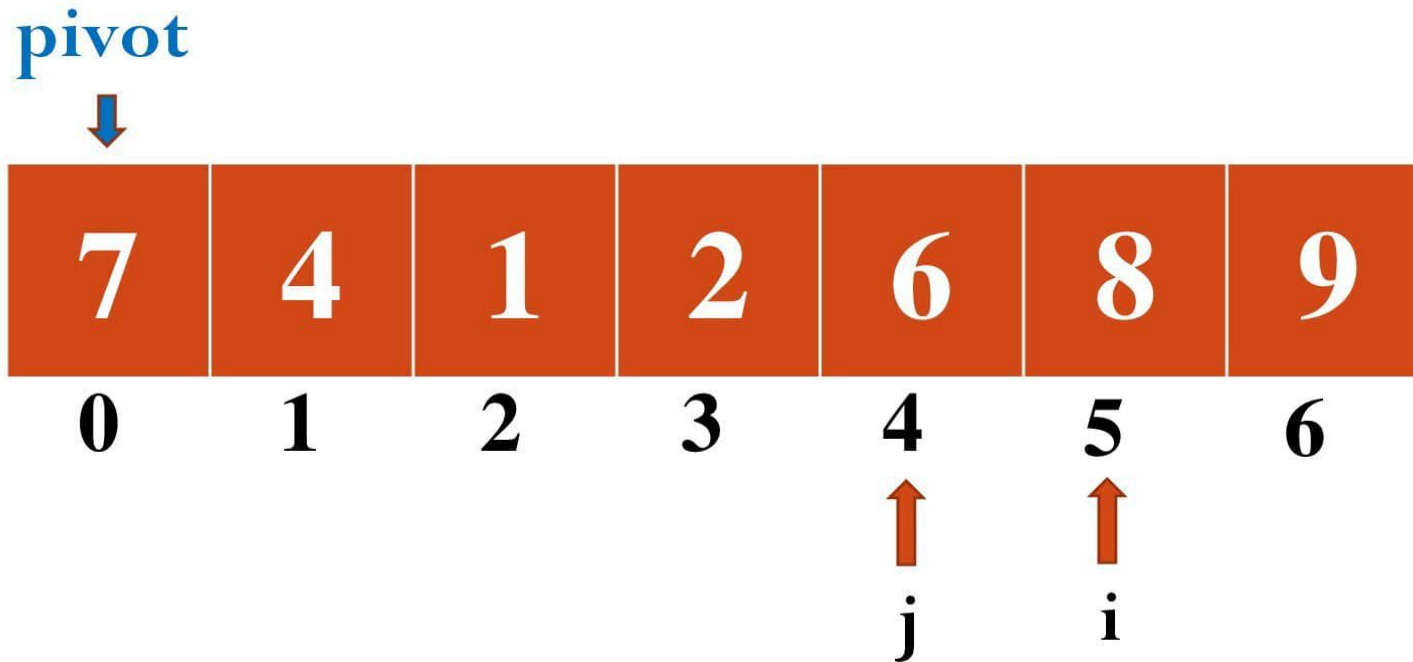


Swap $A[i]$ and $A[j]$

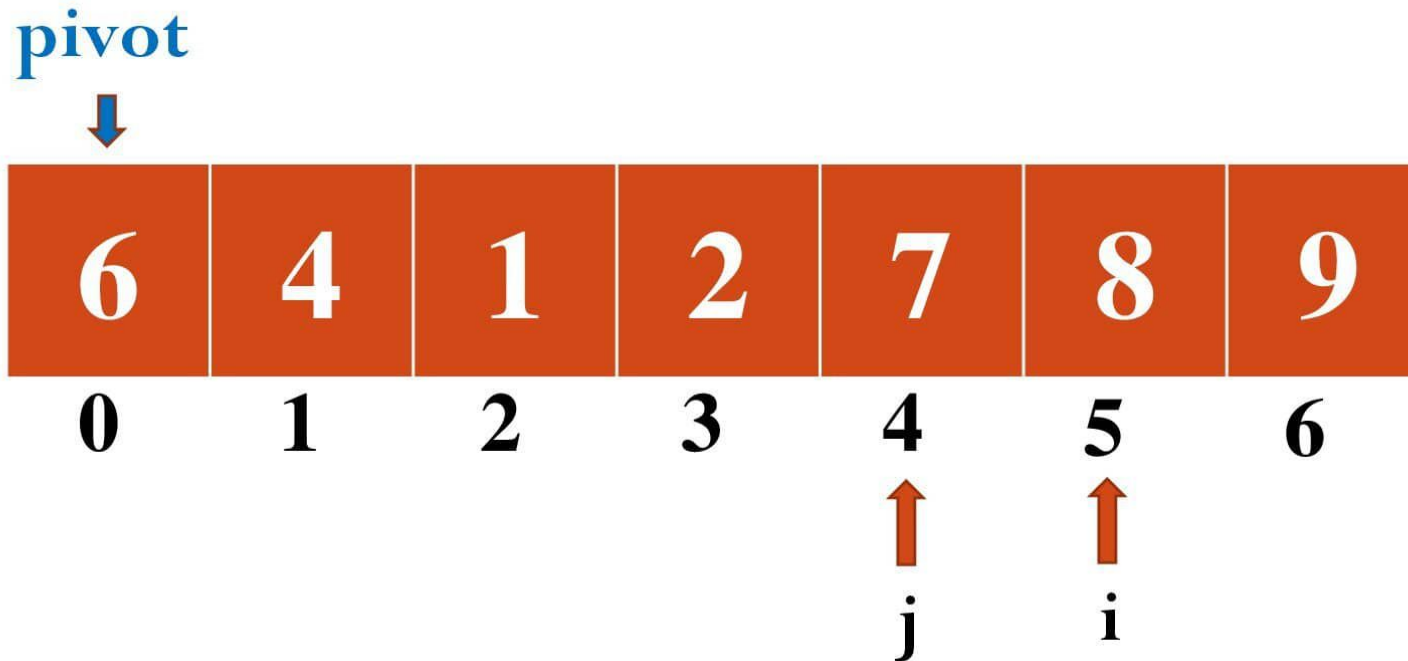
Quick Sort



Quick Sort



Quick Sort

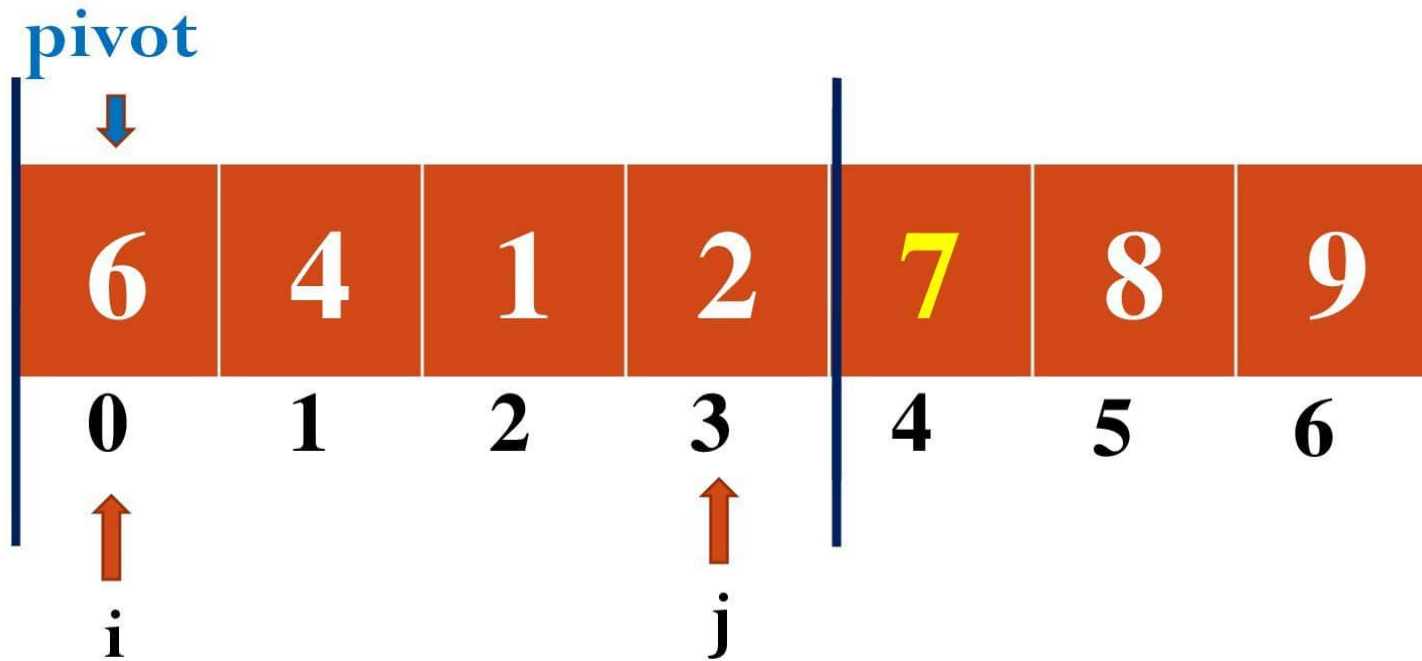


Swap $A[\text{pivot}]$ and $A[j]$

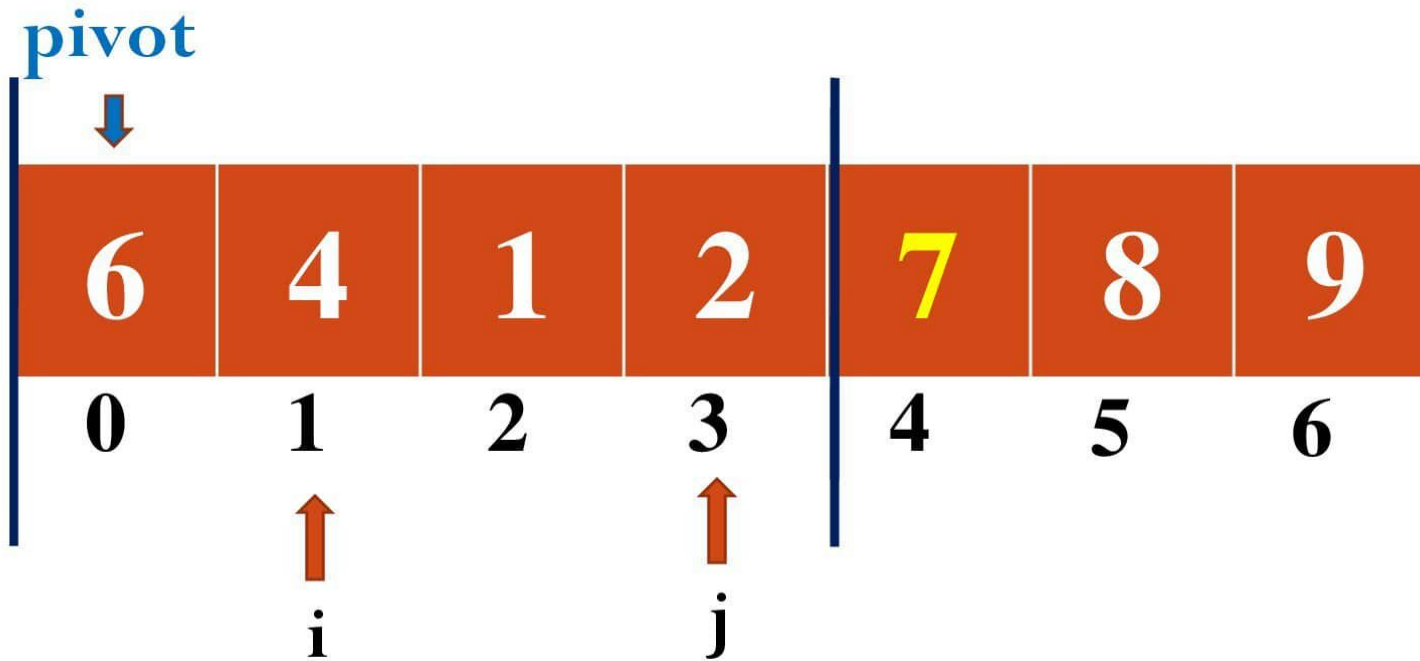
Quick Sort



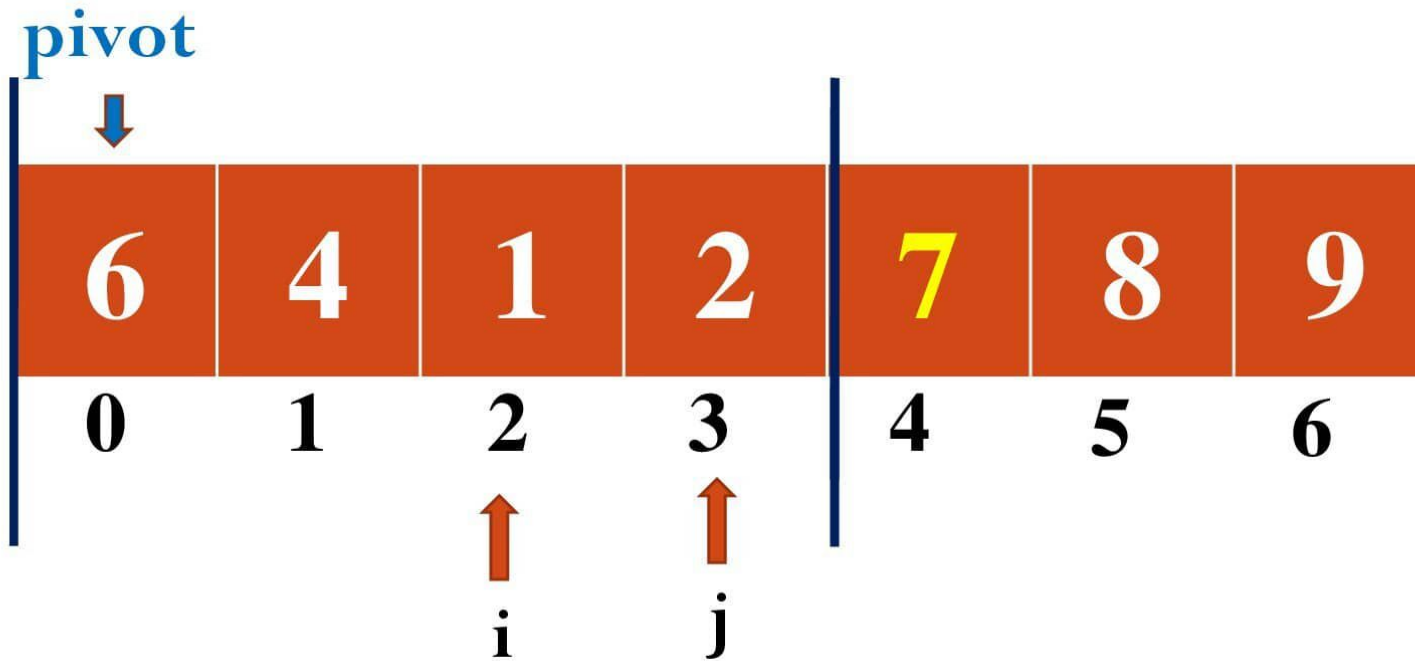
Quick Sort



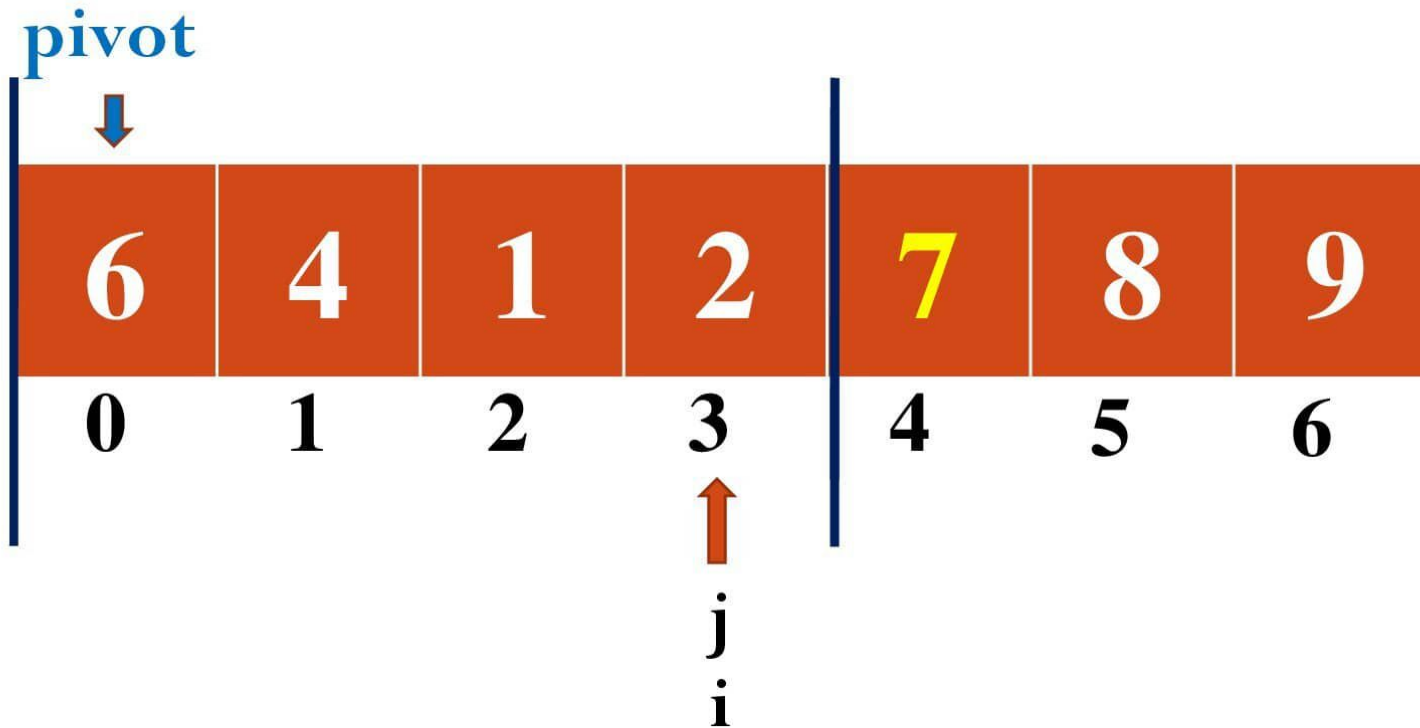
Quick Sort



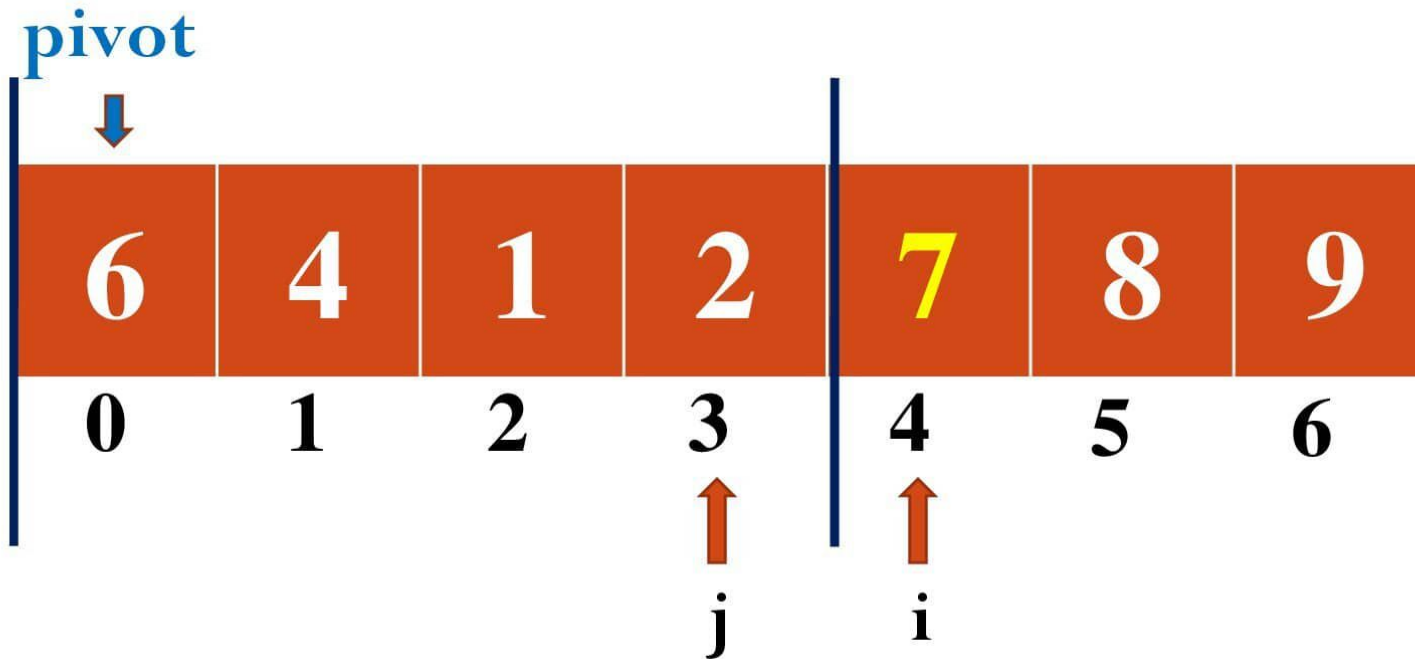
Quick Sort



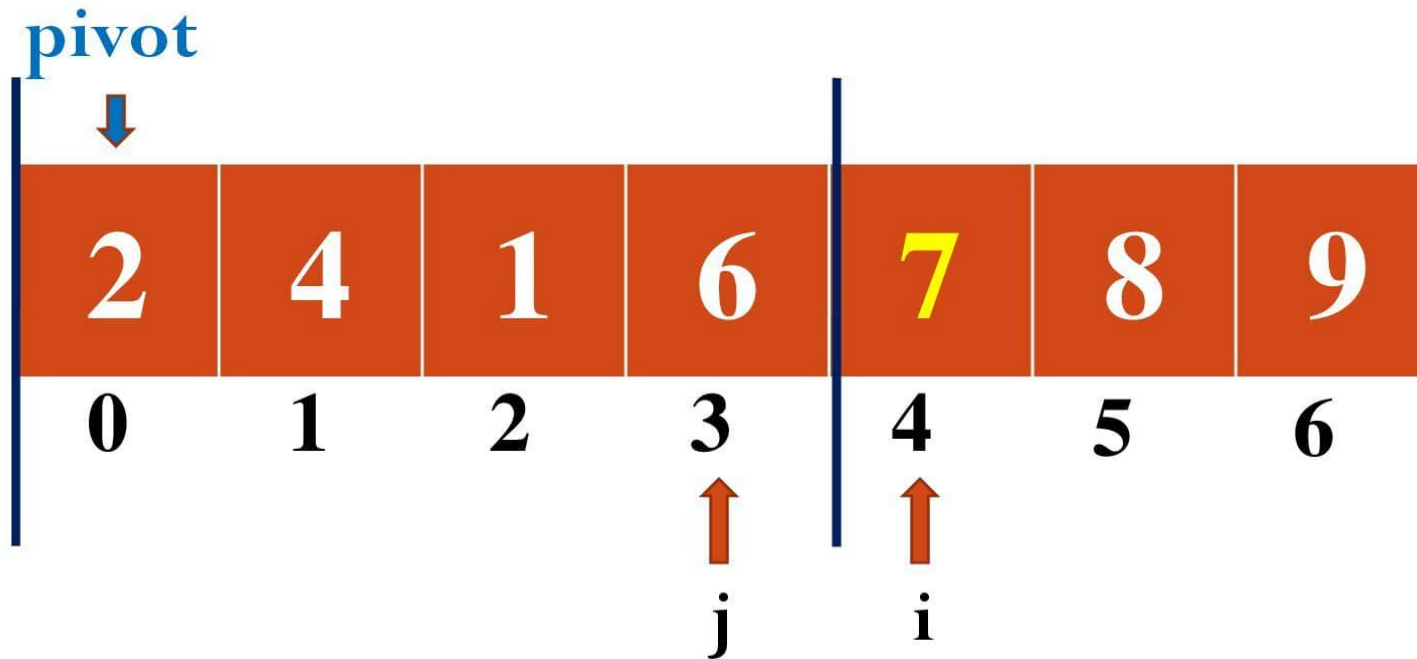
Quick Sort



Quick Sort

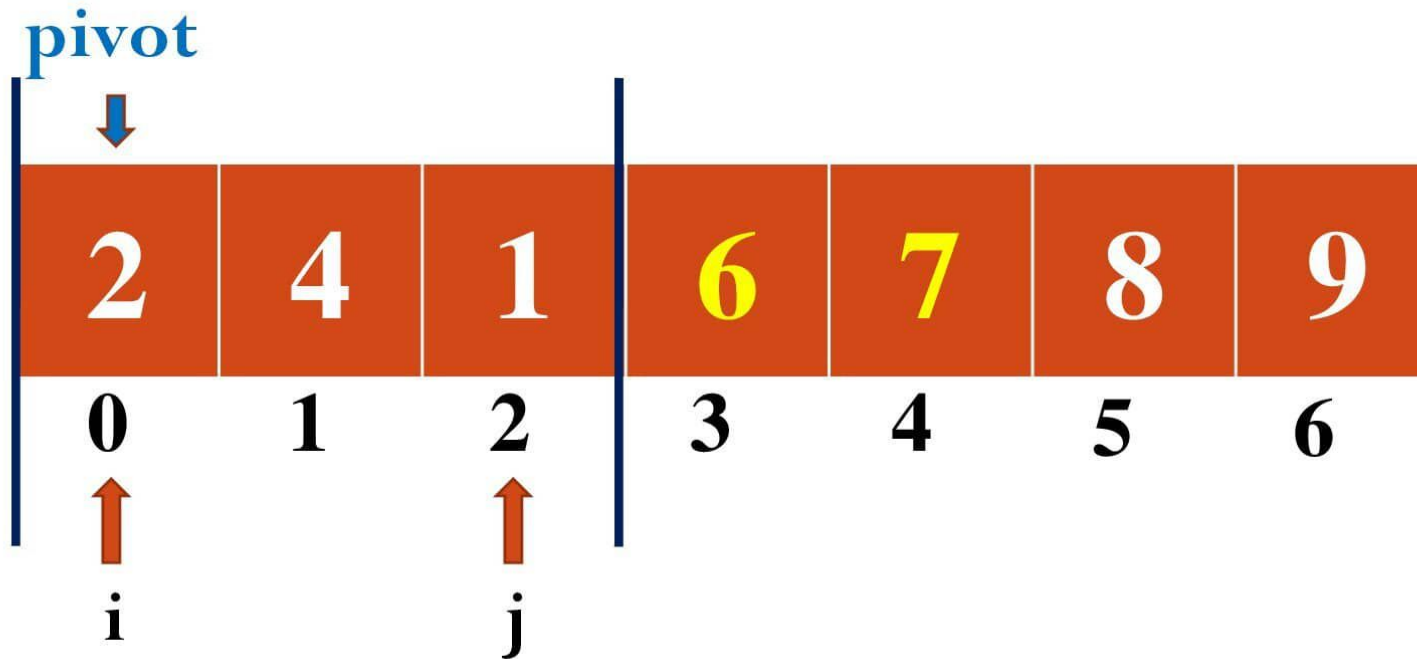


Quick Sort

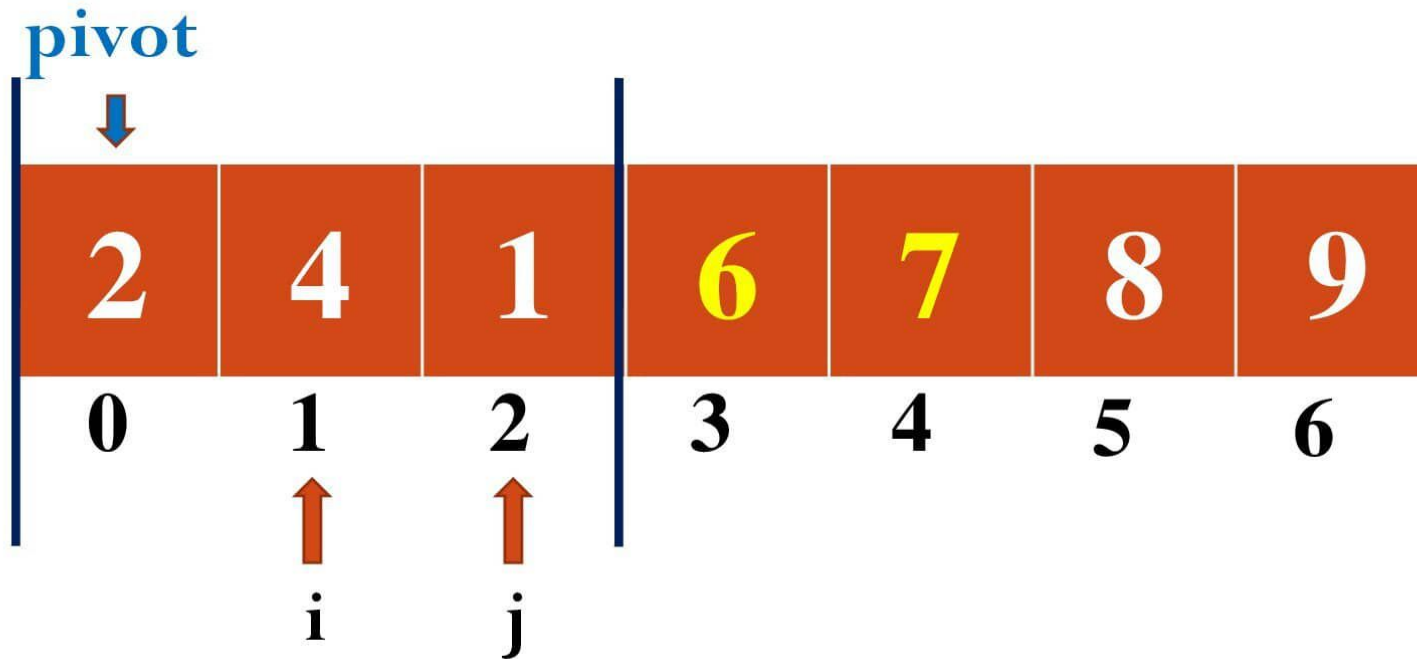


Swap $A[\text{pivot}]$ and $A[j]$

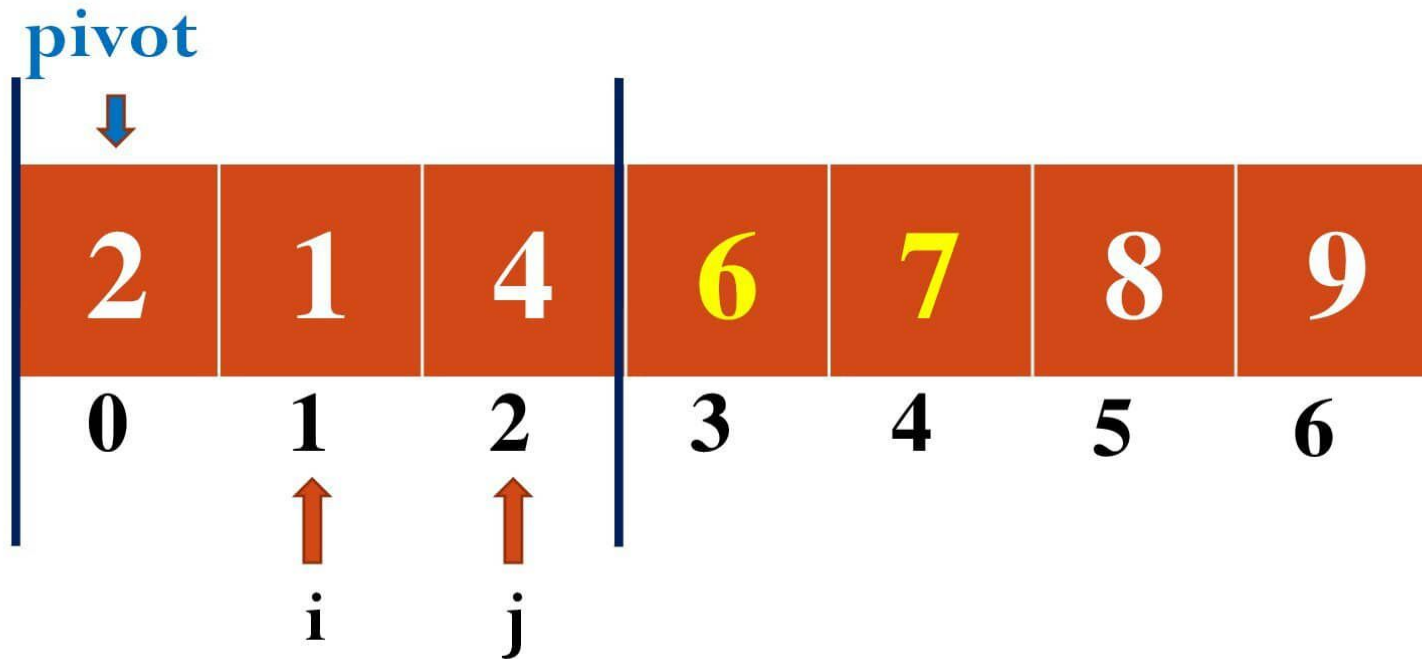
Quick Sort



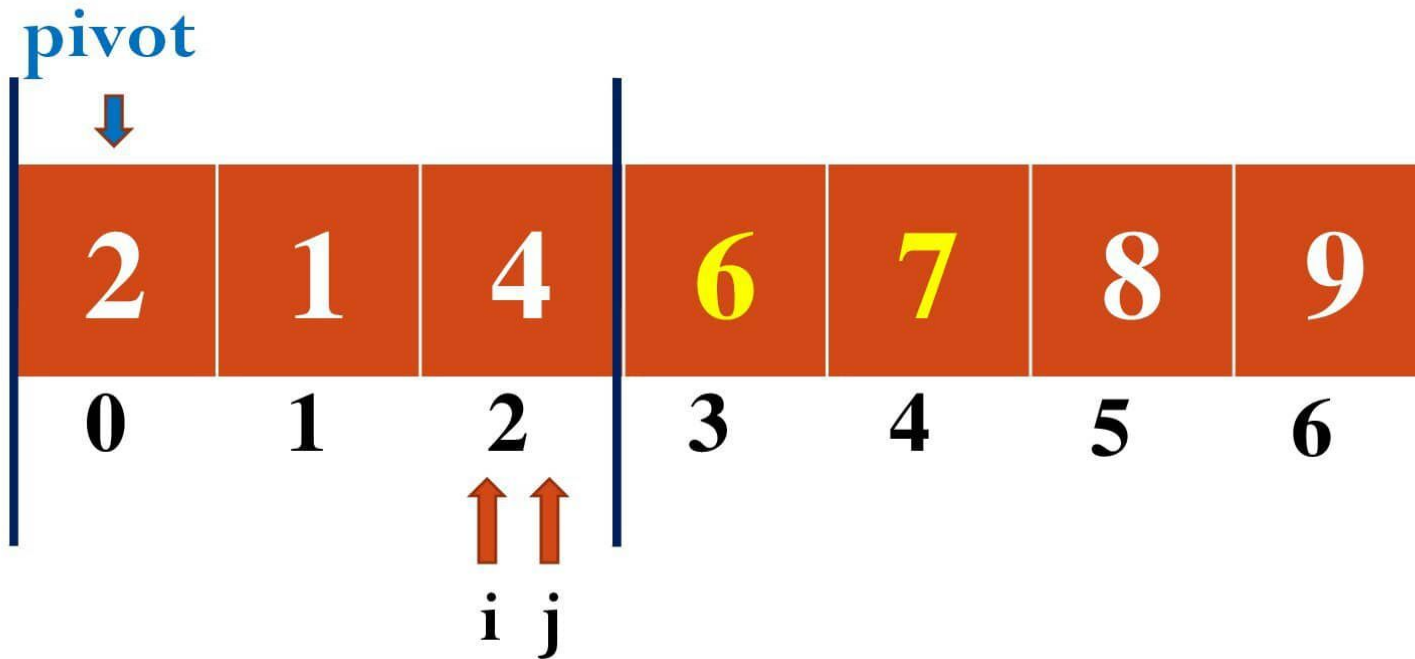
Quick Sort



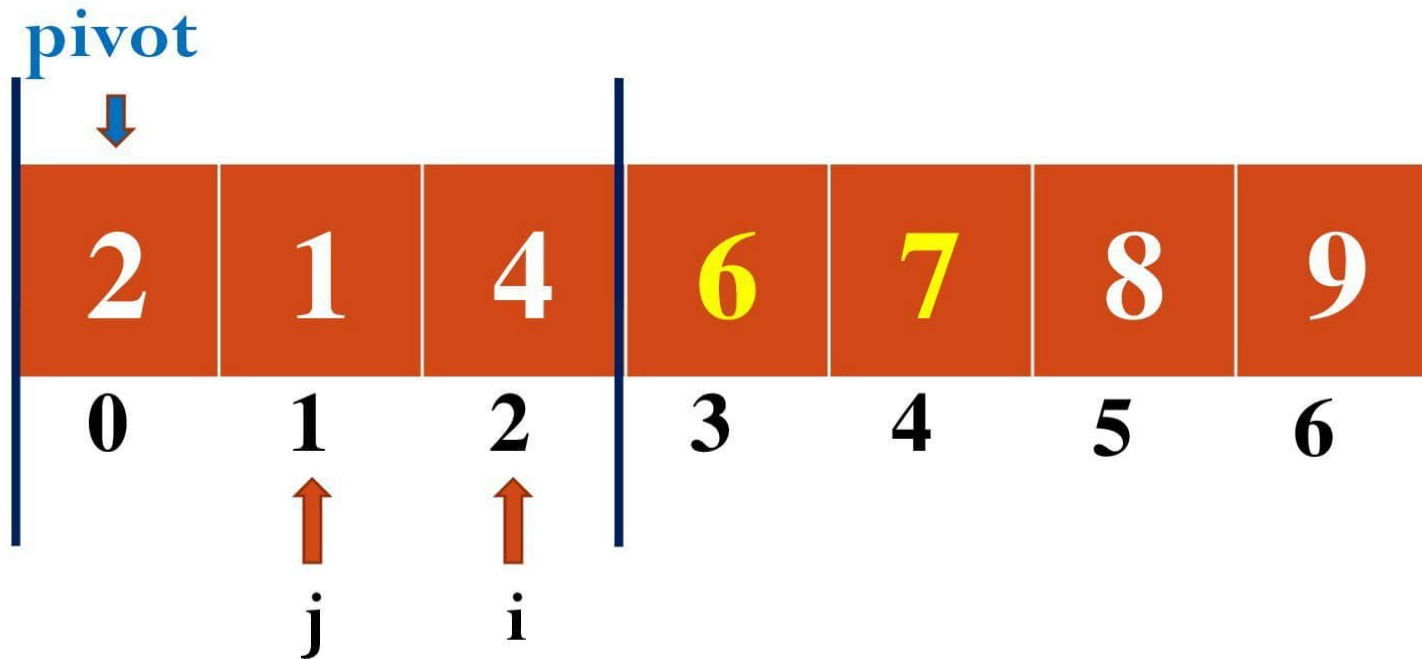
Quick Sort



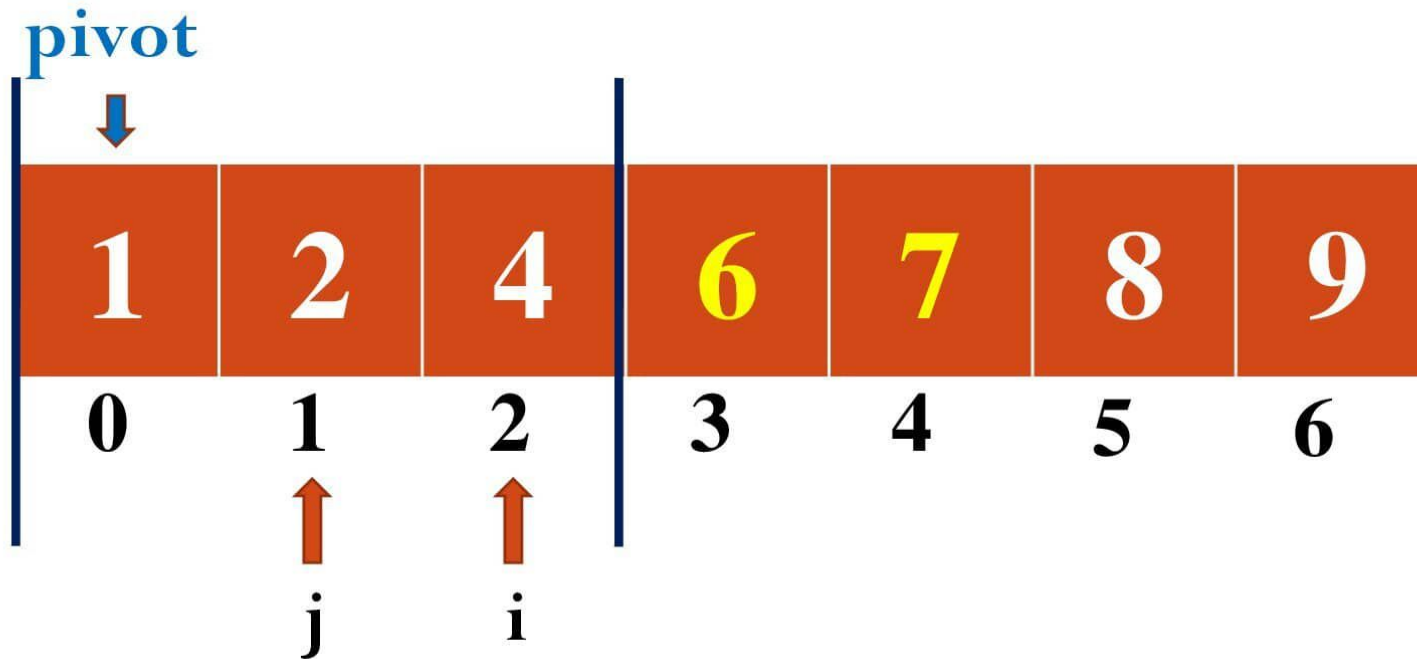
Quick Sort



Quick Sort

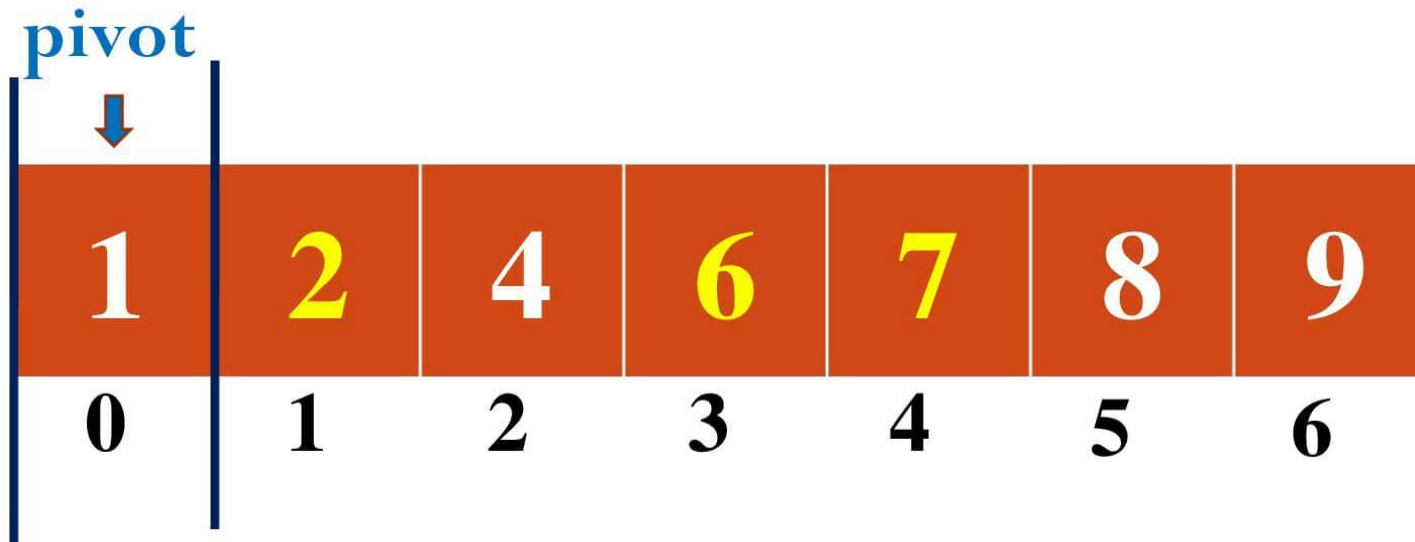


Quick Sort

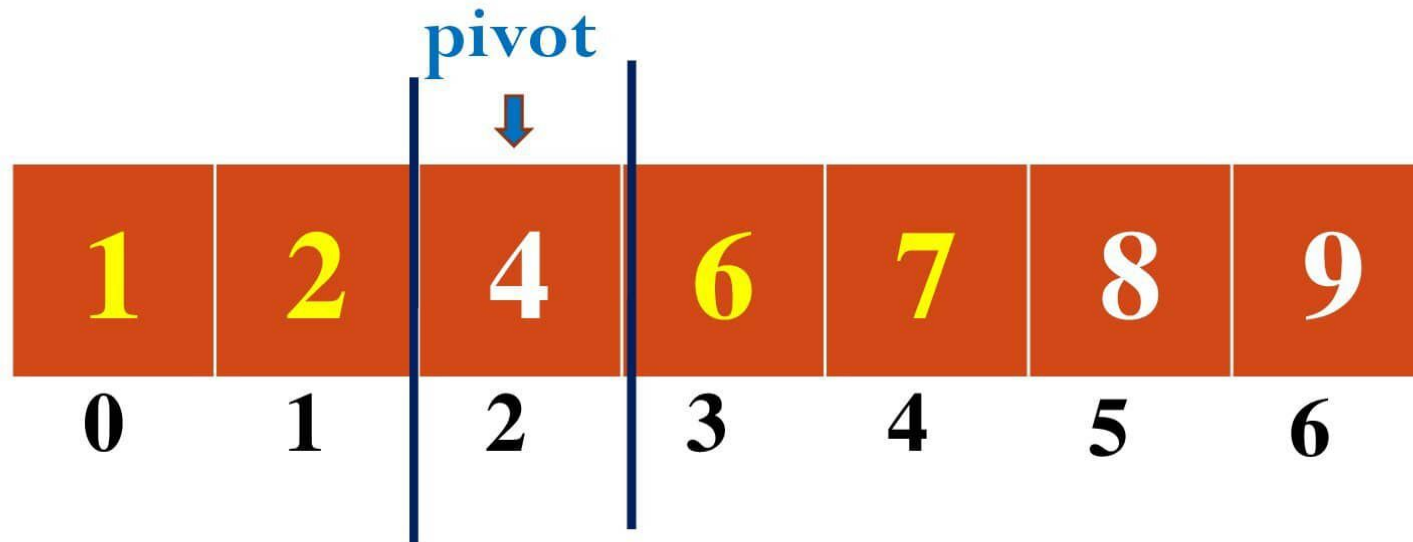


Swap pivot and A[j]

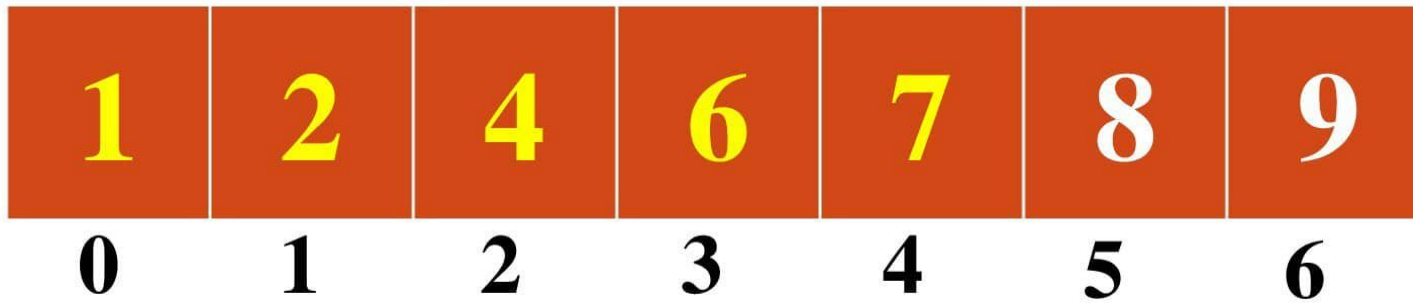
Quick Sort



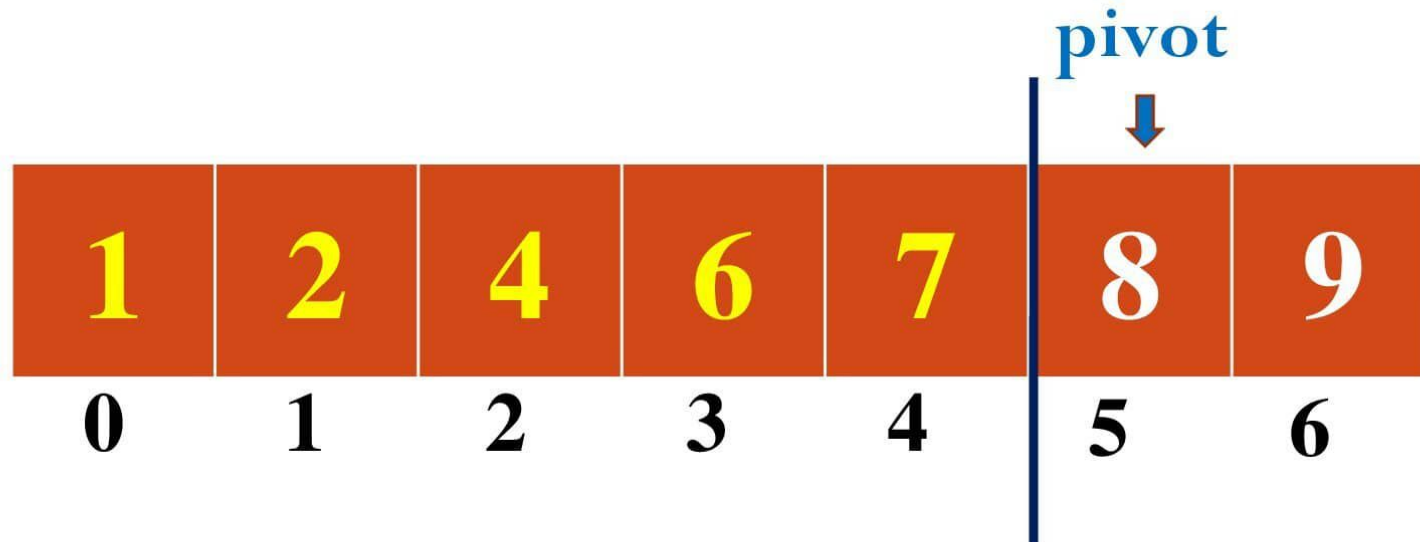
Quick Sort



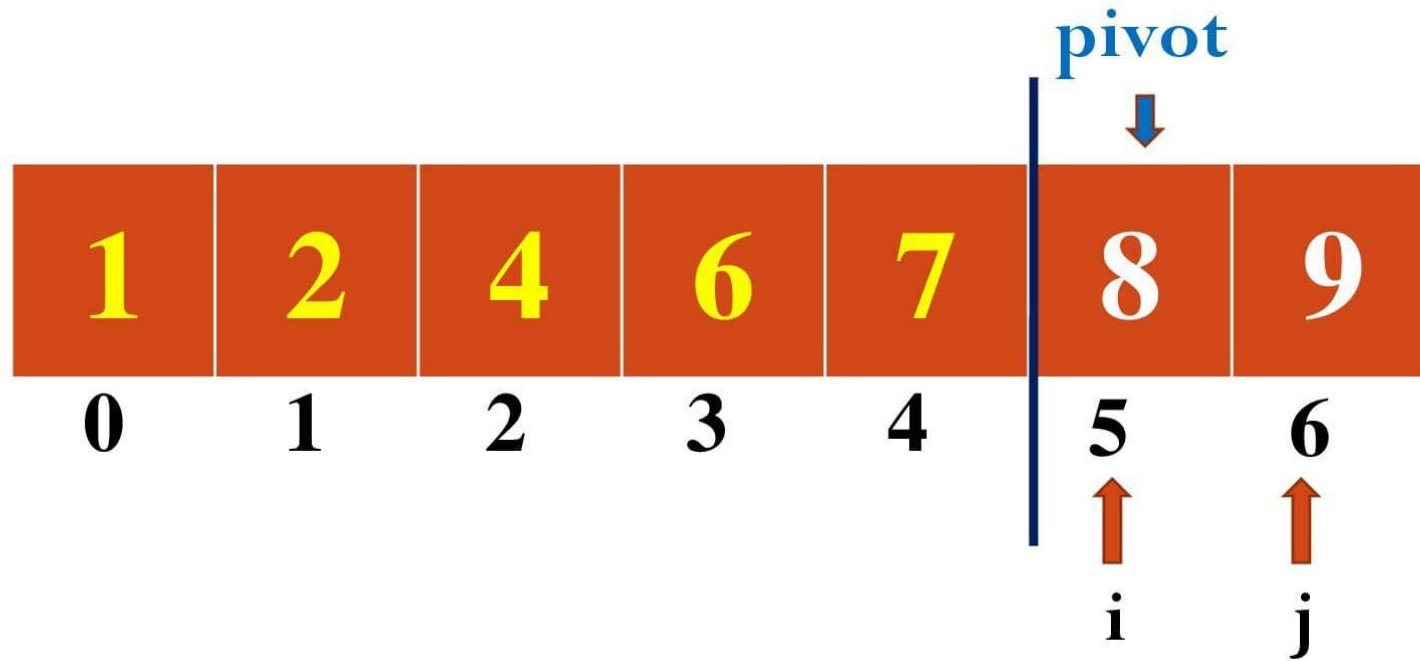
Quick Sort



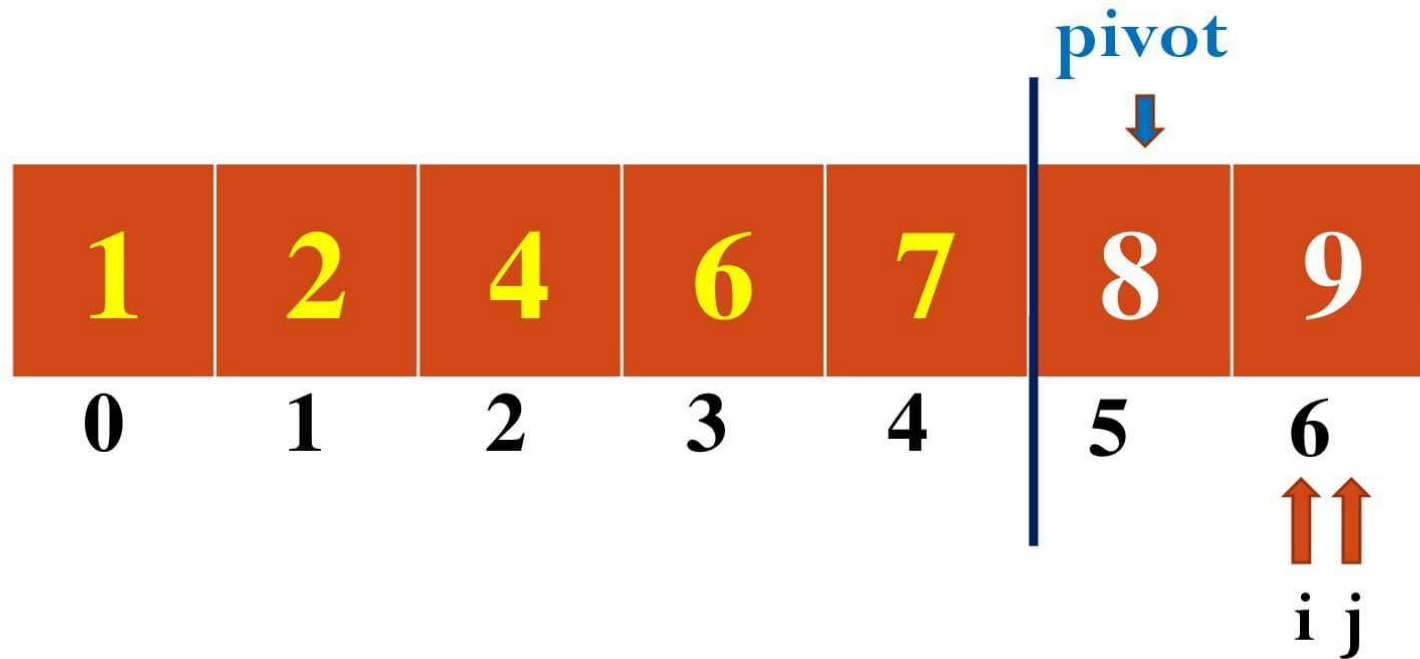
Quick Sort



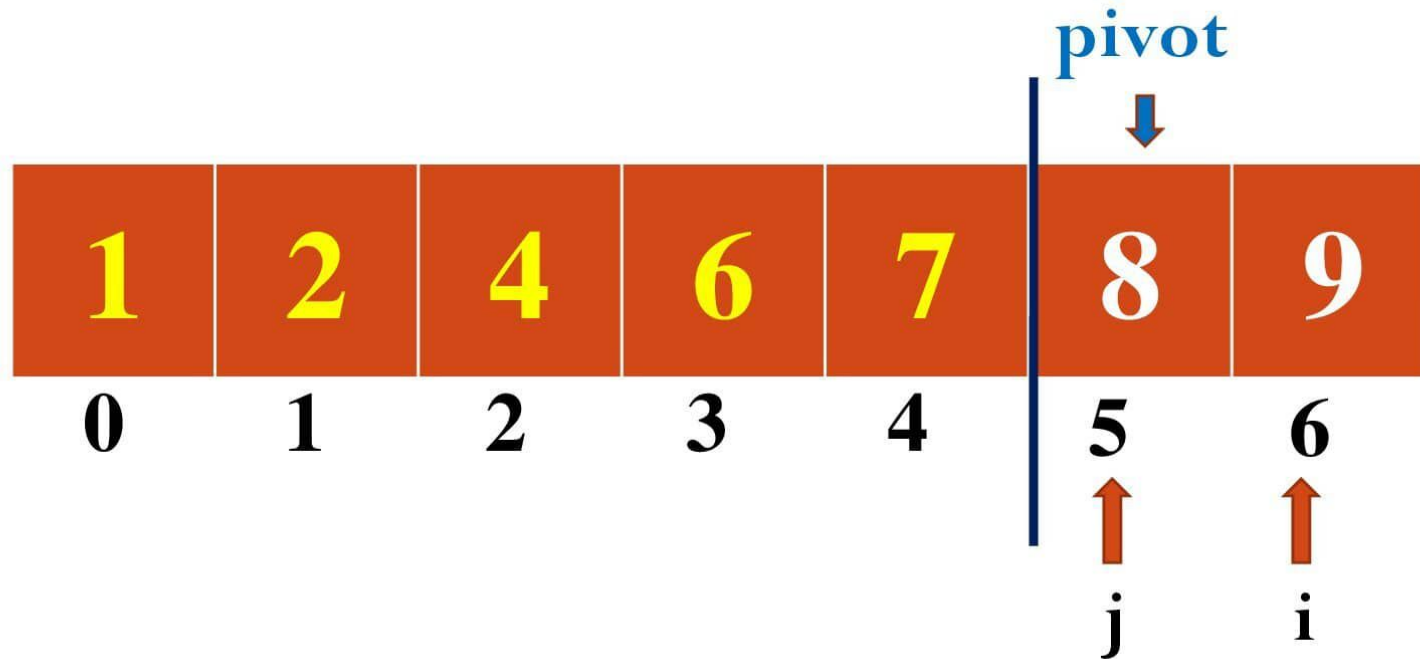
Quick Sort



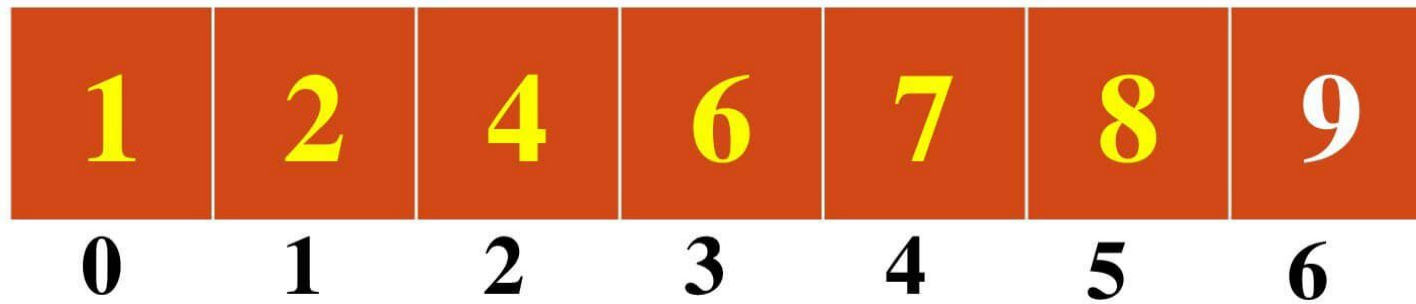
Quick Sort



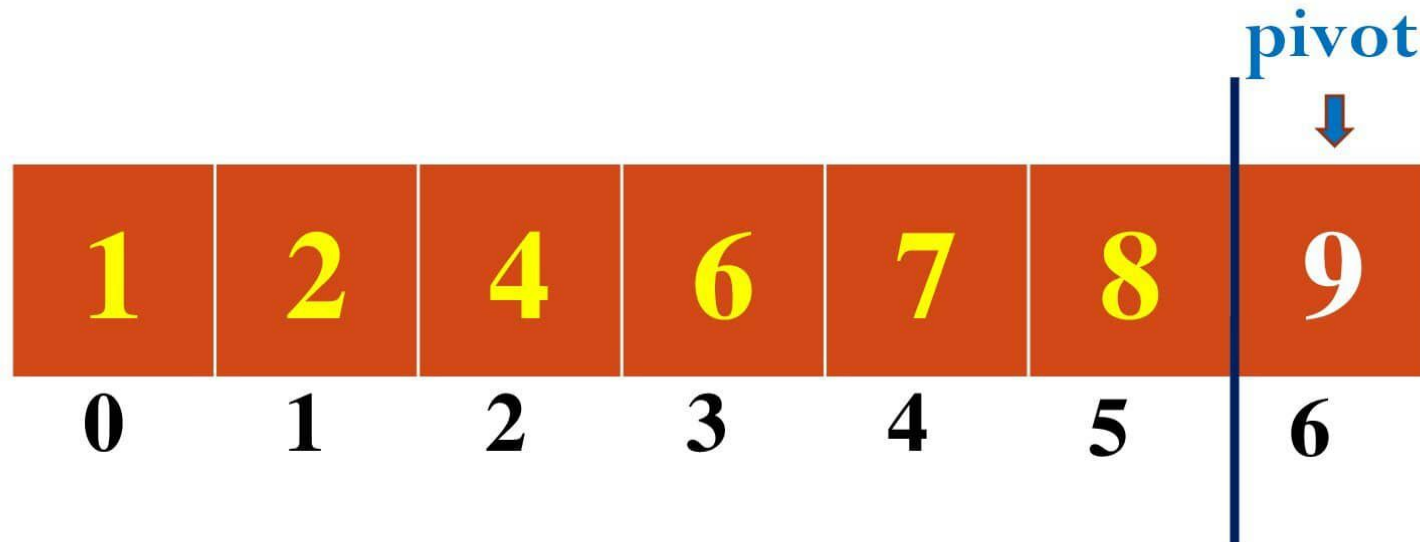
Quick Sort



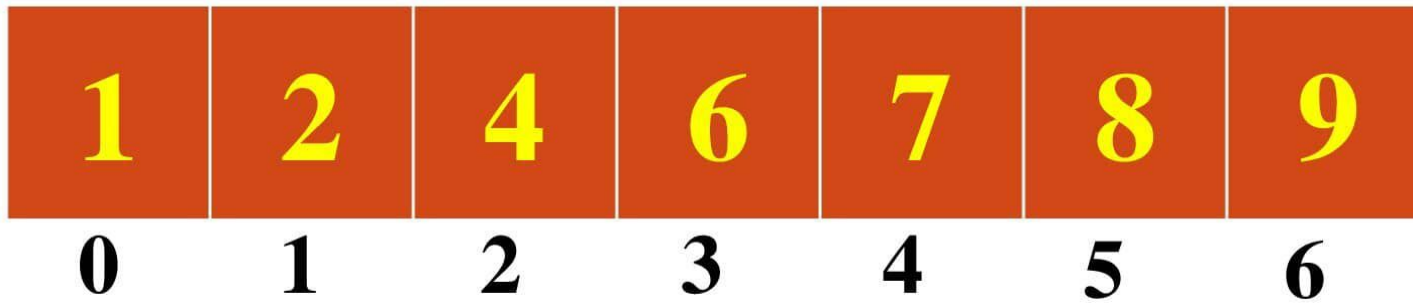
Quick Sort



Quick Sort



Quick Sort



Quick Sort

Algorithm QuickSort(A, low, high)

```
{  
    if low < high then  
    {  
        j = partition(A, low, high);  
        QuickSort(A, low, j-1);  
        QuickSort(A, j+1, high);  
    }  
}
```

Algorithm partition(A, low, high)

```
{  
    i=low, j=high, pivot=low  
    while i < j do  
    {  
        while A[i] ≤ A[pivot] do  
            i++;  
        while(A[j] > A[pivot])  
            j--;  
        if i < j then  
            Swap A[i] and A[j]  
    }  
    Swap A[j] and A[pivot]  
    return j  
}
```